

# Supplementary Appendices to Rationalizing Choice with Multi-Self Models

Attila Ambrus  
Harvard

Kareen Rozen  
Yale

This document contains supplementary appendices to “[Rationalizing Choice with Multi-Self Models](#)” by Ambrus and Rozen. The main paper is referenced throughout as AR.

## A Genericity

In this Appendix we formally establish that for any fixed scaling function  $\phi(\alpha)$  the property that an additive and scale-invariant aggregator is not degenerate holds generically.

Let  $F^*$  denote the set of additive and scale-invariant aggregators. In order to define a topology on  $F^*$ , we transform the latter set of aggregators to a convenient representation. Note that for a fixed scaling function, specifying the aggregated utilities of  $n$  alternatives for selves in the  $n$ -dimensional simplex determines the aggregated utilities of  $n$  alternatives for all possible selves over  $n$  alternatives, since any self is a scalar multiple of exactly one self from the simplex. Hence, with respect to a grand set of alternatives with three elements, there is a natural bijection  $\beta$  between additive and scale-invariant aggregators, and the set of pairs of operators

$$\Omega = (O_1, O_2 | O_1 : \Delta_2 \rightarrow \mathbb{R}^2; O_2 : \Delta_3 \rightarrow \mathbb{R}^3),$$

where  $O_1$  determines how a self’s utilities get aggregated in pairs, and  $O_2$  determines how a self’s utilities get aggregated in the triple.

Define metric  $d$  on  $\Omega$  such that the distance between  $(O_1, O_2)$  and  $(O'_1, O'_2)$  is defined as  $\max_{i=1,2} \sup_{x \in \mathbb{R}^i} |O_i(x) - O'_i(x)|$ .

**Theorem 1.** *Given the topology induced by  $d$ , the pairs of operators in  $\Omega$  that are associated with non-degenerate aggregators in  $F^*$  is open and dense relative to  $\Omega$ .*

*Proof.* For ease of exposition, let

$$\Gamma_1^l(f, v) = f(a, \{a, c\}, v) - f(c, \{a, c\}, v),$$

$$\Gamma_2^l(f, v) = f(a, \{a, b\}, v) - f(b, \{a, b\}, v) + f(b, \{b, c\}, v) - f(c, \{b, c\}, v),$$

$$\Gamma_2^l(f, v) = f(a, \{a, b, c\}, v) - f(b, \{a, b, c\}, v) + [f(a, \{a, b, c\}, v) - f(c, \{a, b, c\}, v)],$$

$$\Gamma_2^l(f, v) = [f(a, \{a, b\}, v) - f(b, \{a, b\}, v)] + [f(a, \{a, c\}, v) - f(c, \{a, c\}, v)],$$

for every  $v \in F^*$ . Note that  $\Gamma_i^j(v)$  stands for side  $j$  of the equation in condition  $i$  in the definition of a degenerate aggregator, given aggregator  $f$  and self  $v$ .

1. (*Openness*). Suppose that for aggregator  $f$  there is a self  $u$  over a triple such that neither of the equalities in the definition of a degenerate aggregator hold with equality. Note that  $u$  cannot be an indifferent self. Let  $\varepsilon_i = \Gamma_i^l(f, v) - \Gamma_i^r(f, v)$  for  $i \in \{1, 2\}$ , and let  $\varepsilon = \max(|\varepsilon_1|, |\varepsilon_2|)$ . Next, for every  $i, j \in \{a, b, c\}$  such that  $i \neq j$ , let  $\alpha^{ij}$  be such that  $\alpha^{ij}(u(i), u(j)) \in \Delta^2$ . Note that the terms  $\alpha^{ij}$  are uniquely defined. Similarly, let  $\alpha^{abc}$  be such that  $\alpha^{abc}(u(a), u(b), u(c)) \in \Delta^3$ . Let  $\alpha = \max(|\alpha^{ab}|, |\alpha^{ac}|, |\alpha^{bc}|, |\alpha^{abc}|)$ . Since  $u$  is not an indifferent self,  $\alpha > 0$ . Then for  $\delta < \frac{\varepsilon}{8\alpha}$  it holds that  $\Gamma_i^l(f', v) \neq \Gamma_i^r(f', v)$  for  $i \in \{1, 2\}$  for every  $f'$  such that  $|\beta(f) - \beta(f')| < \delta$ , since each term given  $f'$  in the above inequalities can differ from the corresponding term given  $f$  by at most  $\frac{\varepsilon}{8}$ .
2. (*Denseness*). Let  $\delta > 0$ . Consider a self  $u \in \Delta_3$  over  $\{a, b, c\}$  such that  $u(a) > u(b) > u(c)$ . For every  $i, j \in \{a, b, c\}$  such that  $i \neq j$ , let  $\alpha^{ij}$  be such that  $\alpha^{ij}(u(i), u(j)) \in \Delta^2$ . Let  $\alpha = \max(|\alpha^{ab}|, |\alpha^{ac}|, |\alpha^{bc}|)$ .

If for an aggregator  $f$  neither of the equalities in the definition of a degenerate aggregator hold, then the aggregator is by definition non-degenerate, hence there is trivially a point in the  $\delta$ -neighborhood of  $\beta(f)$  that corresponds to a non-degenerate aggregator. Otherwise let  $\varepsilon \in (0, \frac{\delta}{\alpha})$  be such that  $\varepsilon \neq |\Gamma_i^l(f, v) - \Gamma_i^r(f, v)|$  for  $i \in \{1, 2\}$ .

Consider now any  $f' \in F^*$  for which

- (a) For triples,  $f'$  is equivalent to  $f$
- (b) For a pair  $\{x, y\}$ , given any utility function  $v$  over  $\{x, y\}$  for which  $v(x) \geq v(y)$ ,  $f'(x, \{x, y\}, v) = f(x, \{x, y\}, v)$  and  $f'(y, \{x, y\}, v) = f(y, \{x, y\}, v)$  if  $v(x) - v(y) < u(a) - u(c)$ , but  $f'(x, \{x, y\}, v) = f(x, \{x, y\}, v) + \varepsilon$  and  $f'(y, \{x, y\}, v) = f(y, \{x, y\}, v)$  if  $v(x) - v(y) \geq u(a) - u(c)$ .

In words, with respect to selves for which the utility difference between the elements of the pair is at least  $u(a) - u(c)$  the aggregated utility is  $\varepsilon > 0$  higher than what  $f$  yields for the preferred alternative (while it is the same for the other alternative) - otherwise  $f'$  is equivalent to  $f$ . By construction,  $|\beta(f') - \beta(f)| < \delta$ . Also note that

$$\Gamma_1^l(f', v) = \Gamma_1^l(f, v) + \varepsilon,$$

$$\Gamma_1^r(f', v) = \Gamma_1^r(f, v),$$

$$\Gamma_2^l(f', v) = \Gamma_2^l(f, v),$$

and

$$\Gamma_2^r(f', v) = \Gamma_2^r(f, v) + \varepsilon.$$

Then  $\varepsilon \neq |\Gamma_i^l(f, v) - \Gamma_i^r(f, v)|$  for  $i \in \{1, 2\}$  implies that  $\Gamma_i^l(f', v) \neq \Gamma_i^r(f', v)$  for  $i \in \{1, 2\}$ . Hence,  $f'$  is non-degenerate. ■

## B Approximate triple-solvability

For some aggregators a tighter upper bound can be provided for the minimum number of selves needed to rationalize a choice function, through a weakening of the triple-solvability requirement. In particular, it suffices for triple-solvability to hold only *approximately*, which can yield a triple-basis with a smaller number of selves. For ease of exposition we only state this property for additively separable aggregators.

**Definition 1.** We say  $\hat{U} \in \mathcal{U}(\{a, b, c\})$  is a  $(\delta, \varepsilon)$ -approximate triple-basis for  $f$  with respect to  $\{a, b, c\}$  if  $f(a, \{a, b\}, \{a, b, c\}, \hat{U}) = f(b, \{a, b\}, \{a, b, c\}, \hat{U}) + \delta$  and

$$|f(x, A, \{a, b, c\}, \hat{U}) - f(y, A, \{a, b, c\}, \hat{U})| < \varepsilon$$

for all other  $A \subseteq \{a, b, c\}$  and  $x, y \in A$ .

That is, a collection of selves  $U$  is a  $(\delta, \varepsilon)$ -approximate triple basis for  $f$  if given choice set  $\{a, b\}$  the aggregated utility of  $U$  for  $a$  is exactly  $\delta$  higher than the aggregated utility of  $b$ , while  $U$  is  $\varepsilon$ -indifferent among all alternatives given every other choice set.

We say that an aggregator  $f$  is *approximately triple-solvable with  $k$  selves* if there is  $\bar{\delta} > 0$  such that exists a  $(\delta, \varepsilon)$ -approximate triple-basis with  $k$  selves for every  $\delta < \bar{\delta}$  and  $\varepsilon > 0$ . That is, for approximate triple-solvability we do not require that the collection of selves in the triple is exactly indifferent between all elements in choice sets other than  $\{a, b\}$ , only that they can be arbitrarily close to being indifferent.

AR-Theorem 1 can then be modified as follows.

**Theorem 2.** Suppose  $f$  satisfies P1-P6 and P9, and is approximately triple-solvable with  $k_f$  selves. Then, for any finite set of alternatives  $X$ , and any choice function  $c : P(X) \rightarrow X$  that exhibits at most  $\frac{n-1}{k_f}$  IIA-violations,  $f$  can rationalize  $c$  with  $n$  selves.

*Proof.* The only difference compared to the proof of [AR-Theorem 1](#) is in the construction of selves. Recall the definition of  $(I_j)_{j=1,\dots,j^*}$  from the proof of [AR-Theorem 1](#). Let  $\delta_1 \in (0, \bar{\delta})$ . Define iteratively  $\delta_j$  for  $j \in \{2, \dots, j^* + 1\}$  such that  $\delta_j \in (0, \frac{\delta_{j-1}}{IIA(c)+1})$ . Define a self  $u^X$  such that  $u^X$  is  $\delta_{j^*+1}$ -indifferent and the preference ordering of the self is  $c(X) \succ c(X \setminus \{c(X)\}) \succ \dots$ . Let

$$\delta^{**} = \min_{x \neq y \in X, A \ni x, y} |f(x, A, X, u^X)| - |f(y, A, X, u^X)|.$$

Finally, let  $\varepsilon \in (0, \frac{\delta^{**}}{|X|})$ . Then for every  $j \in \{1, \dots, j^*\}$  and  $A \in I_j$  construct a collection of selves  $U^A \in \mathcal{U}(X)$  the following way: take a  $(\delta_j, \varepsilon)$ -approximate triple-basis  $U$ , and define  $U^A$  by defining, for each  $u_i \in U$ , a self  $u_i^A \in U^A$  by

$$u_i^A(x) = \begin{cases} u_i(a) & x = c(A) \\ u_i(b) & x \in A \setminus \{c(A)\} \\ u_i(c) & x \in X \setminus A. \end{cases}$$

Proving the collection of selves consisting of  $u^X$  and  $U^A$  for each  $A \in \bigcup_{j=1}^{j^*} I_j$  rationalizes  $c$  is analogous to the proof in [AR-Theorem 1](#). ■

## C Relaxing P6

Our main results can be extended to aggregators violating P6, that is, to aggregators that depend in a nontrivial way on alternatives unavailable in a given choice set. However, the appropriate definition of triple-solvability is more complicated.

The main complication arising in the absence of P6 is that triple-solvability needs to be defined on a general  $X$ , as opposed to just a triple  $\{a, b, c\}$ . It is convenient to introduce the following notation: for any triple  $\{a, b, c\}$ , any basic set of alternatives  $X \supset \{a, b, c\}$ , and any self  $u$  defined on  $\{a, b, c\}$ , define the set  $E(u, X) = \{\hat{u} : X \rightarrow \{u(a), u(b), u(c)\} | \hat{u}(x) = u(x) \forall x \in \{a, b, c\}\}$ . In words,  $E(u, X)$  is the set of extensions of  $u$  from  $\{a, b, c\}$  to  $X$  for which each element in  $X \setminus \{a, b, c\}$  receives the same utility as either  $a$  or  $b$  or  $c$ . Similarly, for any  $U = (u_1, \dots, u_m) \in \mathcal{U}(\{a, b, c\})$ , let  $E(U, X) = \{(\hat{u}_1, \dots, \hat{u}_m) | \hat{u}_i \in E(u_i, X) \text{ for all } i \in \{1, \dots, m\}\}$ .

**Definition 2.** We say  $U \in \mathcal{U}(\{a, b, c\})$  is a universal triple-basis for  $f$  if for any  $X \supset \{a, b, c\}$  the following holds: for all  $\hat{U} \in E(U, X)$ ,  $f(a, \{a, b\}, X, \hat{U}) > f(b, \{a, b\}, X, \hat{U})$ , and  $f(\cdot, A, X, \hat{U})$  is constant for all other  $A \subseteq \{a, b, c\}$ .

A universal triple-basis solves the triple  $\{a, b, c\}$  whenever the utilities of unattainable elements

don't differ from utilities of elements in  $\{a, b, c\}$ , for all selves in the triple-basis. An aggregator  $f$  is *universally triple-solvable* if the following condition is satisfied.

**Condition** (Universal triple-solvability of  $f$ ) There exists a triple  $\{a, b, c\}$  and  $k \in \mathbb{Z}_+$  such that for every  $\delta > 0$  there is a  $\delta$ -indifferent  $U \in \mathcal{U}^k(\{a, b, c\})$  constituting a universal triple-basis for  $f$  with respect to  $\{a, b, c\}$ .

It is easy to see that for aggregators satisfying P6, universal triple-solvability is equivalent to triple-solvability. If  $f$  satisfying P1-P5 is universally triple-solvable with  $k$  selves, then the same construction can be applied as in the proof of [AR-Theorem 1](#) to obtain an analogous lower bound on the set of choice functions that  $f$  can rationalize with a given number of selves. The proof of this result is analogous to the proof of [AR-Theorem 1](#) and hence omitted.

**Theorem 3.** *Suppose  $f$  satisfies P1-P5 and is universally triple-solvable wrt to  $X$  with  $k_f$  selves. Then, using  $n$  selves,  $f$  can rationalize any choice function, on any grand set of alternatives  $X$ , that exhibits at most  $\frac{n-1}{k_f}$  IIA-violations.*