

Beckmann's approach to multi-item multi-bidder auctions*

Alexander Kolesnikov[†]

Fedor Sandomirskiy[‡]

Aleh Tsyvinski[§]

Alexander P. Zimin[¶]

Abstract

We consider the problem of revenue-maximizing Bayesian auction design with several i.i.d. bidders and several items. We show that the auction-design problem can be reduced to the problem of continuous optimal transportation introduced by Beckmann (1952). We establish the strong duality between the two problems and demonstrate the existence of solutions. We then develop a new numerical approximation scheme that combines multi-to-single-agent reduction and the majorization theory insights to characterize the solution.

*We are grateful to Federico Echenique, Luciano Pomatto, Joseph Root, and Omer Tamuz for conversations that inspired this work. We thank Kim Border, Benjamin Brooks, Sergiu Hart, Jason Hartline, Andreas Kleiner, Alejandro Manelli, and Robert McCann for useful discussions.

[†]HSE University. Alexander Kolesnikov acknowledges the support of RSF Grant №22-21-00566 <https://rscf.ru/en/project/22-21-00566/>. The article was prepared within the framework of the HSE University Basic Research Program.

[‡]Caltech and HSE University. Fedor Sandomirskiy thanks Linde Institute at Caltech, National Science Foundation (grant CNS 1518941), and the HSE University Basic Research Program for their support.

[§]Yale.

[¶]MIT and HSE University.

1 Introduction

The current understanding of multi-bidder multi-item revenue-maximizing auctions is far from being complete even in the basic setting of several bidders competing for several items and having i.i.d. additive utilities over them. Only the case of one item and several bidders was analyzed completely (Myerson, 1981). A seemingly innocent problem with one bidder and several items already turns out to be notoriously difficult to analyze, optimal mechanisms are known only in a few particular cases and exhibit complicated structure (Armstrong and Rochet, 1999; Rochet and Stole, 2003; Daskalakis, 2015). Essentially, nothing was known about optimal auctions, where both the number of bidders and the number items exceed one.

The main contribution of our paper is an unexpected connection between the problem of auction design and an optimal transportation problem in the classic model of Beckmann (1952). In contrast to the prevalent Monge-Kantorovich approach to transportation, Beckmann’s paradigm of “continuous transportation” captures the trajectories along which a commodity is moving continuously from producers to consumers.

Let $\pi_p(x)$ and $\pi_c(x)$ be the density of production and consumption at a given geographical location x thought as a point in some Euclidean space and $\rho = \rho(x)$ be a weight function. Let $c = c(x)$ be a vector field representing the direction and the intensity of the commodity flow such that the local imbalances created in regions where $\pi_p \neq \pi_c$ are compensated in the steady state by the ρ -weighted flow. This compensation boils down to a condition that the divergence $\text{div}[\rho \cdot c]$ must be equal to $\pi_p - \pi_c$. Let $\Phi(c(x))$ be the local cost of transportation. For given π_p, π_c, ρ , and Φ , Beckmann’s problem is to find the flow c that compensates the supply-demand imbalances and has the minimal total cost:

$$\min_{c: \text{div}[\rho \cdot c] + \pi_c - \pi_p = 0} \int \Phi(c(x)) \rho(x) dx. \quad (1)$$

We demonstrate that a dual problem to revenue maximization takes the form of Beckmann’s problem (1) with a particular cost function Φ and marginals π_p and π_c satisfying a certain majorization constraint. Similar generalizations of Beckmann’s problems are related to mean-field limits of the Wardrop equilibria for congested optimal transportation games (Santambrogio, 2015; Carlier, 2012) but have not appeared in the context of auction design.

We establish the strong duality, namely, the optimal revenue of the auctioneer equals the optimal value of the dual problem. The strong duality is especially useful if combined with the existence of a solution, i.e., if the optima is attained. This combination enables complementary slackness conditions essential both for

constructing explicit solutions and for analysis of structural properties. We demonstrate both the existence of an optimal auction and an optimal vector field c , which is one of the most technically challenging parts of the paper. Specifically, the existence of c is not guaranteed unless we allow for singular vector fields represented by vector measures. One might think that singular vector fields is a technical peculiarity, but we demonstrate that singular solutions emerge even in toy examples.

We then propose a new numerical approximation scheme that allows us to conduct simulations that were previously out of reach. The approach relies on a combination of multi-to-single-agent reduction of [Cai et al. \(2012\)](#) and [Alaei et al. \(2019\)](#) avoiding the curse of dimensionality at the cost of dealing with a non-linear feasibility constraint, majorization theory insights ([Kleiner et al., 2021](#)) to handle this constraint, and cutting-edge numerical methods to handle optimization over convex functions that speed up the algorithm in practice. As far as we know, algorithms based on multi-to-single-agent reduction have never been previously implemented. We note that even the advanced neural-network approach of ([Dütting et al., 2019](#)) does not produce the outcome detailed and reliable enough to make structural conclusions because of the curse of dimensionality.

Using our approximation scheme, we compute optimal auctions for several bidders having uniformly distributed values over two items, bound the revenue loss from using sub-optimal designs, and show how the results change as the number of bidders grows.

Single bidder versus multiple bidders. To get more intuition about our approach and to highlight the specific features of the multi-bidder setting, we compare it to the single-bidder benchmark of the monopolist’s problem.

An important advance in understanding the monopolist’s problem was made by [Daskalakis et al. \(2017\)](#) who comprehensively showed how to reduce it to an optimal transportation one. Interestingly, instead of Beckmann’s problem arising in the multi-bidder setting, the dual derived by [Daskalakis et al. \(2017\)](#) is the Monge–Kantorovich optimal transportation problem with a majorization constraint. The Monge–Kantorovich problem has the following form:

$$\min_{\gamma: \gamma_1=\pi_p, \gamma_2=\pi_c} \int |x - y| d\gamma(x, y), \quad (2)$$

where $\pi_p(x) dx$ is the geographical distribution of production, $\pi_c(y) dy$ is the distribution of consumption, and the goal is to find a transportation plan γ such that the total transportation cost given by the integral in (2) is minimal and supply meets demand, i.e., the marginal of γ on the first coordinate is π_p and on the second, π_c . In contrast to Beckmann’s problem, the transportation happens momentarily: only initial and final destinations are captured by the plan γ , not the trajectories connecting them.

We conclude that the revenue-maximization problem for a single-bidder has two differently looking optimal-transportation duals: the Monge–Kantorovich dual and Beckmann’s one. This indicates the connection between the two duals themselves, in particular, their values must be equal. It turns out that for a single bidder, the cost function Φ in Beckmann’s problem can be simplified to $\Phi(x) = |x|$. For this cost function, the values of (1) and (2) are known to coincide by the so-called Beckmann’s duality (Santambrogio, 2015, Section 4.2). The presence of the two duals for the monopolist’s problem is a repercussion of this duality.

For a single bidder, one can use any of the two duals. However, the link between revenue maximization and the Monge–Kantorovich problem turns out to be limited to the single-bidder case. By contrast, the connection to Beckmann’s problem generalizes to any number of bidders and items.

Let us now highlight how our methodology overcomes the limitations of the approach by Daskalakis et al. (2017) in the multi-bidder case. The standard first step in the analysis of the monopolist’s problem is replacing the non-tractable maximization over mechanism via a handy maximization over interim utility functions $u = u(x)$, where x is the buyer’s type (Rochet and Choné, 1998). The Rochet-Choné representation is the starting point for the analysis of Daskalakis et al. (2017). The corresponding optimization problem is of the form

$$\max_u \int \left(\langle \nabla u(x), x \rangle - u(x) \right) \rho(x) \, dx,$$

where the maximum is taken over non-decreasing convex non-negative 1-Lipschitz functions u .

The Rochet-Choné representation can be generalized to multi-bidder problems at the cost of getting an extra constraint, the condition by Border (1991) capturing feasibility of the corresponding interim allocation rule. We rely on a less known equivalent condition by Hart and Reny (2015) resulting in a majorization constraint on the distribution of u ’s gradient. The connection to the majorization theory fuels our analysis and simulations. This theory has multiple recent applications in economic design; see, e.g., (Kleiner et al., 2021; Arieli et al., 2019; Candogan and Strack, 2021; Nikzad, 2022). However, majorization constraints on objective’s gradient have been studied neither in economic nor in mathematical literature.

The presence of the non-local constraint on the gradient does not allow us to get rid of the derivatives of u , which was crucial for the approach of Daskalakis et al. (2017). This obstacle explains why their approach does not generalize to the multi-bidder setting and why our dual problem does not look similar to the Monge-Kantorovich one. The non-local constraint on the gradient is a major complication; it leads to involved functional classes needed to establish strong duality and the existence of a solution to the dual.

Related literature

Linear programs and their duals are ubiquitous in microeconomics and economic design (Vohra, 2011; Bichler, 2017). While classic results focus on the finite-dimensional case corresponding to a finite number of types, the modern literature is increasingly interested in “continuous” infinite-dimensional settings highlighting structural properties of solutions usually hidden behind combinatorial artifacts of discrete type spaces.

Apart from multi-item auctions discussed below, infinite-dimensional linear programs and their duals naturally arise in various contexts, e.g., informationally or distributionally robust auction design (Bergemann et al., 2016; Kogiyigit et al., 2020; Suzdaltsev, 2020), information economics (Kolotilin, 2018; Dworzak and Martini, 2019; Dizdar and Kováč, 2020; Arieli et al., 2021a). Infinite-dimensional programs often have the structure similar to the Monge-Kantorovich optimal transportation, for example, in the context of sorting on the labor market (Boerma et al., 2021), matching with transferable utility and principal-agent problems (Chiappori et al., 2010), econometrics (Galichon, 2021), optimal taxation (Steinerberger and Tsyvinski, 2019), strategic learning and forecasting (Gensbittel, 2015; Arieli et al., 2021b; Guo and Shmaya, 2021). Other economic applications of optimal transport can be found in (Figalli et al., 2011; McCann and Zhang, 2019) and are surveyed by Galichon (2016) and Carlier (2012). For non-linear economic problems, a dual approach sharing some similarity with optimal transportation duality was proposed by Nöldeke and Samuelson (2018). A comprehensive presentation of the mathematical theory of transportation can be found in the books by Santambrogio (2015) and Villani (2009) and in surveys Bogachev and Kolesnikov (2012); Guillen and McCann (2013).

The continuous model of transportation developed by Beckmann (1952) is one of the classical economic models of transport networks that had considerable early popularity. It has not been used much in the recent economic literature with the exception of spatial equilibrium models of Fajgelbaum and Schaal (2020) and Allen and Arkolakis (2014).¹

For infinite-dimensional problems, guessing the form of a dual is usually straightforward and the central questions become whether the duality gap is zero or not (strong versus weak duality) and whether primal and dual solutions exist; both strong duality and the existence are needed for complementary slackness conditions to hold. In auction design, these questions have only been studied in the single-bidder case. Daskalakis et al. (2017) established the connection to optimal transport, demonstrated the strong duality, and showed the existence; their proofs

¹Beckmann’s problem has played an important role in developing the modern mathematical theory of optimal transportation anticipating the dynamic viewpoint which is discussed in Appendix F.

were then simplified by [Kleiner and Manelli \(2019\)](#). For several bidders, [Gianakopoulos and Koutsoupas \(2018\)](#) partially relaxed the incentive-compatibility constraint and got a weakly dual problem sharing some similarity with the maximal flow one. In contrast to our paper, they did not discuss the issue of existence (non-zero duality gap diminishes the importance of this question) and did not use the multi-to-single-agent reduction so that the dimension of the problem grows with the number of bidders. For a finite number of types, a strongly dual problem resembling the maximal flow one was derived by [Cai et al. \(2019\)](#). Since their result is finite-dimensional, the existence questions are mute and the strong duality becomes a consequence of the standard linear programming duality.

Even for a single bidder, optimal multi-item auctions can be complex and require non-linear pricing of a continuum of fractional bundles; explicit answers are known in a few particular cases such as uniformly or exponentially distributed values ([Daskalakis et al., 2017](#)). A primal approach of [Haghpanah and Hartline \(2021\)](#) based on virtual surplus maximization provides an alternative to optimal transportation technique of [Daskalakis et al. \(2017\)](#) and, in some cases, pins down an optimal mechanism, e.g., it shows when pure bundling is optimal in the single-bidder case; see also ([Hartline, 2013](#), Chapter 8). Instead of looking for optimal mechanisms the literature has mainly focused on either showing that a simple mechanism can guarantee a certain fraction of the optimal revenue or asking how well one can approximate the optimal mechanism withing a certain parametric class; see representative papers ([Hart and Reny, 2019](#); [Babaioff et al., 2020](#)) and ([Hart and Nisan, 2017](#); [Babaioff et al., 2021](#)). The only explicitly solved multi-item auction with several bidders assumes that bidders' valuations are binary ([Yao, 2017](#)).

2 Model

We work in the standard setting of Bayesian auction design with quasilinear bidders having i.i.d. additive utilities over items.

There is a set $\mathcal{B} = \{1, 2, \dots, B\}$ of $B \geq 1$ bidders and a set $\mathcal{I} = \{1, 2, \dots, I\}$ of $I \geq 1$ items. We assume that the items are divisible and normalize the total amount of each item to one unit. As usual, indivisible items can be made divisible by interpreting fractional amounts as probability shares.

Bidders treat the items as perfect substitutes and, hence, bidders' preferences are modelled by additive utility functions quasi-linear in money. The utility function of a bidder $b \in \mathcal{B}$ receiving a bundle $p_b \in \mathbb{R}_+^{\mathcal{I}}$ of items for a price t_b takes the form

$$\langle p_b, x_b \rangle - t_b,$$

where $\langle \cdot, \cdot \rangle$ is the standard dot product in $\mathbb{R}^{\mathcal{I}}$ and the vector $x_b \in \mathbb{R}_+^{\mathcal{I}}$ specifies b 's maximal willingness to pay for each of the items. The vector x_b can be seen

as bidder b 's type and constitutes the bidder's private information. Each bidder's type x_b belongs to the set of types $X = [0, 1]^{\mathcal{I}}$.

We assume that the fraction of bidders of different types in the population is described by a density ρ positive on X and zero beyond. The bidders are chosen from this population independently and, hence, the types $x_b \in X$, $b \in \mathcal{B}$, are i.i.d. draws with the distribution μ where $d\mu(x_b) = \rho(x_b) dx_b$. The auctioneer and bidders know ρ and each bidder observes the realization of her own type.

A mechanism which we also refer to as auction is given by a collection of bundles $P = (P_b(x))_{b \in \mathcal{B}}$ and transfers $T = (T_b)_{b \in \mathcal{B}}$ for each profile of types $x = (x_b)_{b \in \mathcal{B}}$. Formally, a mechanism (P, T) is a measurable map $X^{\mathcal{B}} \rightarrow \mathbb{R}_+^{\mathcal{I} \times \mathcal{B}} \times \mathbb{R}^{\mathcal{B}}$:

$$(x_b)_{b \in \mathcal{B}} \rightarrow \left(P_b((x_b)_{b \in \mathcal{B}}), T_b((x_b)_{b \in \mathcal{B}}) \right)_{b \in \mathcal{B}}.$$

Here P_b is the bundle received by a bidder $b \in \mathcal{B}$ and T_b is the amount of money she pays to the auctioneer. A mechanism is feasible if for any profile of types $(x_b)_{b \in \mathcal{B}}$

$$\sum_{b \in \mathcal{B}} P_{b,i}((x_b)_{b \in \mathcal{B}}) \leq 1 \quad \text{for all items } i \in \mathcal{I}, \quad (3)$$

i.e., the auctioneer has only one unit of each item to sell and so a mechanism cannot allocate more than one unit.

The auctioneer aims to design an auction maximizing the expected revenue $\sum_{b \in \mathcal{B}} T_b$. Bidders' types are their private information and a bidder may misreport her type if this brings her higher utility. Similarly, participation is voluntary and bidders may decide not to take part in the auction if they do not expect this to be profitable. Hence, providing incentives for truthful behavior and participation becomes design constraints. To formalize them, compute the expected allocation and transfer faced by a bidder b of a given type x_b assuming that others report their types truthfully:

$$\bar{P}_b(x_b) = \int_{X^{\mathcal{B} \setminus \{b\}}} P_b((x_b)_{b \in \mathcal{B}}) \cdot \left(\prod_{d \in \mathcal{B} \setminus \{b\}} \rho(x_d) \right) dx_1 \cdots dx_{b-1} dx_{b+1} \cdots dx_B, \quad (4)$$

$$\bar{T}_b(x_b) = \int_{X^{\mathcal{B} \setminus \{b\}}} T_b((x_b)_{b \in \mathcal{B}}) \cdot \left(\prod_{d \in \mathcal{B} \setminus \{b\}} \rho(x_d) \right) dx_1 \cdots dx_{b-1} dx_{b+1} \cdots dx_B. \quad (5)$$

Such one-bidder marginals (\bar{P}_b, \bar{T}_b) of the original mechanism (P, T) are known as its reduced forms or interim mechanisms. The reduced mechanism for a bidder b captures how her expected utility depends on her type and her report, i.e., all the information relevant to her: if her type is x_b and she reports to be of type x'_b , while other bidders remain truthful, b 's expected utility takes the form

$$\langle \bar{P}_b(x'_b), x_b \rangle - \bar{T}_b(x'_b).$$

A mechanism is called Bayesian incentive-compatible if truth-telling is a Bayesian equilibrium, i.e., no bidder b has an incentive to misreport her values if others report truthfully. Formally,

$$\langle \bar{P}_b(x_b), x_b \rangle - \bar{T}_b(x_b) \geq \langle \bar{P}_b(x'_b), x_b \rangle - \bar{T}_b(x'_b) \quad (6)$$

for all $x_b, x'_b \in X$ and $b \in \mathcal{B}$.

A mechanism is called individually rational if no bidder wants to abstain from participation, i.e., nobody gets a negative expected utility. Formally,

$$\langle \bar{P}_b(x_b), x_b \rangle - \bar{T}_b(x_b) \geq 0 \quad (7)$$

for all $x_b \in X$ and $b \in \mathcal{B}$.

The auctioneer's design problem takes the following form.

Auctioneer's problem: *maximize the expected revenue*

$$\int_{X^{\mathcal{B}}} \left(\sum_{b \in \mathcal{B}} T_b((x_b)_{b \in \mathcal{B}}) \right) \cdot \left(\prod_{b \in \mathcal{B}} \rho(x_b) \right) dx_1 \cdots dx_B \quad (8)$$

over *individually-rational Bayesian incentive-compatible feasible mechanisms* (P, T) .

In the case of a single bidder ($B = 1$), the auctioneer's problem becomes the multi-item monopolist's problem. Note that for $B = 1$, the reduced mechanism coincides with the original one, i.e., $\bar{P}_1 \equiv P_1$ and $\bar{T}_1 \equiv T_1$. In what follows, we will use the monopolist's problem as a benchmark and, in particular, connect our characterization to the one obtained by [Daskalakis et al. \(2017\)](#).

3 Multi-bidder version of Rochet-Choné representation

A common starting point for the analysis of the monopolist's problem is its equivalent representation derived in [Rochet and Choné \(1998\)](#). We first recall their insight in the single-bidder setting and then describe its extension to the general case of $B \geq 1$ bidders.

3.1 1-bidder case (monopolist's problem)

With each mechanism (P, T) , we can associate the interim utility function $u(x) = \langle P(x), x \rangle - T(x)$, i.e., the expected utility obtained by a bidder of type x . Following [Rochet and Choné \(1998\)](#), the monopolist's problem can be rewritten as

a maximization over the utility function u under some constraints. Bayesian incentive compatibility and individual rationality boil down to u being a convex non-negative function. The allocation probabilities $P(x)$ are given by the gradient $\nabla u(x)$. Hence, $\langle \nabla u(x), x \rangle$ is the utility that the bidder derives from the allocated items. As the total utility is $u(x)$, the difference $\langle \nabla u(x), x \rangle - u(x)$ is the payment that goes to the monopolist. Consequently, the monopolist's problem reduces to maximizing

$$\int_X \left(\langle \nabla u(x), x \rangle - u(x) \right) \rho(x) \, dx, \quad (9)$$

over convex $u : X \rightarrow \mathbb{R}_+$ such that $\nabla u(x) \in [0, 1]^I$. The last condition originates from the requirement of feasibility: for each item i , the allocated amount

$$P_i(x) = \frac{\partial u}{\partial x_i}(x) \quad (10)$$

has to be between 0 and 1.

3.2 The case of $B \geq 1$ bidders (auctioneer's problem)

Consider now the auction-design problem with $B \geq 1$ bidders. We show that this problem can be reduced to an optimization problem that is similar to the monopolist's problem but the feasibility constraint $\frac{\partial u}{\partial x_i}(x) \leq 1$ on the gradient's values is replaced by a non-local majorization condition on the distribution of the gradient.

Definition 1 (Majorization²). For a pair of measures ν and ν' , we say that ν majorizes ν' if $\int \varphi \, d\nu \geq \int \varphi \, d\nu'$ for any convex non-decreasing function φ . A random variable ξ majorizes ξ' if the distribution of ξ majorizes that of ξ' . We write $\nu \succeq \nu'$ and $\xi \succeq \xi'$.

Informally, majorization means that ν can be obtained from ν' by combining mean-preserving spreads with moving mass to higher values.

As we will see, the auctioneer's problem with B bidders is equivalent to the following one.

Multi-bidder Rochet-Choné problem: *maximize*

$$B \cdot \int_X \left(\langle \nabla u(x), x \rangle - u(x) \right) \rho(x) \, dx \quad (11)$$

²Majorization is not to be confused with a closely related notion of dominance with respect to the convex order also known as the Blackwell order and second-order stochastic dominance. The latter corresponds to taking any convex φ , not necessarily non-decreasing. For probability measures, convex dominance implies that ν and ν' have the same mean, while for majorization, the majorizing measure can have a higher mean, i.e., $\int t \, d\nu(t) \geq \int t \, d\nu'(t)$.

over convex non-decreasing functions $u : X \rightarrow \mathbb{R}_+$ with $u(0) = 0$ and such that for all $i \in \mathcal{I}$

$$\frac{\partial u}{\partial x_i}(\chi) \preceq \xi^{B-1}, \quad (12)$$

where $\chi \in X$ is distributed with the density ρ and ξ is uniformly distributed on $[0, 1]$.

Let us clarify the meaning of the condition (12). Each component of the gradient $\frac{\partial u}{\partial x_i}(\chi)$ is treated there as a random variable by assuming that the argument $\chi \in X$ is random and distributed with the density ρ and the distribution of this random variable must be majorized by the distribution of ξ^{B-1} , where ξ is uniform on $[0, 1]$. An equivalent way to write this condition is to assume that for any non-decreasing convex φ

$$\int_X \varphi \left(\frac{\partial u}{\partial x_i}(x) \right) \rho(x) dx \leq \int_0^1 \varphi(z^{B-1}) dz. \quad (13)$$

Note that we do not assume that the function u is smooth and, hence, the partial derivative $\frac{\partial u}{\partial x_i}(x)$ may not exist for some x . Despite this fact, the optimization problem (11) is well-defined since the gradient of a convex function exists almost everywhere and integration with respect to an absolutely continuous measure is not sensitive to the behavior of the integrand on sets of zero Lebesgue measure; see Appendix A for basics of convex analysis.

Proposition 1. *The optimal revenue in the auctioneer’s problem (8) and the value of the multi-bidder Rochet-Choné problem (11) coincide and the optima in both problems are attained.*

A proof of this proposition is contained in Appendix B. Below we discuss the underlying ideas and why representing the auctioneer’s problem by (11) is useful.

Proposition 1 allows one to treat auctions with a different number of bidders in a similar way.³ However, the single-bidder case, $B = 1$, is special. By plugging in $\varphi(z) = \max\{0, z - 1\}$ to (13), we see that majorization implies $\frac{\partial u}{\partial x_i}(x) \leq 1$ for any number of bidders. For one bidder, however, the reverse implication also holds as the right-hand side of (13) is equal to $\varphi(1)$ and φ is monotone. Consequently, the dominance condition on the gradient’s distribution boils down to the pointwise condition on the gradient’s values and we obtain the classic Rochet-Choné representation (9) used by Daskalakis et al. (2017).

For $B > 1$, the majorization constraint becomes non-local and non-trivially restricts the distribution of the gradient rather than its pointwise values. As we

³Treating the number of bidders B in (11) as a continuous parameter, one can even interpolate between auctions with different numbers of bidders.

will see in Section 4, this non-locality is a dramatic complication compared to the single-bidder case. Another complication is that the existence of optimal multi-item multi-bidder auctions has not been known and our results required rather sophisticated functional analytic arguments which are developed in the appendix.

The representation (11) makes apparent the connection of the auctioneer’s problem to the theory of optimal transportation and to the majorization theory, which fuel our analysis in Section 4. The majorization theory offers a universal toolbox to study optimization problems under majorization constraints and has recently been recognized as a powerful tool applicable to various problems of economic design (Kleiner et al., 2021). Modern optimal transportation literature also deals with non-classical settings involving majorization constraints; see, e.g., (Henry-Labordère, 2017; Gozlan et al., 2017). In Section 5, we demonstrate that Proposition 1, combined with optimal-transportation insights, leads to an algorithm allowing us to compute optimal auctions with unprecedented precision.

The majorization constraint originates as a repercussion of the original feasibility constraint when the auctioneer’s problem is rewritten as maximization over reduced forms. The idea of rewriting the auctioneer’s problem this way has recently gained popularity in algorithmic mechanism design as it drastically reduces the dimension; see Alaei et al. (2019) and the discussion in Section 5 below. However, the resulting problem has not been written in terms of the utility function u making the similarity with the classic Rochet-Choné representation apparent. More importantly, the literature has previously relied on the feasibility constraint on reduced forms from the original paper of Border (1991), while we use an equivalent condition formulated in terms of majorization and found by Hart and Reny (2015); see also (Kleiner et al., 2021; Gershkov et al., 2021).

We now discuss the ideas underlying Proposition 1 in more detail.

Maximization over reduced-form mechanisms. The first idea is to replace the maximization over mechanisms (P, T) in the auctioneer’s problem by the maximization over the corresponding reduced forms $(\bar{P}_b, \bar{T}_b)_{b \in \mathcal{B}}$. The constraints of Bayesian incentive-compatibility and individual rationality are originally formulated in these terms and the revenue objective rewrites as the sum of revenues collected by each of the one-bidder mechanisms (\bar{P}_b, \bar{T}_b) . An easy symmetrization argument shows that we can assume that all the one-bidder mechanisms are the same for all bidders: $(\bar{P}_b, \bar{T}_b) = (\bar{P}, \bar{T})$. Thus the auctioneer’s problem reduces to maximization of B times the revenue of a Bayesian incentive-compatible individually-rational one-bidder mechanism (\bar{P}, \bar{T}) , i.e., to the monopolist’s problem to which the standard Rochet-Choné representation (9) can be applied. However, we get an extra constraint on (\bar{P}, \bar{T}) originating from the feasibility constraint on a B -bidder mechanism (P, T) . It takes a form of the majorization condition (12) which we discuss in the next item.

Border's constraint via majorization. A one-bidder mechanism (\bar{P}, \bar{T}) is a feasible reduced form if there is a feasible B -bidder mechanism (P, T) such that its reduced form (\bar{P}_b, \bar{T}_b) coincides with (\bar{P}, \bar{T}) for any bidder $b \in \mathcal{B}$.

The first characterization of feasible reduced forms was proved by [Border \(1991\)](#). [Hart and Reny \(2015\)](#) showed that Border's condition can be restated in terms of majorization. Namely, they proved that a single-bidder mechanism (\bar{P}, \bar{T}) is feasible reduced form if and only if for all items $i \in \mathcal{I}$

$$\bar{P}_i(\chi) \preceq \xi^{B-1}, \quad (14)$$

where χ is distributed with the density ρ and ξ is uniformly distributed on $[0, 1]$. The upper bound in (14) is rather intuitive as it corresponds to the reduced form of a mechanism (P, T) allocating each item i to the bidder $b \in \mathcal{B}$ with the highest x_i . In other words, any reduced form is majorized by the reduced form of the efficient allocation rule. [Hart and Reny \(2015\)](#) proved the result for $I = 1$ item. For $I \geq 1$, the same dominance condition has to be applied to each of the components of $\bar{P} = (\bar{P}_i)_{i \in \mathcal{I}}$ since the original feasibility constraint for (P, T) restricts the allocation of each item separately; see ([Cai et al., 2012](#)).

As in the classic Rochet-Choné formula, the allocation probabilities can be recovered as the gradient of the bidder's interim utility (10). Hence, the Border's constraint (14) rewrites as the condition (12) on function u in the multi-bidder Rochet-Choné representation.

The existence of optimal solutions. The established equivalence between the auctioneer's problem and the multi-bidder Rochet-Choné representation allows us to construct a solution to one based on a solution to the other. Hence, it is enough to show that the optimum is attained for the problem (11). This follows from a compactness argument. The set of feasible u is compact and the objective is continuous in the topology of the space of continuous functions. Hence, the optimal u exists since a continuous functional attains its maximal value on a compact set.

A subtle point in this argument is in choosing the right topology. One might think that the topology of continuous functions is too weak to control the gradient and preserve the condition (12) on the gradient's distribution in the limit. Indeed, differentiability is too fine to be preserved by the continuous-function topology. However, thanks to the fact that feasible u are convex, the local property of differentiability can be replaced with a lower bound by an appropriate affine function (see the definition of subdifferential in [Appendix A](#)) which is respected by continuous-function limits.

4 Duality

In Section 3, we saw that the auctioneer’s problem can be reduced to the multi-bidder Rochet-Choné problem (Proposition 1), which is an infinite-dimensional convex program. In this section, we show that for any number of bidders, the dual to this program is a version of Beckmann’s transportation problem (Beckmann, 1952).

Although the Rochet-Choné problem with $B \geq 2$ and $B = 1$ bidders look similar, the dual problem that we find is entirely different from the one identified by Daskalakis et al. (2017) in the case of $B = 1$ bidder. The difference originates from the feasibility constraint in the Rochet-Choné problem. For one bidder, it is given by the pointwise upper bound on the gradient $\frac{\partial u}{\partial x_i}(x) \leq 1$ and is thus equivalent to the 1-Lipschitz property of u . This allows one to get rid of u ’s derivatives completely, which is the crucial step in the approach of Daskalakis et al. (2017). This simplification has no analogue for $B \geq 2$ bidders, where the pointwise bound is replaced by the non-local majorization constraint (12). This non-locality becomes a major obstacle for establishing strong duality and, especially, for proving the existence of a solution to the dual problem. Our approach overcomes this obstacle. In the case of $B = 1$ bidder, both approaches are applicable and we demonstrate that the dual of Daskalakis et al. (2017) can be deduced from our dual. We note that demonstrating strong duality with non-local constraints is new not just for the economic literature but also to the broader mathematical context.

We first discuss the essence of Beckmann’s problem and then define the problem formally paying attention to nuances needed to establish the connection to the auctioneer’s problem. We are given a cost function Φ , densities of production $\pi_p(x)$ and consumption $\pi_c(x)$ of a certain commodity at every geographical location $x \in X$, where X is a subset of an Euclidean space, and a weight-function ρ on X . The goal is to find a transportation flow having the minimal cost and compensating supply-demand imbalance $\pi = \pi_c - \pi_p$ in the steady state, i.e., such that the total ρ -weighted inflow in each region is equal to the difference between supply and demand in this region. The direction and intensity of the flow are represented by a vector field $c = c(x)$ and compensation of imbalances boils down to the following identity that must be satisfied by the divergence:⁴ $\text{div}[\rho \cdot c](x) + \pi(x) = 0$. The problem is to minimize the total cost $\int_X \Phi(c)\rho(x) dx$ over all such vector-fields.

In the application to the auctioneer’s problem, the set of geographical locations

⁴The intuition is as follows. The imbalances are compensated if for any region A (it is enough to consider infinitesimally-small cubes), the outflow through the boundary is equal to the total imbalance in A . By the Gauss theorem, the outflow equals $\int_A \text{div}[\rho \cdot c] dx$. We end up with the condition $\int_A \text{div}[\rho \cdot c] dx = \int_A (\pi_p - \pi_c) dx$ which holds for any A and thus the integrands must be equal.

X will coincide with the set of types $X = [0, 1]^I$ and the weight ρ will be the density of types' distribution. The supply-demand imbalance π will be given by a signed measure which may have singularities. Accordingly, we need to allow the divergence to become a measure as well. To explain the intuition behind the formal definition, for a moment assume that ρ is smooth and equals zero on the boundary of X . Then, using the Gauss theorem or just integrating by parts, we obtain that

$$\int_X \langle \nabla u(x), c(x) \rangle \rho(x) dx = - \int_X u(x) \cdot \operatorname{div}[\rho \cdot c] dx \quad (15)$$

for any smooth function u (there is no term corresponding to the contribution of the boundary of X as we assumed that ρ vanishes there). This formula suggests the formal definition. For a vector field c and weight ρ , the ρ -divergence $\operatorname{div}_\rho[c]$ is a measure on X such that the integration-by-parts relation

$$\int_X \langle \nabla u(x), c(x) \rangle \rho(x) dx = - \int_X u(x) d(\operatorname{div}_\rho[c])(x) \quad (16)$$

holds for any smooth u . In general, the contribution of the boundary cannot be neglected and so $\operatorname{div}_\rho[c]$ may have boundary singularities even for smooth c and⁵ ρ .

Beckmann's problem. *The set of geographical locations is $X = [0, 1]^I$. Spacial imbalance of production and consumption is given by a signed measure π on X ; the imbalance is assumed to have local nature, i.e., $\pi(X) = 0$. Given a convex cost function $\Phi: \mathbb{R}^I \rightarrow \mathbb{R} \cup \{+\infty\}$ and a density $\rho: X \rightarrow \mathbb{R}_+$, the goal is to minimize the cost $\int_X \Phi(c(x)) \cdot \rho(x) dx$ over continuously differentiable vector fields $c: X \rightarrow \mathbb{R}^I$ such that $\operatorname{div}_\rho[c] + \pi = 0$. The value of Beckmann's problem is denoted by*

$$\operatorname{Beck}_\rho(\pi, \Phi) = \inf_{c: \operatorname{div}_\rho[c] + \pi = 0} \int_X \Phi(c(x)) \cdot \rho(x) dx. \quad (17)$$

If there are no smooth c such that $\operatorname{div}_\rho[c] + \pi = 0$, i.e., the minimization is over an empty set, we assume that $\operatorname{Beck}_\rho(\pi, \Phi) = +\infty$.

We now connect Beckmann's problem to auctions. For this purpose, we make the imbalance π a free parameter satisfying a majorization constraint. To describe this constraint, consider the revenue objective in the Rochet-Choné problem (11) and get rid of derivatives via integration by parts

$$\int_X (\langle \nabla u(x), x \rangle - u(x)) \rho(x) dx = -u(0) + \int_X u(x) dm(x), \quad (18)$$

⁵A similar use of measure-valued derivatives can be found in (Ambrosio et al., 2000).

where m is a signed measure such that this identity holds for any smooth u . We consider the following majorization constraint on⁶ π :

$$\pi \succeq m. \quad (19)$$

A similar constraint appears in the single-bidder result by [Daskalakis et al. \(2017\)](#) who refer to m as the transform measure.

To define the cost function Φ , consider a collection $(\varphi_i)_{i \in \mathcal{I}}$ of non-decreasing convex functions on \mathbb{R}_+ with $\varphi_i(0) = 0$. Let φ_i^* be the Legendre transform of φ_i , i.e., $\varphi_i^*(y) = \sup_x \langle x, y \rangle - \varphi_i(x)$; see [Appendix A](#). The cost function Φ is separable and takes the following form

$$\Phi(c) = \sum_{i \in \mathcal{I}} \varphi_i^*(|c_i|). \quad (20)$$

We note that the higher is φ_i , the lower is φ_i^* and so is the cost Φ .

Theorem 1. *In the auctioneer's problem (8) with $|\mathcal{B}| = B \geq 1$ bidders, $|\mathcal{I}| = I \geq 1$ items, and bidders' types distributed on $X = [0, 1]^{\mathcal{I}}$ with positive density ρ , the optimal revenue coincides with*

$$B \cdot \inf_{\substack{(\varphi_i)_{i \in \mathcal{I}}, \\ \pi \succeq m}} \left[\text{Beck}_\rho(\pi, \Phi) + \sum_{i \in \mathcal{I}} \int_0^1 \varphi_i(z^{B-1}) \, dz \right], \quad (21)$$

where Φ is given by (20) and $\varphi_i: \mathbb{R}_+ \rightarrow \mathbb{R}_+ \cup \{+\infty\}$ are non-decreasing convex functions with $\varphi_i(0) = 0$ for each item $i \in \mathcal{I}$.

Theorem 1 is a particular case of a more general duality result ([Theorem 5](#)) proved in [Appendix C](#). The proof goes in two steps. First, we prove a partial duality result ([Theorem 3](#)) internalizing the majorization constraint. It can be interpreted as the equivalence between the auctioneer's problem and the monopolist's problem with adversarial production costs. We derive a novel a priori bound on the solutions of the latter problem ([Proposition 2](#)) with a clear economic interpretation: the monopolist can guarantee a non-negative revenue not only ex-ante but ex-post at no cost. Then, with this a priori bound, we deduce the complete duality. A byproduct of the proof is that one can assume that the vector field c in Beckmann's problem from (21) has non-negative components.

Let us see why the minimization problem (21) is well-defined, i.e., why we minimize over a non-empty set. We need to demonstrate that there is always

⁶The definition of majorization ([Definition 1](#)) is applicable to multidimensional signed measures. In particular, (19) means that $\int u \, d\pi \geq \int u \, dm$ for any convex non-decreasing u on X .

$\pi \succeq m$ such that $\operatorname{div}_\rho[c] + \pi = 0$ for some smooth vector field c , and so Beckmann's problem has a finite value. It turns out that we can always take $\pi = -\operatorname{div}_\rho[x]$ and $c(x) = x$. Let us demonstrate that the majorization condition $\int_X u(x) d\pi(x) \geq \int_X u(x) dm(x)$ holds. We rewrite both sides by the definitions of the divergence and m and get

$$\int_X \langle x, \nabla u(x) \rangle \rho(x) dx \geq u(0) + \int_X (\langle x, \nabla u(x) \rangle - u(x)) \rho(x) dx.$$

The dot-product terms cancel out and we end up with an equivalent inequality $\int_X u(x) \rho(x) dx \geq u(0)$ that holds for any non-decreasing u . We conclude that the problem (21) has a finite value. Moreover, we obtain that the auctioneer's optimal revenue is upper-bounded by

$$B \cdot \inf_{(\varphi_i)_{i \in \mathcal{I}}} \sum_{i \in \mathcal{I}} \left(\int_X \varphi_i^*(x_i) \rho(x) dx + \int_0^1 \varphi_i(z^{B-1}) dz \right). \quad (22)$$

In this upper bound, the minimization splits into a family of I identical one-dimensional minimization problems, one for each item $i \in I$. They can be solved explicitly and the resulting bound corresponds to full surplus extraction; see Appendix D.1.

4.1 Weak duality and complementary slackness

Strong duality results such as Theorem 1 can be seen as a combination of two statements: that the value of the primal problem is at most the value of the dual (weak duality) and that the gap between the values is zero. While the weak duality is always an easy part of the proof, this part is insightful as it explains the form of the dual and leads to complementary slackness conditions.

Let us see why the weak duality holds, i.e., why the optimal revenue is upper-bounded by (1). We know that the optimal revenue equals to $B \cdot \int_X (\langle x, \nabla u(x) \rangle - u(x)) \rho(x) dx$ for some convex non-decreasing function u with $u(0) = 0$ and such that Border's constraint (13) is satisfied (Proposition 1). Hence, the optimal revenue does not exceed

$$B \cdot \left[\int_X (\langle x, \nabla u(x) \rangle - u(x)) \rho(x) dx + \sum_{i \in \mathcal{I}} \left(\int_0^1 \varphi_i(z^{B-1}) dz - \int_X \varphi_i \left(\frac{\partial u}{\partial x_i}(x) \right) \rho(x) dx \right) \right] \quad (23)$$

for any non-decreasing convex functions φ_i on \mathbb{R}_+ with $\varphi_i(0) = 0$ (each term in the sum is non-negative by Border's constraint). The first integral can be rewritten as follows

$$\begin{aligned} \int_X (\langle x, \nabla u(x) \rangle - u(x)) \rho(x) dx &= \int_X u(x) dm(x) \\ &\leq \int_X u(x) d\pi(x) = \int_X \langle \nabla u(x), c(x) \rangle \rho(x) dx, \end{aligned} \quad (24)$$

where m is the transform measure from (18), π is an arbitrary measure such that $\pi \succeq m$ and c is any vector field such that $\text{div}_\rho[c] + \pi = 0$. The first equality holds by the definition of the transform measure, the inequality holds thanks to convexity of u , and the last equality is by the definition of divergence (16). The Fenchel inequality (inequality (36) in Appendix A) applied to $\psi_i(t) = \varphi_i(|t|)$ implies the following bound on the last integrand

$$\langle \nabla u(x), c(x) \rangle \leq \sum_{i \in \mathcal{I}} \varphi_i^*(|c_i(x)|) + \sum_{i \in \mathcal{I}} \varphi_i \left(\frac{\partial u}{\partial x_i}(x) \right), \quad (25)$$

where we used that $\psi_i^*(t) = \varphi_i^*(|t|)$ and non-negativity of u 's partial derivatives. Replacing the first summand in (23) by the resulting upper bound, we see that the terms with partial derivatives of u cancel out and the revenue is bounded from above by

$$B \cdot \left[\int_X \left(\sum_{i \in \mathcal{I}} \varphi_i^*(|c_i(x)|) \right) \rho(x) dx + \sum_{i \in \mathcal{I}} \int_0^1 \varphi_i(z^{B-1}) dz \right]$$

for all convex φ_i with $\varphi_i(0) = 0$, all measures $\pi \succeq m$, and smooth vector fields c such that $\text{div}_\rho[c] + \pi = 0$. Taking infimum over all such φ_i , π , and c , we conclude that the optimal revenue cannot exceed the right-hand side of (21) thus establishing the weak duality.

Complementary slackness conditions are a byproduct of the above computation. Let u^{opt} , φ_i^{opt} , π^{opt} , and c^{opt} be the optima in the primal Rochet-Choné problem (11), the dual problem (21), and internal Beckmann's problem, respectively. We know that u^{opt} exists by Proposition 1 and the existence of the rest of the optima is discussed below. For now, we assume that all of them exist. Under this assumption, the only way the value of the primal problem can be equal to the value of the dual (21) is if each inequality in the derivation of the weak duality holds as equality at u^{opt} , φ_i^{opt} , π^{opt} , and c^{opt} . Namely, each term in the sum from (23) must be zero, and the inequality in (24) together with the Fenchel inequalities used to derive (25) must all be equalities. These observations, combined with the complementary slackness condition for the Fenchel inequality (see Appendix A), lead to the following corollary.

Corollary 1 (Complementary slackness). *Optimal u^{opt} , functions φ_i^{opt} , measure π^{opt} , and vector field c^{opt} satisfy the following family of conditions:*

$$\int_X \varphi_i^{\text{opt}} \left(\frac{\partial u^{\text{opt}}}{\partial x_i}(x) \right) \rho(x) \, dx = \int_0^1 \varphi_i^{\text{opt}}(z^{B-1}) \, dz \quad (26)$$

$$\int_X u^{\text{opt}}(x) \, dm(x) = \int_X u^{\text{opt}}(x) \, d\pi^{\text{opt}}(x) \quad (27)$$

$$c_i^{\text{opt}}(x) \in \partial \varphi_i^{\text{opt}} \left(\frac{\partial u^{\text{opt}}}{\partial x_i}(x) \right) \quad (28)$$

In the last condition, ∂ denotes the subdifferential (35) and the inclusion holds for ρ -almost all $x \in X$.

4.2 Existence

Whether the optima exist or not may seem a technical peculiarity. The importance of this question is justified by the complementary slackness conditions (Corollary 1) which can be written down only if both primal and dual problems attain their optima.

We know that the optimal value of the Rochet-Choné problem (11) is attained at some u^{opt} . It turns out that the family of optimal functions φ_i^{opt} in the dual problem (21) also always exists and corresponds to an optimal strategy of an adversary in the auxiliary monopolist's problem with adversarial production costs discussed in Appendix C.1.

We note that Beckmann's problem is prone to absence of an optimal smooth vector field c^{opt} even for standard cost functions such as $\Phi(c) = \|c\|$. A workaround is to allow for generalized vector fields by replacing a smooth vector field c by a vector measure ς . Then the optimal vector measure ς is known to exist provided that the supply-demand imbalance π is absolutely continuous and, moreover, ς itself turns out to be absolutely continuous (Santambrogio, 2015, Theorem 4.16).

An additional complication of our setting is that the transform measure m typically has singularities on the boundary of X that are inherited by $\pi \succeq m$. As a result, to guarantee existence, we are forced to allow for vector-measures containing singular components.

Working with vector measures creates a new complication. The divergence of a vector measure may not be a measure anymore and can only be defined in the space of generalized functions (Ambrosio et al., 2000). As $\pi = -\text{div}_\rho[c]$, following this approach we would need to allow π to become a generalized function. This complication can be avoided by reformulating the constraint on the vector field bypassing π .

Consider the set \mathcal{C}^{mes} of non-negative vector measures $\varsigma = (\varsigma_i)_{i \in \mathcal{I}}$ satisfying the following condition

$$\int_X (\langle \nabla u(x), x \rangle - u(x)) \cdot \rho(x) dx \leq \sum_{i \in \mathcal{I}} \int_X \frac{\partial u}{\partial x_i}(x) d\varsigma_i(x) \quad (29)$$

for any smooth non-decreasing convex $u : X \rightarrow \mathbb{R}_+$ with $u(0) = 0$. By the Lebesgue decomposition theorem, each ς_i can be represented as the sum of the component that is absolutely continuous with respect to $\rho(x) dx$ and the singular one. We get

$$d\varsigma_i = c_i(x) \cdot \rho(x) dx + d\varsigma_i^{\text{sing}}(x). \quad (30)$$

If the singular component is absent and $c = (c_i)_{i \in \mathcal{I}}$ is smooth, we can define $\pi = -\text{div}_\rho[c]$ and see that the condition (29) is equivalent to the familiar majorization condition (19) on π .

The following extension of Theorem 1 guarantees that the optimum in the dual is attained. It is proved in Appendix C.

Theorem 2 (Extended dual). *The optimal revenue in the auctioneer's problem (8) coincides with*

$$B \cdot \min_{\substack{(\varphi_i)_{i \in \mathcal{I}}, \\ \varsigma \in \mathcal{C}^{\text{mes}}}} \sum_{i \in \mathcal{I}} \left(\varsigma_i^{\text{sing}}(X) + \int_X \varphi_i^*(c_i(x)) \rho(x) dx + \int_0^1 \varphi_i(z^{B-1}) dz \right) \quad (31)$$

and the minimum is attained. Here B is the number of bidders, c_i and $\varsigma_i^{\text{sing}}$ are given by (30), and φ_i are non-decreasing convex functions with $\varphi_i(0) = 0$ for each item $i \in \mathcal{I}$.

Note that the objectives in Theorems 1 and 2 match one another except for the fact that some mass in the extended dual can be transferred from the vector field c to the singular component of the vector measure. This additional flexibility turns out to be crucial for the existence of the optimum.

One may think that the appearance of singular measures is an artifact of a particular proof technique and that singularities do not appear at least in nice examples. This intuition turns out to be wrong and singular measures happen to reflect the essence of the problem. In Appendix D, we solve the dual problem explicitly for two uniform items and $B = 1$ bidder and see that, even in this simplest case, there are singularities on the boundary of the set of types X .

Theorem 2 allows us to write down the complementary slackness conditions without making an extra assumption that the optima exist.

Corollary 2 (Extended complementary slackness). *Consider optimal u^{opt} , $(\varphi_i^{\text{opt}})_{i \in \mathcal{I}}$, and ς^{opt} and decompose ς^{opt} into absolutely-continuous and singular components as in (30). Then all the previously discussed complementary slackness conditions (27), (28), and (26) hold. Moreover, there is one more condition:*

$$\frac{\partial u^{\text{opt}}}{\partial x_i}(x) = 1 \quad (32)$$

for $\varsigma_i^{\text{opt}, \text{sing}}$ -almost all x . In particular, u^{opt} has a partial derivative with respect to x_i for $\varsigma_i^{\text{opt}, \text{sing}}$ -almost all x .

4.3 Relation to Daskalakis et al. (2017) for $B = 1$ bidder

In Theorem 1, we saw that the dual to the auctioneer's problem is given by Beckmann's transportation problem for any number of bidders $B \geq 1$. For $B = 1$ bidder, Daskalakis et al. (2017) derived another dual taking a form of the Monge-Kantorovich optimal transportation problem (2). It is not surprising that the duals for $B \geq 2$ bidders and $B = 1$ bidder do not share any similarity as the feasibility constraint for several bidders becomes non-local and so the approach of Daskalakis et al. (2017) is not applicable. Here we focus on the case of $B = 1$ bidder, where both approaches can be used and so the lack of similarity between the two duals may seem surprising.

It turns out that the dual from Theorem 1 can be simplified in the single-bidder case. Indeed, $z^{B-1} \equiv 1$ for $B = 1$ bidder and so the second integral reduces to $\int_0^1 \varphi_i(z^{B-1}) dz = \varphi_i(1)$. We obtain that the value of the auctioneer's problem is equal to

$$\inf_{\substack{(\varphi_i)_{i \in \mathcal{I}}, \\ \pi \succeq m}} \left[\text{Beck}_\rho(\pi, \Phi) + \sum_{i \in \mathcal{I}} \varphi_i(1) \right]$$

with $\Phi(c) = \sum_i \varphi_i^*(|c_i|)$. This expression can be further simplified. The lower the cost function Φ in Beckmann's problem is, the lower is its value. By increasing φ_i pointwise, we decrease its conjugate φ_i^* . Hence, keeping $\varphi_i(1)$ fixed, the best choice given the requirements of convexity and $\varphi_i(0) = 0$ is the linear function: $\varphi_i(t) = \varphi_i(1) \cdot t$ on $[0, 1]$ and $\varphi_i(t) = +\infty$ for $t > 1$. Optimization over $\varphi_i(1)$ gives $\varphi_i(1) = 0$ and thus the conjugate $\varphi_i^*(t) = t$ for all t .

We obtain Beckmann's problem with the cost function given by l^1 -norm $\Phi(c) = \|c\|_1 = \sum_{i \in \mathcal{I}} |c_i|$. Importantly, this cost function is 1-homogeneous. Beckmann's problem with a 1-homogeneous Φ has a peculiar property: its value does not depend on the density ρ provided that it is smooth and positive, i.e., $\text{Beck}_\rho(\pi, \Phi) = \text{Beck}_1(\pi, \Phi)$, where in the second problem the density is equal

to 1. This property holds, since for any feasible vector field c in the second problem, $c' = \rho \cdot c$ is a feasible vector field in the first problem with the same value.

Corollary 3. *For $B = 1$ bidder whose type is distributed according to a smooth positive density ρ , the optimal revenue of the auctioneer (8) is equal to*

$$\inf_{\pi \succeq m} \text{Beck}_1(\pi, \|\cdot\|_1), \quad (33)$$

where the cost function is given by l^1 -norm $\|c\|_1 = \sum_{i \in \mathcal{I}} |c_i|$.

Beckmann’s problem with the Lebesgue reference measure and the cost function $\|\cdot\|_1$ is an exception where the Beckmann’s problem is known to be connected to the Monge-Kantorovich one.⁷ The so-called Beckmann’s duality states that, for any π ,

$$\text{Beck}_1(\pi, \|\cdot\|_1) = \min_{\gamma: \gamma_1 = \pi_c, \gamma_2 = \pi_p} \int \|x - y\|_1 d\gamma(x, y),$$

where π_c and π_p are the positive and the negative parts of π , respectively, and the minimum is taken over positive measures γ on $X \times X$ with marginals π_c and π_p (Santambrogio, 2015, Section 4.2). Combining this identity with Corollary 3, we obtain the dual problem in the form of Daskalakis et al. (2017).

Corollary 4 (Daskalakis et al. (2017)). *For $B = 1$ bidder whose type is distributed according to a smooth positive density ρ , the optimal revenue of the auctioneer (8) is equal to*

$$\inf_{\pi \succeq m} \min_{\gamma: \gamma_1 = \pi_c, \gamma_2 = \pi_p} \int \|x - y\|_1 d\gamma(x, y). \quad (34)$$

5 Simulations

In this section, we rely on numerical simulations to illustrate what the optima in the primal and dual problems look like in the benchmark case of $B \geq 2$ bidders with i.i.d. uniformly distributed values over $I = 2$ items.

5.1 Algorithm

Here we describe high-level ideas behind the algorithm. The detailed description and proofs can be found in Appendix E. As discussed in the introduction, finding a solution numerically is far from being straightforward: although the auctioneer’s

⁷More generally, there is a connection between Beckmann’s problem and congested optimal transportation problems of Monge-Kantorovich type; see the discussion in Appendix F.

problem is a linear program in a functional space, any reasonable discretization of it cannot be handled by modern LP solvers because of the curse of dimensionality. Indeed, if an agent can have n different values for each of I items, then the mechanism should specify an allocation and transfers for each of $(n^I)^B$ profiles of types which becomes computationally intractable already for two items, $n = 100$, and $B = 2$ agents or for $n = 10$, and $B = 4$ agents.

We escape the curse of dimensionality by dealing with the Rochet-Choné problem (11) which is equivalent to the auctioneer’s problem by Proposition 1. For n points in the discretization, the dimension of the Rochet-Choné problem is constant in the number of bidders B . This observation lies at the heart of algorithmic multi-to-single-agent reduction proposed (but not implemented) by Cai et al. (2012) and Alaei et al. (2019). The reduction in the dimensionality comes at the cost of complexity of the feasibility constraint: the classic form of this constraint by Border (1991) leads to exponentially many inequalities and the two papers propose distinct ad hoc constructions reducing this number to polynomial.

We rely on Border’s constraint in its majorization form (12) which is convex but non-linear. A natural linearization is suggested by the relation between majorization and martingales well-known to economists working on information design (Blackwell, 1951). We use this relation in the following form: a measure ν on $[0, 1]$ majorizes ν' if and only if there is a distribution γ on $[0, 1]^2$ with marginals ν on y and ν' on x and such that $\int y d\gamma(y | x) \geq x$ for γ -almost all x , where $\gamma(y | x)$ denotes the conditional measure on y given x (Shaked and Shanthikumar, 2007, Theorem 4.A.5).⁸

Considering (u, γ) as unknowns, we obtain a linear optimization problem equivalent to (11). Discretization of this problem leads to a number of constraints polynomial in n . In Appendix E, we demonstrate that the values of the discretized problems are guaranteed to converge to the true value as the discretization becomes finer and finer. Our approach is inspired by Ekeland and Moreno-Bromberg (2010) and, to the best of our knowledge, we are the first to obtain such approximation guarantees in multi-item auction design. To speed up the computation in practice, we adapt insights from Oberman (2013) to handle the incentive-compatibility constraint; see Appendix E for details.

5.2 Results

The algorithm was implemented in Python using the LP solver from Gurobi library. Simulations were run on Amazon EC2 instance m6i.16xlarge with 64 vCPUs with 3rd generation Intel Xeon Scalable cores and 256 GB of memory. For $n = 200 \times 200$

⁸Equivalently, there is a supermartingale (ξ, η) such that ξ is distributed according to ν' and η , according to ν .

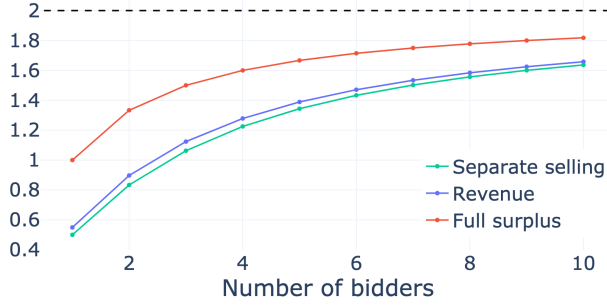


Figure 1: Revenue as a function of the number of bidders B for two items with i.i.d. values uniform on $[0, 1]$. Graphs from bottom to top: selling separately (light-green), selling optimally (blue), full surplus extraction (red), limit for $B \rightarrow \infty$ (the dashed line).

points in discretization, $I = 2$ items and $B = 2$ bidders, the computation required 83 minutes of real time and 20 hours of user time.

We focused on the case of two items with independent values uniformly distributed on $[0, 1]$. We computed how the optimal revenue depends on the number of bidders B . Figure 1 depicts this dependence. Naturally, the optimal revenue is lower-bounded by the revenue obtained from selling the items separately using Myerson’s optimal auction and upper-bounded by the revenue that the auctioneer would get if she could extract the full surplus.⁹ We see that the advantage from using the optimal auction is substantial for small number of bidders and it is maximal for $B = 2$ where the optimal mechanism increases the revenue by 5.7%. For large number of bidders, the use of optimal auction is not justified as selling the items separately leads to almost full surplus extraction.¹⁰

Let us take a closer look at the case of $B = 2$ bidders. The solution to the primal problem is shown in Figure 2 depicting the probability to receive the first item as a function of bidder’s values (x_1, x_2) , i.e., the optimal reduced allocation rule $\bar{P}_1^{\text{opt}}(x_1, x_2) = \frac{\partial}{\partial x_1} u^{\text{opt}}(x_1, x_2)$. The probability for the second item can be obtained by symmetry: $\bar{P}_2^{\text{opt}}(x_1, x_2) = \bar{P}_1^{\text{opt}}(x_2, x_1)$. The discontinuity that we see in Figure 2 correspond to the multi-dimensional reserve price: the minimal x_1 to receive a non-zero portion of the first item non-linearly depends on x_2 unless x_2 is high enough. Comparing Figure 2 to plots by Dütting et al. (2019), we see that neural-network approach not using multi-to-single-agent reduction is prone

⁹Revenue of Myerson’s auction run for each item separately is $2B \int_{0.5}^1 (2x - 1)x^{B-1} dx$ while the full surplus is $2 \left(1 - \frac{1}{B+1}\right)$.

¹⁰A posted price mechanism for the grand bundle extracts $1 - O\left(\frac{1}{B}\right)$ fraction of the full surplus, as $B \rightarrow \infty$.

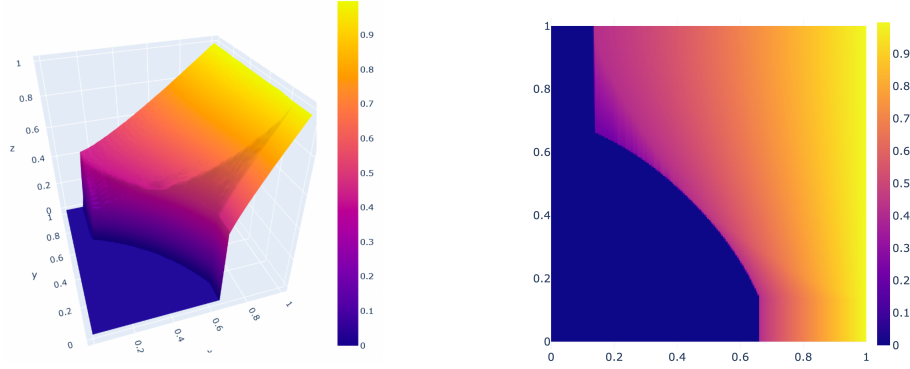


Figure 2: The probability to receive the first item as a function of bidder's values (x_1, x_2) in the optimal 2-bidder 2-item auction with i.i.d. values uniform on $[0, 1]$.

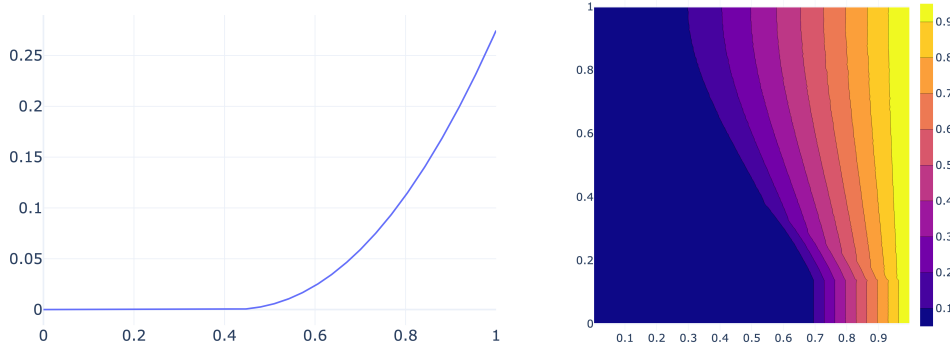


Figure 3: The optimal solution to the dual problem: functions $\varphi_1 = \varphi_2$ (left) and a contour plot of the first component of the vector field $c = (c_1, c_2)$ from Beckmann's problem (right).

to smoothing artefacts. The solution to the dual problem is shown in Figure 3. The contour plot demonstrates the first component c_1^{opt} of the optimal vector field c^{opt} ; the second component can be obtained by $c_2^{\text{opt}}(x_1, x_2) = c_1^{\text{opt}}(x_2, x_1)$. By the complementary slackness condition (28), we have $c_i^{\text{opt}}(x) \in \partial\varphi_i^{\text{opt}}\left(\frac{\partial u^{\text{opt}}}{\partial x_i}(x)\right)$ and so one could expect that the vector field inherits the discontinuity of ∂u . The optimal vector field turns out to be smooth because the optimal φ_i are zero in the discontinuity region.

References

D. R. Adams and L. I. Hedberg. *Function Spaces and Potential Theory*, volume 314 of *Grundlehren der mathematischen Wissenschaften*. Springer, 1999.

- S. Alaei, H. Fu, N. Haghpahan, J. Hartline, and A. Malekian. Efficient computation of optimal auctions via reduced forms. *Mathematics of Operations Research*, 44(3):1058–1086, 2019.
- C. Aliprantis and K. Border. *Infinite Dimensional Analysis: A Hitchhiker’s Guide (3rd Edition)*. Springer Berlin Heidelberg New York, 2006.
- T. Allen and C. Arkolakis. Trade and the topography of the spatial economy. *The Quarterly Journal of Economics*, 129(3):1085–1140, 2014.
- L. Ambrosio, N. Fusco, and D. Pallara. *Functions of bounded variation and free discontinuity problems*. Courier Corporation, 2000.
- I. Arieli, Y. Babichenko, R. Smorodinsky, and T. Yamashita. Optimal persuasion via bi-pooling. *Available at SSRN 3511516*, 2019.
- I. Arieli, Y. Babichenko, F. Sandomirskiy, and O. Tamuz. Feasible joint posterior beliefs. *Journal of Political Economy*, 129(9):2546–2594, 2021a.
- I. Arieli, Y. Babichenko, F. Sandomirskiy, and O. Tamuz. Persuasion as transportation. *online soon*, 2021b.
- M. Armstrong and J.-C. Rochet. Multi-dimensional screening:: A user’s guide. *European Economic Review*, 43(4-6):959–979, 1999.
- M. Babaioff, N. Immorlica, B. Lucier, and S. M. Weinberg. A simple and approximately optimal mechanism for an additive buyer. *Journal of the ACM (JACM)*, 67(4):1–40, 2020.
- M. Babaioff, Y. A. Gonczarowski, and N. Nisan. The menu-size complexity of revenue approximation. *Games and Economic Behavior*, 2021.
- A. Beck. *First-order methods in optimization*. SIAM, 2017.
- M. Beckmann. A continuous model of transportation. *Econometrica: Journal of the Econometric Society*, pages 643–660, 1952.
- D. Bergemann, B. A. Brooks, and S. Morris. Informationally robust optimal auction design. 2016.
- M. Bichler. *Market design: a linear programming approach to auctions and matching*. Cambridge University Press, 2017.
- D. Blackwell. Comparison of experiments. In *Proceedings of the Second Berkeley Symposium on Mathematical Statistics and Probability*, pages 93–102. University of California Press, 1951.

- J. Boerma, A. Tsyvinski, and A. P. Zimin. Sorting with team formation. Technical report, National Bureau of Economic Research, 2021.
- V. I. Bogachev and A. V. Kolesnikov. The Monge–Kantorovich problem: achievements, connections, and perspectives. *Russian Math. Surveys*, 67(5):785–890, Oct. 2012. doi: 10.1070/RM2012v067n05ABEH004808.
- K. C. Border. Implementation of reduced form auctions: A geometric approach. *Econometrica: Journal of the Econometric Society*, pages 1175–1187, 1991.
- Y. Cai, C. Daskalakis, and S. M. Weinberg. An algorithmic characterization of multi-dimensional mechanisms. In *Proceedings of the forty-fourth annual ACM symposium on Theory of computing*, pages 459–478, 2012.
- Y. Cai, N. R. Devanur, and S. M. Weinberg. A duality-based unified approach to bayesian mechanism design. *SIAM Journal on Computing*, 50(3):STOC16–160, 2019.
- O. Candogan and P. Strack. Optimal disclosure of information to a privately informed receiver. *arXiv preprint arXiv:2101.10431*, 2021.
- G. Carlier. Optimal transportation and economic applications. *Lecture Notes*, 2012.
- P.-A. Chiappori, R. J. McCann, and L. P. Nesheim. Hedonic price equilibria, stable matching, and optimal transport: equivalence, topology, and uniqueness. *Economic Theory*, 42(2):317–354, 2010.
- C. Daskalakis. Multi-item auctions defying intuition? *ACM SIGecom Exchanges*, 14(1):41–75, 2015.
- C. Daskalakis, A. Deckelbaum, and C. Tzamos. Strong duality for a multiple-good monopolist. *Econometrica*, 85(3):735–767, 2017.
- D. Dizdar and E. Kováč. A simple proof of strong duality in the linear persuasion problem. *Games and Economic Behavior*, 122:407–412, 2020.
- P. Dütting, Z. Feng, H. Narasimhan, D. Parkes, and S. S. Ravindranath. Optimal auctions through deep learning. In *International Conference on Machine Learning*, pages 1706–1715. PMLR, 2019.
- P. Dworczak and G. Martini. The simple economics of optimal persuasion. *Journal of Political Economy*, 127(5):1993–2048, 2019.

- I. Ekeland and S. Moreno-Bromberg. An algorithm for computing solutions of variational problems with global convexity constraints. *Numerische Mathematik*, 115(1):45–69, 2010.
- P. D. Fajgelbaum and E. Schaal. Optimal transport networks in spatial equilibrium. *Econometrica*, 88(4):1411–1452, 2020.
- A. Figalli, Y.-H. Kim, and R. J. McCann. When is multidimensional screening a convex program? *Journal of Economic Theory*, 146(2):454–478, 2011.
- A. Galichon. *Optimal transport methods in economics*. Princeton University Press, 2016.
- A. Galichon. A survey of some recent applications of optimal transport methods to econometrics. *arXiv preprint arXiv:2102.01716*, 2021.
- F. Gensbittel. Extensions of the cav (u) theorem for repeated games with incomplete information on one side. *Mathematics of Operations Research*, 40(1):80–104, 2015.
- A. Gershkov, B. Moldovanu, P. Strack, and M. Zhang. A theory of auctions with endogenous valuations. *Journal of Political Economy*, 129(4):1011–1051, 2021.
- Y. Giannakopoulos and E. Koutsoupias. Duality and optimality of auctions for uniform distributions. *SIAM Journal on Computing*, 47(1):121–165, 2018.
- N. Gozlan, C. Roberto, P.-M. Samson, and P. Tetali. Kantorovich duality for general transport costs and applications. *Journal of Functional Analysis*, 273(11):3327–3405, 2017.
- N. Guillen and R. McCann. Five lectures on optimal transportation: Geometry, regularity and applications. In *Analysis and Geometry of Metric Measure Spaces: Lecture Notes of the 50th Séminaire de Mathématiques Supérieures (SMS), Montréal, 2011*, chapter 6, pages 145–180. CRM Proceedings & Lecture Notes, 2013. doi: 10.1090/crmp/056/06.
- Y. Guo and E. Shmaya. Costly miscalibration. *Theoretical Economics*, 16(2):477–506, 2021.
- N. Haghpahan and J. Hartline. When is pure bundling optimal? *The Review of Economic Studies*, 88(3):1127–1156, 2021.
- S. Hart and N. Nisan. Approximate revenue maximization with multiple items. *Journal of Economic Theory*, 172:313–347, 2017.

- S. Hart and P. J. Reny. Implementation of reduced form mechanisms: a simple approach and a new characterization. *Economic Theory Bulletin*, 3(1):1–8, 2015.
- S. Hart and P. J. Reny. The better half of selling separately. *ACM Transactions on Economics and Computation (TEAC)*, 7(4):1–18, 2019.
- J. D. Hartline. Mechanism design and approximation. *Book draft*, 122:1, 2013.
- P. Henry-Labordère. *Model-free hedging: A martingale optimal transport viewpoint*. CRC Press, 2017.
- A. Kleiner and A. Manelli. Strong duality in monopoly pricing. *Econometrica*, 87(4):1391–1396, 2019.
- A. Kleiner, B. Moldovanu, and P. Strack. Extreme points and majorization: Economic applications. *Econometrica*, 89(4):1557–1593, 2021.
- Ç. Koçyiğit, G. Iyengar, D. Kuhn, and W. Wiesemann. Distributionally robust mechanism design. *Management Science*, 66(1):159–189, 2020.
- A. Kolotilin. Optimal information disclosure: A linear programming approach. *Theoretical Economics*, 13(2):607–635, 2018.
- A. M. Manelli and D. R. Vincent. Bundling as an optimal selling mechanism for a multiple-good monopolist. *Journal of Economic Theory*, 127(1):1–35, 2006.
- R. McCann and K. S. Zhang. On concavity of the monopolist’s problem facing consumers with nonlinear price preferences. *Communication on pure and applied mathematics*, 72(7):1386–1423, 2019. doi: <https://doi.org/10.1002/cpa.21817>.
- R. B. Myerson. Optimal auction design. *Mathematics of operations research*, 6(1):58–73, 1981.
- A. Nikzad. Constrained majorization: Applications in mechanism design. *Available at SSRN 4030091*, 2022.
- G. Nöldeke and L. Samuelson. The implementation duality. *Econometrica*, 86(4):1283–1324, 2018.
- A. M. Oberman. A numerical method for variational problems with convexity constraints. *SIAM Journal on Scientific Computing*, 35(1):A378–A396, 2013.
- J.-C. Rochet and P. Choné. Ironing, sweeping, and multidimensional screening. *Econometrica*, pages 783–826, 1998.

- J.-C. Rochet and L. A. Stole. The economics of multidimensional screening. *Econometric Society Monographs*, 35:150–197, 2003.
- R. T. Rockafellar. *Convex analysis*. Princeton university press, 2015.
- F. Santambrogio. Optimal transport for applied mathematicians. *Birkäuser, NY*, 55(58-63):94, 2015.
- M. Shaked and J. G. Shanthikumar. *Stochastic orders*. Springer Science & Business Media, 2007.
- S. Steinerberger and A. Tsyvinski. Tax mechanisms and gradient flows. Technical report, National Bureau of Economic Research, 2019.
- V. Strassen. The existence of probability measures with given marginals. *The Annals of Mathematical Statistics*, 36(2):423–439, 1965.
- A. Suzdaltsev. An optimal distributionally robust auction. *arXiv preprint arXiv:2006.05192*, 2020.
- C. Villani. *Optimal transport: old and new*, volume 338. Springer, 2009.
- R. V. Vohra. *Mechanism design: a linear programming approach*, volume 47. Cambridge University Press, 2011.
- A. C.-C. Yao. Dominant-strategy versus bayesian multi-item auctions: Maximum revenue determination and comparison. In *Proceedings of the 2017 ACM Conference on Economics and Computation*, pages 3–20, 2017.

A Convex analysis basics

Throughout the paper, we consider convex functions on $[0, 1]$, X , \mathbb{R} , or $\mathbb{R}^{\mathcal{I}}$ taking values in $\mathbb{R} \cup \{+\infty\}$. Here we briefly remind the reader some important facts and definitions.

The subdifferential of a convex function f is defined by

$$\partial f(x) = \{\tau : f(y) \geq f(x) + \langle \tau, y - x \rangle, \quad \forall y\}. \quad (35)$$

Partial derivatives of $f(x)$ at $x \in \mathbb{R}^d$ (if exist) are denoted by

$$f_{x_i}(x) = \frac{\partial f}{\partial x_i}(x).$$

The gradient $\nabla f(x)$ is the vector of partial derivatives

$$\nabla f(x) = (f_{x_i}(x))_{i=1,\dots,d}.$$

If f is differentiable at x , then the subdifferential $\partial f(x)$ consists of just one element: $\partial f(x) = \{\nabla f(x)\}$. By the Alexandrov theorem, a convex function is twice differentiable except for a set of zero Lebesgue measure; see (Villani, 2009, Theorem 14.1). In particular, the gradient $\nabla f(x)$ is defined almost everywhere and hence the integrals of the gradient with respect to an absolutely continuous measure are well-defined even for non-smooth f .

The Legendre transform also known as Fenchel's conjugate of a convex function f is a convex function given by

$$f^*(y) = \sup_x (\langle x, y \rangle - f(x)).$$

We will widely use the Fenchel inequality

$$f(x) + f^*(y) \geq \langle x, y \rangle \quad (36)$$

and the corresponding “complementary slackness” condition taking the following form: $f(x) + f^*(y) = \langle x, y \rangle$ if and only if $y \in \partial f(x)$ and $x \in \partial f^*(y)$.

B Proof of Proposition 1 (Rochet-Choné representation of auctioneer's problem)

Recall that the auctioneer's problem (8) is to maximize the revenue

$$\int_{X^B} \left(\sum_{b \in B} T_b((x_b)_{b \in B}) \right) \cdot \left(\prod_{b \in B} \rho(x_b) \right) dx_1 \cdots dx_B \quad (37)$$

over individually-rational Bayesian incentive-compatible feasible mechanisms. The multi-bidder Rochet-Choné problem (11) is to maximize

$$B \cdot \int_X (\langle \nabla u(x), x \rangle - u(x)) \rho(x) dx \quad (38)$$

over convex non-decreasing functions $u : X \rightarrow \mathbb{R}_+$ with $u(0) = 0$ and such that, for all $i \in \mathcal{I}$,

$$\frac{\partial u}{\partial x_i}(x) \preceq \xi^{B-1}, \quad (39)$$

where $\chi \in X$ is distributed with the density ρ and ξ is uniformly distributed on $[0, 1]$. Our goal is to prove that the values of the two optimization problems coincide and both maxima are attained. The proof relies on a sequence of lemmas.

It will be convenient to work with a version of the Rochet-Choné problem where the constraint $u(0) = 0$ is relaxed (the requirements that u is non-decreasing and takes only non-negative values remain).

Lemma 1. *The constraint $u(0) = 0$ in the Rochet-Choné problem (38) can be relaxed without affecting the value and whether the optimum is attained or not.*

Proof. It is enough to show that for any feasible u with $u(0) > 0$, there is a feasible \tilde{u} with $\tilde{u}(0) = 0$ and the same or higher value of the objective. Defining $\tilde{u}(x) = u(x) - u(0)$ completes the proof. \square

Let us demonstrate that, for any feasible solution to (37), there is a feasible solution to (38) with the relaxed constraint $u(0) = 0$ and vice versa. This will imply that the two problems have the same values and, moreover, the optima are attained or not attained simultaneously.

Lemma 2. *For any individually-rational Bayesian incentive-compatible feasible mechanism (P, T) from (37), there exists a function u satisfying all the constraints of the Rochet-Choné problem (38) except for, possibly, $u(0) = 0$ and such that the revenue of (P, T) is equal to the value of (38) at u .*

Proof. Consider the symmetrization of the mechanism (P, T) over all permutations of bidders:

$$P_{b,i}^{\text{sym}}((x_k)_{k \in \mathcal{B}}) = \frac{1}{|N|!} \sum_{\sigma \in S_{\mathcal{B}}} P_{\sigma(b),i}((x_{\sigma(b)})_{b \in \mathcal{B}})$$

$$T_i^{\text{sym}}((x_k)_{k \in \mathcal{B}}) = \frac{1}{|N|!} \sum_{\sigma \in S_{\mathcal{B}}} T_{\sigma(b)}((x_{\sigma(b)})_{b \in \mathcal{B}}),$$

where $S_{\mathcal{B}}$ denotes the set of all permutations σ of the set of bidders \mathcal{B} . The symmetrization $(P^{\text{sym}}, T^{\text{sym}})$ results in the same revenue and inherits all the properties of (P, T) . By symmetry, all the bidders contribute equally to the revenue and so the revenue can be rewritten as

$$B \cdot \int_{X^{\mathcal{B}}} T_1^{\text{sym}}((x_b)_{b \in \mathcal{B}}) \cdot \left(\prod_{b \in \mathcal{B}} \rho(x_b) \right) dx_1 \cdots dx_B = B \cdot \int_X \overline{T}_1^{\text{sym}}(x) \cdot \rho(x) dx,$$

where $(\overline{P}_b^{\text{sym}}, \overline{T}_b^{\text{sym}})$ denotes bidder b 's reduced mechanism (reduced mechanisms are the same for all the bidders by symmetry). Define $u(x)$ as the average utility of a bidder of type $x \in X$ in $(P^{\text{sym}}, T^{\text{sym}})$:

$$u(x) = \left\langle \overline{P}_1^{\text{sym}}(x), x \right\rangle - \overline{T}_1^{\text{sym}}(x).$$

By the definition of incentive compatibility,

$$\langle \overline{P}_b^{\text{sym}}(y), y \rangle - \overline{T}_b^{\text{sym}}(y) \geq \langle \overline{P}_b^{\text{sym}}(x), y \rangle - \overline{T}_b^{\text{sym}}(x).$$

Thus

$$u(y) \geq u(x) + \langle \overline{P}_1^{\text{sym}}(x), y - x \rangle. \quad (40)$$

We conclude that $u(y) = \max_{x \in X} \left(u(x) + \langle \overline{P}_1^{\text{sym}}(x), y - x \rangle \right)$ and, hence, u is a convex function as the pointwise maximum of a family of affine functions. Comparing (40) to the definition of the subdifferential of a convex function (35), we see that $\overline{P}_1^{\text{sym}}(x)$ belongs to the subdifferential $\partial u(x)$. For Lebesgue-almost all x , the gradient $\nabla f(x)$ of a convex function f is well-defined and the subdifferential $\partial f(x)$ coincides with the singleton $\{\nabla f(x)\}$. Therefore,

$$\overline{P}_1^{\text{sym}}(x) = \nabla u(x) \quad (41)$$

for almost all x . By the definition of u , we can express $\overline{T}_1^{\text{sym}}(x)$ as follows:

$$\overline{T}_1^{\text{sym}}(x) = \langle \overline{P}_1^{\text{sym}}(x), x \rangle - u(x) = \langle \nabla u(x), x \rangle - u(x),$$

where the second equality holds almost everywhere. Thus

$$B \cdot \int_X \overline{T}_1^{\text{sym}}(x) \cdot \rho(x) dx = B \cdot \int_X \left(\langle \nabla u(x), x \rangle - u(x) \right) \rho(x) dx,$$

i.e., u gives the same value to (38) as (P, T) to (37). We already know that u is convex. It remains to check that u is non-negative, monotone, and that it satisfies the majorization constraint (39). Non-negativity is immediate since, by the definition, $u \geq 0$ is equivalent to individual rationality of $(P^{\text{sym}}, T^{\text{sym}})$. By (41), u is a convex function with the gradient having non-negative components almost everywhere. Hence, u is non-decreasing.

To check (39), note that P^{sym} can be seen as a family of I allocation rules $P^{\text{sym}, i} = (P_{b,i})_{b \in \mathcal{B}}$, one for each item $i \in \mathcal{I}$. The reduced allocation $\overline{P}_b^{\text{sym}, i}: X \rightarrow \mathbb{R}_+$ for this one-item rule is equal to the corresponding component of $\overline{P}_b^{\text{sym}}$.

By the Border theorem in the form of [Hart and Reny \(2015\)](#), a function $f: X \rightarrow [0, 1]$ coincides with a reduced form \overline{Q}_b of some bidder-symmetric feasible one-item mechanism (Q, S) if and only if $f(\chi) \preceq \xi^{B-1}$, where $\chi \in X$ is distributed with the density ρ and ξ is uniformly distributed on $[0, 1]$.

Applying Border's theorem, we conclude that

$$\overline{P}_b^{\text{sym}, i}(\chi) \preceq \xi^{B-1}.$$

Since $\overline{P_b^{\text{sym},i}}$ is equal to $\overline{P_{b,i}^{\text{sym}}}$ and the latter coincides with $\frac{\partial u}{\partial x_i}$ by (41), we obtain the desired condition (39). To summarize, for any (P, T) , we constructed u giving the same value to the Rochet-Choné problem and satisfying all its constraints (without $u(0) = 0$ which was shown to be redundant). \square

Now we show how to construct (P, T) starting from u .

Lemma 3. *For any u satisfying the constraints of the Rochet-Choné problem (38) except for, possibly, $u(0) = 0$, there exists an individually-rational Bayesian incentive-compatible feasible mechanism (P, T) such that its revenue (37) is equal to the value of (38) at u .*

Proof. The proof reverses the construction used to prove Lemma 2. Consider a function f^i equal to the component of u 's gradient corresponding to an item $i \in \mathcal{I}$, i.e., $f^i = \frac{\partial u}{\partial x_i}$. We assume that f^i is defined for all $x \in X$: whenever the gradient is not well-defined, we select f^i arbitrarily so that the vector $f = (f^i(x))_{i \in \mathcal{I}}$ belongs to the subdifferential $\partial u(x)$. The function f^i is non-negative as u is monotone and $f(\chi)$ is majorized by ξ^{B-1} since $\frac{\partial u}{\partial x_i}(\chi)$ is. Thus, by Border's theorem, there exists a feasible one-item allocation $P^i: X^B \rightarrow \mathbb{R}_+^B$ such that $\overline{P_b^i} = f^i(x)$ for any bidder b .

Define the mechanism (P, T) as follows. The items are allocated by applying P^i to each $i \in \mathcal{I}$, i.e., $P_{b,i} = P_b^i$. The transfers T are given by

$$T_b((x_b)_{b \in \mathcal{B}}) = \langle f(x_b), x_b \rangle - u(x_b).$$

Thus (P, T) is feasible and the reduced mechanisms satisfy

$$\overline{P}_b(x) = f(x) \quad \text{and} \quad \overline{T}_b(x) = \langle f(x), x \rangle - u(x) \quad (42)$$

for any bidder b . As $f = \nabla u$ almost everywhere, the second identity in (42) implies that the revenue of (P, T) coincides with the value of (38) at u . It remains to check that (P, T) is individually rational and Bayesian incentive-compatible. Individual rationality reads as $\langle \overline{P}_b(x), x \rangle - \overline{T}_b(x) \geq 0$. By (42), the left-hand side equals $u(x)$ and so individual rationality follows from non-negativity of u . To show incentive-compatibility, recall that $f(x)$ is an element of the subdifferential of u and so

$$u(x') \geq u(x) + \langle f(x), x' - x \rangle.$$

By (42), this inequality rewrites as

$$\langle \overline{P}_b(x_b), x_b \rangle - \overline{T}_b(x_b) \geq \langle \overline{P}_b(x'_b), x_b \rangle - \overline{T}_b(x'_b),$$

which is exactly the condition of incentive-compatibility for (P, T) . Thus (P, T) is an individually-rational Bayesian incentive-compatible feasible mechanism with revenue equal to the value of the Rochet-Choné objective at u . \square

The above lemmas imply that the values of problems (37) and (38) coincide. To prove that the optima are attained we need the following pair of lemmas.

Let $\mathcal{U}_{\text{Lip},1}$ be the set of all convex non-decreasing functions $u: X \rightarrow \mathbb{R}_+$ with $u(0) = 0$ satisfying 1-Lipschitz condition $|u(x) - u(y)| \leq \sum_{i \in \mathcal{I}} |x_i - y_i|$ and endowed with the topology of the set of continuous functions.

Lemma 4. *The Rochet-Choné objective*

$$B \cdot \int_X \left(\langle \nabla u(x), x \rangle - u(x) \right) \rho(x) \, dx \quad (43)$$

is a continuous functional over the set $\mathcal{U}_{\text{Lip},1}$.

Proof. Let $u^{(n)} \rightarrow u$ be a uniformly convergent sequence of functions from $\mathcal{U}_{\text{Lip},1}$. Any limiting point y of any sequence $\{y^{(n)}\}$ such that $y^{(n)} \in \partial u^{(n)}(x)$, belongs to $\partial u(x)$. Indeed, for every z one has

$$u^{(n)}(z) \geq u^{(n)}(x) + \langle y^{(n)}, z - x \rangle$$

by definition of the subdifferential. From the convergence $u^{(n)} \rightarrow u$ and $y^{(n)} \rightarrow y$ one gets

$$u(z) \geq u(x) + \langle y, z - x \rangle$$

for all z , hence, $y \in \partial u(x)$. Since the subdifferential $\partial u^{(n)}(x)$ coincides with the gradient $\{\nabla u^{(n)}(x)\}$ for all n and almost all x , we get that $\nabla u^{(n)}(x)$ converges to $\nabla u(x)$ almost everywhere.

Thus the convergence of $u^{(n)} \rightarrow u$ in the topology of continuous functions implies the convergence of integrands in (43) almost everywhere. To deduce the continuity of the functional, we need to show that taking the limit commutes with the integration. This follows from the Lebesgue dominated convergence theorem. To apply this theorem, it remains to show that, in addition to convergence almost everywhere, the sequence of integrands is bounded. Since $u^{(n)}$ is a convergent sequence of continuous functions, $\sup_{n \in \mathbb{N}, x \in X} |u^{(n)}(x)| < \infty$ and $\langle \nabla u^{(n)}(x), x \rangle$ is bounded by I thanks to the 1-Lipschitz property. We obtain boundedness of the sequence of integrands and conclude that the Rochet-Choné objective is continuous. \square

The next lemma shows that the feasible set in the Rochet-Choné problem is a compact subset of $\mathcal{U}_{\text{Lip},1}$.

Lemma 5. *The set of convex non-decreasing functions $u: X \rightarrow \mathbb{R}_+$ with $u(0) = 0$ satisfying the majorization condition (39) is a compact subset of $\mathcal{U}_{\text{Lip},1}$.*

Proof. Since the upper bound in (39) is a random variable taking values in $[0, 1]$, we see that the gradient of a function u from the statement of the lemma takes values in $[0, 1]^Z$ and thus such u belongs to $\mathcal{U}_{\text{Lip},1}$.

To prove the compactness of the set of such u , note that the set $\mathcal{U}_{\text{Lip},1}$ is a set of uniformly bounded uniformly equicontinuous functions. Hence, any sequence of functions from $\mathcal{U}_{\text{Lip},1}$ contains a convergent subsequence. Thus, to prove compactness of a subset of $\mathcal{U}_{\text{Lip},1}$, it is enough to check closedness of this subset. If $u^{(n)} \in \mathcal{U}_{\text{Lip},1}$ is a sequence of functions converging uniformly to some u , we know that their gradients $\nabla u^{(n)}$ converge to ∇u almost surely (see the proof of Lemma 4). As the gradients are bounded, their distributions converge weakly. Therefore, if $u^{(n)}$ satisfy the majorization condition (39), it is also satisfied by the limit u . We obtain closedness and thus compactness. \square

Now the proof of Proposition 1 is almost immediate.

Proof of Proposition 1. By Lemma 1, the value of the Rochet-Choné problem does not change if we relax the constraint $u(0) = 0$. Lemma 2 implies that the value of the Rochet-Choné problem with the relaxed constraint is at least the value of the auctioneer's problem, while Lemma 3 gives the opposite inequality. Thus the values of the Rochet-Choné and the auctioneer's problems are equal.

By Lemmas 4 and 5, the Rochet-Choné problem can be seen as maximization of a continuous functional over a compact set. Therefore, this problem attains its optimum, i.e., the optimal u exists. By Lemma 3, we can find a mechanism (P, T) such that the auctioneer's revenue is the same as the value of the Rochet-Choné objective. Thus the optimum in the auctioneer's problem is also attained, i.e., the optimal auction exists as well. \square

C Duality and proofs

In this section, we prove Theorems 1 and Theorem 2 establishing the strong dual to the auctioneer's problem. The proof is split into two big parts. First, we derive a partial dual problem internalizing Border's feasibility constraint (39). This problem is interpreted as a problem of a monopolist facing adversarial production costs; a result which may be of independent interest. In terms of this problem, we formulate a novel a priori bound on solutions, our main technical tool. Next, relying on this tool, we proceed with proving the theorems.

By Proposition 1, we know that the auctioneer's problem is equivalent to the multi-bidder Rochet-Choné problem where the distribution of u 's gradient is majorized by a particular distribution depending on the number of bidders. As our arguments do not depend on the exact form of the dominating distribution, in this

section we allow for general majorizing distributions and, consequently, the results of this section extend Theorems 1 and Theorem 2 to general majorization.

Let us describe the generalized Rochet-Choné problem and introduce some useful notation along the way. Recall that \mathcal{I} is the set of $I \geq 1$ items and $X = [0, 1]^{\mathcal{I}}$ is the set of bidders' types endowed with a density ρ . We will denote the corresponding distribution by μ so that

$$d\mu(x) = \rho(x) dx$$

and assume that ρ is strictly positive on X .

For a convex function u on X , its gradient is well-defined for almost all x ; see Appendix A. For the gradient's component $\frac{\partial}{\partial x_i} u(x)$ we will sometimes use compact notation $u_{x_i}(x)$. We denote by ν_i the distribution of the gradient's i 'th component $u_{x_i}(\chi)$ assuming that χ has distribution μ .

Rochet-Choné problem with general majorization: *given an absolutely continuous probability measure μ on X and a collection of probability measures $(\eta_i)_{i \in \mathcal{I}}$ on \mathbb{R}_+ , maximize*

$$\int_X \left(\langle \nabla u(x), x \rangle - u(x) \right) d\mu(x) \quad (44)$$

over convex non-decreasing functions $u: X \rightarrow \mathbb{R}_+$ with $u(0) = 0$ and such that for all $i \in \mathcal{I}$

$$\nu_i \preceq \eta_i, \quad (45)$$

where ν_i is the distribution of u_{x_i} .

If all η_i are the same and coincide with the distribution of ξ^{B-1} with ξ uniform on $[0, 1]$, then the problem (44) coincides with the multibidder Rochet-Choné problem (11) up to a factor B in the objective. By Proposition 1, for such choice of η_i , the value of (44) is equal to $\frac{1}{B}$ of the optimal revenue in the auctioneer's problem with B bidders.

C.1 Auctioneer's problem as monopolist's problem with adversarial production costs

Consider a monopolist selling $I \geq 1$ items $i \in \mathcal{I}$ to one buyer whose type x is distributed according to some measure μ on $X = [0, 1]^{\mathcal{I}}$ with density ρ . In contrast to the single-bidder setting considered in Sections 2 and 3, these items have not yet been produced and so deciding on the amount to produce is a part of the monopolist's problem. We assume that the production costs are separable across items and, for each item $i \in \mathcal{I}$, are given by a convex non-decreasing function φ_i . The presence of the production costs $\sum_{i \in \mathcal{I}} \varphi_i(P_i)$ replace the feasibility constraint $P_i \leq 1$ of the monopolist's problem considered in Section 3.1. That model corresponds to a particular case of φ_i equal to 0 on $[0, 1]$ and $+\infty$ outside.

Monopolist's problem with production costs. For each item $i \in \mathcal{I}$, convex non-decreasing production costs $\varphi_i : [0, \infty) \rightarrow \mathbb{R}_+ \cup \{+\infty\}$ with $\varphi_i(0) = 0$ are given. The monopolist aims to maximize the total revenue consisting of the buyer's payment minus the production costs

$$\Phi(u, (\varphi_i)_{i \in \mathcal{I}}) = \int_X \left(\langle \nabla u(x), x \rangle - u(x) \right) d\mu(x) - \sum_{i \in \mathcal{I}} \int_X \varphi_i \left(\frac{\partial u}{\partial x_i}(x) \right) d\mu(x) \quad (46)$$

over convex non-decreasing functions $u : X \rightarrow \mathbb{R}_+$ with¹¹ $u(0) = 0$.

Let $\text{Rev}^{\text{opt}}[(\varphi_i)_{i \in \mathcal{I}}]$ be the value of the problem (46), i.e., the maximal revenue the monopolist can achieve. Since the zero mechanism corresponding to $u \equiv 0$ is feasible, the maximal revenue is non-negative, however, it may be infinite, e.g., if the costs are zero and so the monopolist has an incentive to increase production infinitely.

Consider an adversary who aims to minimize the monopolist's revenue by selecting the production costs but is penalized for choosing high costs. The adversary's objective is to minimize

$$\text{Rev}^{\text{opt}}[(\varphi_i)_{i \in \mathcal{I}}] + \sum_{i \in \mathcal{I}} \int \varphi_i(z) d\eta_i(z) \quad (47)$$

for some given measures η_i .

Theorem 3. Let η_i be probability measures on $[0, 1]$ such that $\eta_i([t, 1]) > 0$ for any $t > 0$. Then the following assertions hold:

- The value of the Rochet-Choné problem with general majorization (44) coincides with the optimal value achieved by the adversary in the minimization problem (47).

In particular, if all η_i are equal to the the distribution of ξ^{B-1} with ξ uniformly distributed on $[0, 1]$, the value of (47) coincides with $\frac{1}{B}$ fraction of the auctioneer's optimal revenue (8) for B bidders.

- The optimum in (47) is attained, i.e., the adversary has an optimal strategy given by lower semicontinuous functions $(\varphi_i^{\text{opt}})_{i \in \mathcal{I}}$.

Let us formulate the result paying attention to functional classes to which u and φ_i belong.

¹¹One can show that this problem is equivalent to maximization of $\int_X (T(x) - \sum_i \varphi_i(P_i)) d\mu(x)$ over individually-rational Bayesian incentive-compatible mechanisms $(P, T) : X \rightarrow \mathbb{R}_+^{\mathcal{I}} \times \mathbb{R}$ (the argument repeats the proof of Proposition 1). We do not rely on this equivalence.

Denote by $\mathcal{U}_{\text{Lip},K}$ the set of non-decreasing convex functions u on X that have $u(0) = 0$ and are K -Lipschitz in the l^1 -norm, i.e. $|u(x) - u(x')| \leq K \sum_i |x_i - x'_i|$. Note that monotonicity and K -Lipschitz properties together are equivalent to the following inequality on partial derivatives

$$0 \leq u_{x_i} \leq K, \quad \forall i \in \mathcal{I},$$

that must hold almost everywhere in X . For a probability measure η on $[0, 1]$, denote by $\mathcal{U}_{\eta,+\infty}$ the set of convex non-decreasing lower semicontinuous functions $\varphi: \mathbb{R}_+ \rightarrow \mathbb{R}_+ \cup \{+\infty\}$ such that $\varphi(0) = 0$, the integral $\int_0^1 \varphi(z) d\eta(z) < \infty$, and $\varphi(z) = +\infty$ for $z > 1$.

Formally, we prove the following identity

$$\begin{aligned} \max_{u \in \mathcal{U}_{\text{Lip},1}, \nu_i \leq \eta_i} \int (\langle x, \nabla u \rangle - u(x)) d\mu &= \\ &= \min_{\varphi_i \in \mathcal{U}_{\eta_i,+\infty}} \max_{u \in \mathcal{U}_{\text{Lip},1}} \left[\Phi(u, (\varphi_i)_{i \in \mathcal{I}}) + \sum_{i \in \mathcal{I}} \int_0^1 \varphi_i(x) d\eta_i(x) \right]. \end{aligned} \quad (48)$$

The proof of Theorem 3 is contained in the next subsection. The high-level idea is to apply a functional minimax theorem to the Lagrangian internalizing the majorization constraint (45). Indeed, interpret $\Phi(u, (\varphi_i)_{i \in \mathcal{I}})$ as the payoff function in a zero-sum game. The maximizer selects u , while the minimizer picks $(\varphi_i)_{i \in \mathcal{I}}$. The minimizer can infinitely penalize the maximizer for a violation of the majorization constraint (45). On the other hand, if the constraint is not violated, the best the minimizer can do is to select $\varphi_i \equiv 0$ on $[0, 1]$ for all i making the payoff equal to the objective of the Rochet-Choné problem with general majorization (44). We conclude that the maxinf-value of the game coincides with the value of the Rochet-Choné problem (44). Similarly, one can show that inf max-value is the optimal value of the adversary's problem (47). Next we apply the following functional minimax theorem which can be found in (Adams and Hedberg, 1999, Theorem 2.4.1).

Theorem 4. *Let X, Y be convex subsets of linear topological spaces. We assume, in addition, that X is a compact Hausdorff space. Let $f: X \times Y \rightarrow (-\infty, +\infty]$ be a function that is lower semicontinuous in x for every $y \in Y$, convex in x , and concave in y . Then*

$$\min_{x \in X} \sup_{y \in Y} f(x, y) = \sup_{y \in Y} \min_{x \in X} f(x, y).$$

By this theorem, we conclude that the maxinf and inf max values coincide. This gives us the first item of Theorem 3. We note that, in contrast to typical

game-theoretic derivations of dual problems, the payoff function Φ is not affine in the strategy u of the maximizer. However, Φ is convex in u which is enough for Theorem 4.

This gives the result with infimum over φ_i instead of minimum. Proving that the minimum is attained is the most difficult part of the proof as the set of minimizer's strategies is not compact and so we cannot use the standard compactness arguments.

C.2 Proof of Theorem 3

In addition to $\mathcal{U}_{\text{Lip},K}$ and $\mathcal{U}_{\eta,+\infty}$ defined above, we will need the following functional spaces:

- \mathcal{U}^p , $p \geq 1$, is the set of non-decreasing convex functions $u: X \rightarrow \mathbb{R}_+$ with $u(0) = 0$ and such that $\int_X |\nabla u|^p d\mu < \infty$, i.e., the gradient of u belongs to $L^p(\mu)$.
- $\mathcal{U}_{\mathbb{R}_+}$ is the set of convex non-decreasing lower semicontinuous functions $\varphi: \mathbb{R}_+ \rightarrow \mathbb{R}_+ \cup \{+\infty\}$ with $\varphi(0) = 0$ such that there exists $t_0 \in \mathbb{R}_+$ with

$$\varphi(t_0) > t_0.$$

Note that lower semicontinuity withing this class simply means that

$$\lim_{s \rightarrow t-} \varphi(s) = \varphi(t),$$

where $t = \inf\{s : \varphi(s) = +\infty\}$.

- $\mathcal{U}_{[0,1]}^{+\infty}$ is the set of all functions $\varphi \in \mathcal{U}_{\mathbb{R}_+}$ such that $\varphi_i(t)$ is finite for $t < 1$ and equal to $+\infty$ for $t > 1$.

The following two simple lemmas provide compactness and continuity properties needed for the proof of Theorem 3.

Lemma 6. *For any given $K \geq 0$, the set $\mathcal{U}_{\text{Lip},K}$ is compact in the uniform convergence topology.*

Proof. The compactness follows from the Arzelà–Ascoli theorem and the obvious fact that convexity and monotonicity are preserved under uniform convergence. \square

The following lemma extends the continuity of the objective obtained in Lemma 4 in the presence of functions φ_i .

Lemma 7. *For any tuple of $\varphi_i \in \mathcal{U}_{\mathbb{R}_+}$ the functional $\Phi(u, \varphi_i)$ is upper semicontinuous in u in the uniform convergence topology on $\mathcal{U}_{\text{Lip}, K}$ for every $K > 0$.*

If, in addition, φ_i do not take value $+\infty$ and are continuous, then $\Phi(u, \varphi_i)$ is continuous in u in the uniform convergence topology on $\mathcal{U}_{\text{Lip}, K}$ for every $K > 0$.

Proof. As it was demonstrated in the proof of Lemma 4, if $u^{(n)} \rightarrow u$ is a uniformly convergent sequence of Lipschitz convex non-decreasing functions on X , then the gradients $\nabla u^{(n)}$ converge to $\nabla u(x)$ almost everywhere. Thus, by the Fatou lemma (Theorem 11.20 in Aliprantis and Border (2006)), it is sufficient to check that

$$\overline{\lim}_n \left[\langle x, \nabla u^{(n)}(x) \rangle - u^{(n)}(x) - \sum_{i \in \mathcal{I}} \varphi_i(u_{x_i}^{(n)}(x)) \right] \leq \langle x, \nabla u(x) \rangle - u(x) - \sum_{i \in \mathcal{I}} \varphi_i(u_{x_i}(x))$$

μ -a.e. and $\langle x, \nabla u^{(n)}(x) \rangle - u^{(n)}(x) - \sum_{i \in \mathcal{I}} \varphi_i(u_{x_i}^{(n)}(x)) \leq C$ for some C . The first inequality follows immediately from μ -a.e. convergence and lower semicontinuity of φ_i . Next, since $u^{(n)}, u_{x_i}^{(n)}$ are nonnegative, one has

$$\langle x, \nabla u^{(n)}(x) \rangle - u^{(n)}(x) - \sum_{i \in \mathcal{I}} \varphi_i(u_{x_i}^{(n)}(x)) \leq \langle x, \nabla u^{(n)}(x) \rangle \leq IK.$$

The second statement of the lemma follows from the Lebesgue dominated convergence theorem applied as in Lemma 4. \square

Now we are ready to prove Theorem 3.

Proof of Theorem 3.

Step 1. The optimum in the left-hand side of (48) is attained. Indeed, by Proposition 1, we know that the value of the Rochet-Choné problem attains its value and this problem coincides with the left-hand side of (48). The proposition establishes the result for a particular choice of majorizing measures ν_i originating from Border's condition, but the extension to arbitrary ν_i is straightforward.

Step 2. Let us rewrite our problem in the minimax form and apply Theorem 4:

$$\begin{aligned} & \max_{u \in \mathcal{U}^1, \nu_i \leq \eta_i} \int (\langle x, \nabla u \rangle - u(x)) d\mu = \\ &= \max_{u \in \mathcal{U}_{\text{Lip}, 1}} \inf_{\varphi_i \in \mathcal{U}_{\eta_i, +\infty}} \left[\int \left(\langle x, \nabla u \rangle - u(x) - \sum_{i \in \mathcal{I}} \varphi_i(u_{x_i}) \right) d\mu + \sum_{i \in \mathcal{I}} \int_0^1 \varphi_i(x) d\eta_i(x) \right] \\ &= \inf_{\varphi_i \in \mathcal{U}_{\eta_i, +\infty}} \max_{u \in \mathcal{U}_{\text{Lip}, 1}} \left[\int \left(\langle x, \nabla u \rangle - u(x) - \sum_{i \in \mathcal{I}} \varphi_i(u_{x_i}) \right) d\mu + \sum_{i \in \mathcal{I}} \int_0^1 \varphi_i(x) d\eta_i(x) \right] \\ &= \inf_{\varphi_i \in \mathcal{U}_{\eta_i, +\infty}} \max_{u \in \mathcal{U}_{\text{Lip}, 1}} \left[\Phi(u, \varphi_i) + \sum_{i \in \mathcal{I}} \int_0^1 \varphi_i(x) d\eta_i(x) \right]. \end{aligned}$$

The first equality is obvious, while the second one follows from the minimax principle. Here we use compactness of $\mathcal{U}_{\text{Lip},1}$, linearity in φ_i , concavity in u (follows from convexity of φ_i), the upper semicontinuity was established in Theorem 2.

Step 3. We construct such a family of functions (φ_i) that the infimum in the right-hand side of (48) is prospectively reached on them. Consider a sequence of tuples of functions $\{(\varphi_i^{(n)})_{i \in \mathcal{I}}\}_n \subset \mathcal{U}_{\eta_i, +\infty}$ such that

$$\begin{aligned} \lim_{n \rightarrow \infty} \max_{u \in \mathcal{U}_{\text{Lip},1}} \left[\Phi(u, \varphi_i^{(n)}) + \sum_{i \in \mathcal{I}} \int_0^1 \varphi_i^{(n)}(x) \, d\eta_i(x) \right] = \\ = \max_{u \in \mathcal{U}^1, \nu_i \preceq \eta_i} \int (\langle x, \nabla u \rangle - u(x)) \, d\mu. \end{aligned} \quad (49)$$

Denote by M the optimal value of the objective function

$$M := \max_{u \in \mathcal{U}^1, \nu_i \preceq \eta_i} \int (\langle x, \nabla u \rangle - u(x)) \, d\mu.$$

We may assume that for all n we have

$$2M \geq \max_{u \in \mathcal{U}_{\text{Lip},1}} \left[\Phi(u, \varphi_i^{(n)}) + \sum_{i \in \mathcal{I}} \int_0^1 \varphi_i^{(n)}(x) \, d\eta_i(x) \right].$$

Since $\max_{u \in \mathcal{U}_{\text{Lip},1}} \Phi(u, \varphi_i^{(n)}) \geq 0$ and $\varphi_i^{(n)}(x) \geq 0$ for all x , we conclude that $\int_0^1 \varphi_i^{(n)}(x) \, d\eta_i \leq 2M$ for all $i \in \mathcal{I}$ and for all n . All the functions $\varphi_i^{(n)}$ are non-negative and non-decreasing on $[0, 1]$; therefore, for every $t \in [0, 1)$,

$$\int_0^1 \varphi_i^{(n)}(x) \, d\eta_i \geq \int_t^1 \varphi_i^{(n)}(x) \, d\eta_i \geq \varphi_i^{(n)}(t) \cdot \eta_i([t, 1]).$$

Thus $\varphi_i^{(n)}(t) \leq M_t$ for all $t \in [0, 1)$ and for all n , where $M_t := 2M/\eta_i([t, 1])$.

For every $t \in [0, 1)$, the sequence $\{\varphi_i^{(n)}\}$ is uniformly bounded on $[0, t]$ by the constant M_t . So, applying Helly's principle and passing to subsequences countably many number of times, we can assume that there exists a tuple of functions $(\varphi_i)_{i \in \mathcal{I}}$ defined on $[0, 1)$ such that $\varphi_i^{(n)} \rightarrow \varphi_i$ pointwise on $[0, 1)$.

Each of the functions φ_i is non-negative, non-decreasing, and convex. In particular, $\lim_{t \rightarrow 1-} \varphi_i(t)$ is well defined. We extend the definition of φ_i on $[0, +\infty)$ as follows: define $\varphi_i(1)$ as $\lim_{t \rightarrow 1-} \varphi_i(t) \in \mathbb{R} \cup \{+\infty\}$, and define $\varphi_i(x)$ at every $x \in (1, +\infty)$ to be equal to $+\infty$. The constructed function is lower-semicontinuous. Besides, $\lim_{n \rightarrow \infty} \varphi_i^{(n)}(x) = \varphi_i(x)$ for all $x \in [0, 1)$, and one can easily check that $\liminf_{n \rightarrow \infty} \varphi_i^{(n)}(1) \geq \varphi_i(1)$.

We only need to check that $\varphi_i \in L^1(\eta_i)$ to prove that $\varphi_i \in \mathcal{U}_{\eta_i, +\infty}$. For every i , the sequence of integrals $\{\int_0^1 \varphi_i^{(n)}(x) d\eta_i\}_n$ is bounded. Passing to subsequences, we may additionally assume (and we will use it in the following part of the proof) that each of this sequences converges. By the Fatou lemma,

$$2M \geq \lim_{n \rightarrow \infty} \int_0^1 \varphi_i^{(n)}(x) d\eta_i \geq \int_0^1 \liminf_{n \rightarrow \infty} \varphi_i^{(n)}(x) d\eta_i \geq \int_0^1 \varphi_i(x) d\eta_i. \quad (50)$$

Hence, $\varphi_i \in L^1(\eta_i)$. Thus $\varphi_i \in \mathcal{U}_{\eta_i, +\infty}$.

Step 4. We claim that for all $u \in \mathcal{U}_{\text{Lip}, 1}$ the following inequality holds:

$$\Phi(u, \varphi_i) + \sum_{i \in \mathcal{I}} \int_0^1 \varphi_i(x) d\eta_i(x) \leq M.$$

Fix a function $u \in \mathcal{U}_{\text{Lip}, 1}$. For any $\varepsilon \in (0, 1)$, consider a function $u^\varepsilon = (1 - \varepsilon) \cdot u \in \mathcal{U}_{\text{Lip}, 1 - \varepsilon}$. For all $x \in X$ and for all $i \in \mathcal{I}$, the value $u_{x_i}^\varepsilon(x)$ is not greater than $1 - \varepsilon$. So, for any $i \in \mathcal{I}$, the sequence of functions $\{\varphi_i^{(n)}(u_{x_i}^\varepsilon(x))\}_n$ converges to $\varphi_i(u_{x_i}^\varepsilon(x))$ pointwise almost everywhere. In addition, the inequality $0 \leq \varphi_i^{(n)}(u_{x_i}^\varepsilon(x)) \leq M_{1 - \varepsilon}$ holds for almost all $x \in X$ and for all n ; therefore, it follows from Lebesgue's dominated convergence theorem that

$$\lim_{n \rightarrow \infty} \int \varphi_i^{(n)}(u_{x_i}^\varepsilon) d\mu = \int \varphi_i(u_{x_i}^\varepsilon) d\mu.$$

Combining this with the fact that $\lim_{n \rightarrow \infty} \int_0^1 \varphi_i^{(n)}(x) d\eta_i \geq \int_0^1 \varphi_i(x) d\eta_i$, we conclude that

$$\lim_{n \rightarrow \infty} \left[\Phi(u^\varepsilon, \varphi_i^{(n)}) + \sum_{i \in \mathcal{I}} \int_0^1 \varphi_i^{(n)}(x) d\eta_i(x) \right] \geq \Phi(u^\varepsilon, \varphi_i) + \sum_{i \in \mathcal{I}} \int_0^1 \varphi_i(x) d\eta_i(x).$$

In particular,

$$\begin{aligned} M &= \lim_{n \rightarrow \infty} \max_{v \in \mathcal{U}_{\text{Lip}, 1}} \left[\Phi(v, \varphi_i^{(n)}) + \sum_{i \in \mathcal{I}} \int_0^1 \varphi_i^{(n)}(x) d\eta_i(x) \right] \geq \\ &\geq \lim_{n \rightarrow \infty} \left[\Phi(u^\varepsilon, \varphi_i^{(n)}) + \sum_{i \in \mathcal{I}} \int_0^1 \varphi_i^{(n)}(x) d\eta_i(x) \right] \geq \\ &\geq \Phi(u^\varepsilon, \varphi_i) + \sum_{i \in \mathcal{I}} \int_0^1 \varphi_i(x) d\eta_i(x). \end{aligned}$$

Let $\varepsilon_n = \frac{1}{n}$. For every $i \in \mathcal{I}$, the sequence $\{\varphi_i(u_{x_i}^{\varepsilon_n}(x))\}_n$ is an increasing sequence of non-negative functions that converges to $\varphi_i(u_{x_i}(x))$ pointwise. So, by the Beppo Levi's lemma (Theorem 11.18 in [Aliprantis and Border \(2006\)](#)) we have

$$\lim_{n \rightarrow \infty} \int \varphi_i(u_{x_i}^{\varepsilon_n}(x)) \, d\mu(x) = \int \varphi_i(u_{x_i}(x)) \, d\mu(x).$$

Thus, for all $u \in \mathcal{U}_{\text{Lip},1}$, we have

$$\Phi(u, \varphi_i) + \sum_{i \in \mathcal{I}} \int_0^1 \varphi_i(x) \, d\eta_i(x) = \lim_{n \rightarrow \infty} \left(\Phi(u^{\varepsilon_n}, \varphi_i) + \sum_{i \in \mathcal{I}} \int_0^1 \varphi_i(x) \, d\eta_i(x) \right) \leq M.$$

Since the last inequality holds for all $u \in \mathcal{U}_{\text{Lip},1}$, we conclude that

$$\max_{u \in \mathcal{U}_{\text{Lip},1}} \left[\Phi(u, \varphi_i) + \sum_{i \in \mathcal{I}} \int_0^1 \varphi_i(x) \, d\eta_i(x) \right] \leq M.$$

Thus it follows from the definition of M that the equality holds and the minimum in (48) is reached on a family of functions sequence of functions $(\varphi_i)_{i \in \mathcal{I}}$. \square

C.3 Tools to approach complete duality: main a priori estimate and its corollaries

In the next section, we discuss the complete duality results. The main insight in their proofs allowing us to cope with non-compactness of the problem is an a priori bound on a solution. This bound has a clear economic interpretation in the context of the monopolist's problem with production (46). Here we discuss this bound and its implications.

Informally, the bound is as follows. It states that in the optimal mechanism, the monopolist never gets a negative revenue ex-post, i.e., $\langle \nabla u^{\text{opt}}(x), x \rangle - u(x) - \sum_{i \in \mathcal{I}} \varphi_i(u_{x_i}^{\text{opt}}(x))$ is non-negative.¹² This observation is not elementary as one could possibly expect that by serving those costumers who bring negative profit, the monopolist could extract higher rent from the rest of the population.

We will rely on the notation for functional classes introduced in Section C.2.

Proposition 2. (Main a priori estimate). *Let $(\varphi_i)_{i \in \mathcal{I}}$ be a collection of functions from $\mathcal{U}_{\mathbb{R}_+}$. Then for every function $u \in \mathcal{U}^1$, there exists a non-decreasing convex function \tilde{u} with $\tilde{u}(0) = 0$ such that*

$$\langle x, \nabla \tilde{u}(x) \rangle - \tilde{u}(x) - \sum_{i \in \mathcal{I}} \varphi_i(\tilde{u}_{x_i}(x)) \geq \max \left\{ \langle x, \nabla u(x) \rangle - u(x) - \sum_{i \in \mathcal{I}} \varphi_i(u_{x_i}(x)), 0 \right\} \quad (51)$$

¹²Formulated in terms of monopolist's mechanism $(P^{\text{opt}}, T^{\text{opt}})$, this inequality means $T^{\text{opt}}(x) - \sum_{i \in \mathcal{I}} \varphi_i(P_i^{\text{opt}}(x)) \geq 0$.

for all $x \in X$. In particular, this implies (see Proposition 3) that for any function $u^{\text{opt}} \in \mathcal{U}^1$ maximizing the functional $\Phi(u, \varphi_i)$ over $u \in \mathcal{U}^1$, the inequality

$$\langle x, \nabla u^{\text{opt}}(x) \rangle - u^{\text{opt}}(x) - \sum_{i \in \mathcal{I}} \varphi_i(u_{x_i}^{\text{opt}}(x)) \geq 0 \quad (52)$$

holds almost everywhere.

The main a priori estimate is used in Proposition 3 to show that, for a wide class of functions φ_i , the functional $\Phi(u, \varphi_i)$ attains its maximum on a Lipschitz function u . This fact will be used in approximation Lemma 8 which, together with Proposition 5, help us to justify the minimax principle (Proposition 5) needed to prove complete duality (Theorem 5).

Proof of Proposition 2. Consider the Legendre transform of u

$$u^*(y) = \sup_x (\langle x, y \rangle - u(x)),$$

assuming $u(x) = +\infty$ if $x \notin X$. Next we define

$$v(y) = \max \left\{ u^*(y), \sum_{i \in \mathcal{I}} \varphi_i(y_i) \right\}.$$

Note that v is a lower semicontinuous convex function and $v(0) = 0$. Set

$$\tilde{u} = v^* = \left[\max \left\{ u^*, \sum_{i \in \mathcal{I}} \varphi_i \right\} \right]^*.$$

Then, by the Fenchel–Moreau theorem (Rockafellar, 2015),

$$(\tilde{u})^* = \max \left\{ u^*, \sum_{i \in \mathcal{I}} \varphi_i \right\}$$

and, for every point x where $\nabla \tilde{u}(x)$ exists, one has

$$\langle x, \nabla \tilde{u}(x) \rangle - \tilde{u}(x) - \sum_{i \in \mathcal{I}} \varphi_i(\tilde{u}_{x_i}) = \left((\tilde{u})^* - \sum_{i \in \mathcal{I}} \varphi_i \right) (\nabla \tilde{u}) \geq 0. \quad (53)$$

Consider a point x , such that $\nabla u(x)$ exists and satisfies $\langle x, \nabla u(x) \rangle - u(x) - \sum_{i \in \mathcal{I}} \varphi_i(u_{x_i}(x)) \geq 0$. Equivalently, $(u^* - \sum_{i \in \mathcal{I}} \varphi_i)(\nabla u(x)) \geq 0$. It follows from the theorem about the subdifferential of a maximum of convex functions (Dubovitsky-Milyutin theorem; see Theorem 3.50 in Beck (2017)) that $\partial v(y) =$

$\partial [\max \{u^*, \sum_{i \in \mathcal{I}} \varphi_i\}] (y)$ contains $\partial u^*(y)$ if $u^* \geq \sum_{i \in \mathcal{I}} \varphi_i$. Hence, if x satisfies $(u^* - \sum_{i \in \mathcal{I}} \varphi_i)(\nabla u(x)) \geq 0$, then

$$v(\nabla u(x)) = u^*(\nabla u(x)) \quad \text{and} \quad x \in \partial v(\nabla u(x)).$$

This implies $v(\nabla u(x)) + v^*(x) = \langle \nabla u(x), x \rangle$, hence

$$\tilde{u}(x) = v^*(x) = \langle \nabla u(x), x \rangle - u^*(\nabla u(x)) = u(x)$$

and from the inclusion $x \in \partial v(\nabla u(x))$ we get $\nabla u(x) \in \partial v^*(x) = \partial \tilde{u}(x)$. In particular, if $\nabla \tilde{u}$ exists, then $\nabla u(x) = \nabla \tilde{u}(x)$ and

$$\langle x, \nabla \tilde{u}(x) \rangle - \tilde{u}(x) - \sum_{i \in \mathcal{I}} \varphi_i(\tilde{u}_{x_i}) = \langle x, \nabla u(x) \rangle - u(x) - \sum_{i \in \mathcal{I}} \varphi_i(u_{x_i}). \quad (54)$$

The desired inequality (51) follows from (53) and (54). \square

With the help of the main a priori estimate, we obtain the following a priori bound on the regularity of the optimum.

Proposition 3. *Fix a collection of functions $\varphi_i \in \mathcal{U}_{\mathbb{R}_+}$, $i \in \mathcal{I}$, and consider numbers M_i such that*

$$\varphi_i(M_i) > M_i,$$

which exist by the definition of the class $\mathcal{U}_{\mathbb{R}_+}$. Then there exists a number L depending on M_i and $\varphi_i(M_i)$ such that $\Phi(u, \varphi_i)$ attains its maximum on \mathcal{U}^1 at a function u^{opt} that belongs to $\mathcal{U}_{\text{Lip}, L}$.

Proof. By Proposition 2, any function $v \in \mathcal{U}^1$ can be replaced with a convex non-decreasing function u with $u(0) = 0$ such that

$$\Phi(u, \varphi_i) \geq \Phi(v, \varphi_i) \quad (55)$$

and

$$\langle x, \nabla u(x) \rangle - u(x) - \sum_{i \in \mathcal{I}} \varphi_i(u_{x_i}(x)) \geq 0 \quad (56)$$

for almost all x . Moreover, if the function v does not satisfy inequality (56), then (55) is strict. So, it is enough to check the existence of L depending on φ_i , such that every Lipschitz function u satisfying this inequality is L -Lipschitz.

Indeed, since $u(x) \geq 0$, $u_{x_i}(x) \geq 0$ and $x_i \leq 1$ for all $i \in \mathcal{I}$, assumption (56) implies:

$$\begin{aligned} u_{x_1} + u_{x_2} + \cdots + u_{x_I} - \sum_{i \in \mathcal{I}} \varphi_i(u_{x_i}) &\geq \\ &\geq \langle x, \nabla u(x) \rangle - u(x) - \sum_{i \in \mathcal{I}} \varphi_i(u_{x_i}) \geq 0. \end{aligned} \quad (57)$$

For all $i \in \mathcal{I}$, consider function $\psi_i(x_i) = x_i - \varphi_i(x_i)$. This function is concave and $0 = \psi_i(0) > \psi_i(M_i)$. Hence, ψ_i is decreasing on $[M_i, +\infty)$ and its maximum is reached on $[0, M_i]$. Note that $\psi_i(x_i) \leq x_i$ for all x_i , hence

$$\max_{x \geq 0} \psi_i(x_i) = \max_{0 \leq x_i \leq M_i} \psi_i(x_i) \leq M_i.$$

Inequality (57) can be rewritten in the following form:

$$\sum_{i \in \mathcal{I}} \psi_i(u_{x_i}(x_i)) \geq 0$$

for almost all $x \in X$. Hence, for all $i \in \mathcal{I}$ and almost all $x \in X$,

$$\psi_i(u_{x_i}(x_i)) \geq - \sum_{j \neq i} \psi_j(u_{x_j}(x_j)) \geq - \sum_{j \neq i} M_j. \quad (58)$$

Concavity of ψ_i implies that all $x_i \geq M_i$ satisfy inequality

$$\frac{\psi_i(x_i) - \psi_i(M_i)}{x_i - M_i} \leq \psi'_i(M_i) \leq \frac{\psi_i(M_i) - \psi_i(0)}{M_i - 0} \quad \Leftrightarrow \quad \psi_i(x_i) \leq \frac{x}{M_i} \psi_i(M_i).$$

Hence, if

$$x_i > \max \left\{ M_i, - \frac{\sum_{j \neq i} M_j \cdot M_i}{\psi_i(M_i)} \right\} = \widehat{M}_i,$$

then $\psi_i(x_i) < - \sum_{i \in \mathcal{I}} M_i$.

Hence, inequality (58) implies that $u_{x_i}(x) \leq \widehat{M}_i$ for almost all x . Thus u is L -Lipschitz with $L = \max \{ \widehat{M}_i \}$.

It remains to show that $\Phi(u, \varphi_i)$ attains its maximum on $\mathcal{U}_{\text{Lip}, L}$.

According to estimate (52) we can restrict ourselves to the set of functions $u \in \mathcal{U}^1$ satisfying

$$\langle x, \nabla u(x) \rangle - u(x) - \sum_{i \in \mathcal{I}} \varphi_i(u_{x_i}(x)) \geq 0$$

for almost all $x \in X$. We showed that all such functions u belong to $\mathcal{U}_{\text{Lip}, L}$. This set is compact in uniform convergence topology and $\Phi(u, \varphi_i)$ upper semicontinuous on $\mathcal{U}_{\text{Lip}, L}$. Hence, it reaches its maximum on this set. \square

To prove complete duality, we will need the following weak form of partial duality. The goal is to represent the value in the inf max form so that we can apply the minimax theorem and obtain the maxinf representation, which is done in Proposition 5. The subtlety is that, to apply the minimax theorem, compactness of one

of the spaces is required and so we need to choose carefully a dense minimization subspace \mathcal{Q} in the set of convex one-dimensional functions. In what follows,

$$\mathcal{Q} \subset \mathcal{U}_{\mathbb{R}_+} \quad (59)$$

denotes the set of all increasing, convex functions $\varphi: [0, +\infty) \rightarrow \mathbb{R}_+$ that equal zero at the origin and have bounded derivatives.

Proposition 4. *Under the assumptions of Theorem 3, the following partial duality equation holds:*

$$\max_{u \in \mathcal{U}_{\text{Lip},1}, \nu_i \preceq \eta_i} \int (\langle x, \nabla u \rangle - u(x)) d\mu = \inf_{\varphi_i \in \mathcal{Q}} \max_{u \in \mathcal{U}^1} \left[\Phi(u, \varphi_i) + \sum_{i \in \mathcal{I}} \int_0^1 \varphi_i(x) d\eta_i(x) \right].$$

The key part of the proof is the following lemma.

Lemma 8. *For any family of functions $(\varphi_i)_{i \in \mathcal{I}} \subset \mathcal{U}_{[0,1]}^{+\infty}$ there exist increasing sequences of functions $\{\varphi_i^{(n)}\}_n \subset \mathcal{Q}$, $i \in \mathcal{I}$, such that each sequence $\{\varphi_i^{(n)}\}_n$ converges to φ_i pointwise on $[0, 1]$ and that*

$$\max_{u \in \mathcal{U}_{\text{Lip},1}} \Phi(u, \varphi_i) = \lim_{n \rightarrow \infty} \max_{u \in \mathcal{U}^1} \Phi(u, \varphi_i^{(n)}).$$

Proof of Lemma 8. For every n , denote by $t_{n,i}$ such a point on the interval $[0, 1]$ that $n \in \partial\varphi_i(t_{n,i})$. Such a point exists since $(-\infty, +\infty) = \cup_{t \in [0,1]} \partial\varphi_i(t)$. Denote by $\varphi_i^{(n)}$ the following function:

$$\varphi_i^{(n)}(t) = \begin{cases} \varphi_i(t) & \text{if } t \in [0, t_{n,i}], \\ \varphi_i(t_{n,i}) + n(t - t_{n,i}) & \text{otherwise.} \end{cases}$$

The function $\varphi_i^{(n)}$ is convex; therefore, $\varphi_i^{(n)} \in \mathcal{Q}$. Besides, $\varphi_i^{(n)}(x) \leq \varphi_i(x)$ for $x \in [0, 1]$, and $\varphi_i^{(n)}$ coincides with φ_i on the interval $[0, t_n]$. For each i , the sequence of points $\{t_{n,i}\}_n$ is monotonically increasing and converges to 1; therefore, each sequence $\{\varphi_i^{(n)}\}_n$ is increasing and converges to φ_i pointwise on $[0, 1]$. Finally, since pointwise supremum of lower semicontinuous functions is lower semicontinuous, we conclude that

$$\lim_{n \rightarrow \infty} \varphi_i^{(n)}(1) = \lim_{t \rightarrow 1} \lim_{n \rightarrow \infty} \varphi_i^{(n)}(t) = \lim_{t \rightarrow 1} \varphi_i(t) = \varphi_i(1).$$

Consider any function $v \in \mathcal{U}_{\text{Lip},1}$. For each $i \in \mathcal{I}$, the sequence of non-negative functions $\{\varphi_i^{(n)}(v_{x_i}(x))\}_n$ is monotonically increasing and converges to $\varphi_i(v_{x_i}(x))$ pointwise almost everywhere; therefore, by the monotone convergence theorem

$$\lim_{n \rightarrow \infty} \int \varphi_i^{(n)}(v_{x_i}) d\mu = \int \varphi_i(v_{x_i}) d\mu.$$

So,

$$\Phi(v, \varphi_i) = \lim_{n \rightarrow \infty} \Phi(v, \varphi_i^{(n)}) \leq \lim_{n \rightarrow \infty} \max_{u \in \mathcal{U}^1} \Phi(u, \varphi_i^{(n)}).$$

Since the last inequality holds for all $v \in \mathcal{U}_{\text{Lip},1}$, we conclude that

$$\max_{u \in \mathcal{U}_{\text{Lip},1}} \Phi(u, \varphi_i) \leq \lim_{n \rightarrow \infty} \max_{u \in \mathcal{U}^1} \Phi(u, \varphi_i^{(n)}). \quad (60)$$

Let $u^{(n)}$ be a maximizer of the functional $\Phi(\cdot, \varphi_i^{(n)})$. For every $n \geq 4I$, we have $\varphi_i^{(n)}(1 + \frac{4I}{n}) \geq 4I$ and $1 + \frac{4I}{n} \leq 2$. So, if we denote by M_i the number $1 + \frac{4I}{n}$, by Proposition 3, we have

$$u_{x_i} \leq \max \left\{ M_i, \frac{(I-1) \cdot M_i^2}{\varphi_i(M_i) - M_i} \right\} \leq \max \left\{ 1 + \frac{4I}{n}, \frac{4(I-1)}{4I-2} \right\} = 1 + \frac{4I}{n}.$$

So, $u^{(n)} \in \mathcal{U}_{\text{Lip}, 1+4I/n} \subset \mathcal{U}_{\text{Lip},2}$ for all $n \geq 4I$.

Passing to a subsequence, one can assume that $u^{(n)} \rightarrow u^{(0)}$ uniformly. Since $u_{x_i}^{(n)}(x) \rightarrow u_{x_i}^{(0)}(x)$ pointwise for almost all x , and $u_{x_i}^{(n)} \leq 1 + \frac{4I}{n} \rightarrow 1$, we conclude that $\bar{u}^{(0)} \in \mathcal{U}_{\text{Lip},1}$. By the dominated convergence theorem,

$$\lim_{n \rightarrow \infty} \int [\langle x, \nabla u^{(n)}(x) \rangle - u^{(n)}(x)] d\mu \rightarrow \int [\langle x, \nabla u^{(0)}(x) \rangle - u^{(0)}(x)] d\mu. \quad (61)$$

We claim that for almost all $x \in X$ and for every $i \in \mathcal{I}$ we have

$$\liminf_{n \rightarrow \infty} \varphi_i^{(n)}(u_{x_i}^{(n)}(x)) \geq \varphi_i(u_{x_i}^{(0)}(x)).$$

Let $t_i = \sup\{t: \varphi_i(t) \leq 2I+1\}$. First, since every $\varphi_i^{(n)}$ is a non-decreasing function,

$$\varphi_i^{(n)}(u_{x_i}^{(n)}(x)) \geq \varphi_i^{(n)}(\min(u_{x_i}^{(n)}(x), t_i)).$$

Next, let us check that

$$\liminf_{n \rightarrow \infty} \varphi_i^{(n)}(\min(u_{x_i}^{(n)}(x), t_i)) \geq \liminf_{n \rightarrow \infty} \varphi_i(\min(u_{x_i}^{(n)}(x), t_i)).$$

Indeed, if $t_i < 1$, then $t_{n,i} \geq t_i$ for all large enough n . Therefore,

$$\varphi_i^{(n)}(\min(u_{x_i}^{(n)}(x), t_i)) = \varphi_i(\min(u_{x_i}^{(n)}(x), t_i))$$

for all large enough n . Otherwise, suppose that $t_i = 1$. Then $\varphi_i(1) \leq 2I+1 < +\infty$; therefore, by the lower semicontinuity of φ_i for any $\varepsilon > 0$ there exists a point $p_i < 1$ such that $\varphi_i(p_i) \geq \varphi_i(1) - \varepsilon$. Then for all n such that $t_{n,i} \geq p_i$ the

inequality $\varphi_i^{(n)}(x) - \varphi_i(x) \geq -\varepsilon$ holds for all $x \in [0, 1]$. Indeed, if $x \leq t_{n,i}$, then $\varphi_i^{(n)}(x) = \varphi_i(x)$. Otherwise, $x \geq t_{n,i} \geq p_i$; therefore,

$$\varphi_i^{(n)}(x) \geq \varphi_i^{(n)}(p_i) = \varphi_i(p_i) \geq \varphi_i(x) - \varepsilon.$$

Thus, in the case $t_i = 1$, the inequality

$$\liminf_{n \rightarrow \infty} \varphi_i^{(n)}(\min(u_{x_i}^{(n)}(x), t_i)) \geq \liminf_{n \rightarrow \infty} \varphi_i(\min(u_{x_i}^{(n)}(x), t_i)) - \varepsilon$$

holds for all $\varepsilon > 0$. Letting ε tend to 0, we obtain the desired one.

Finally, we check that $u_{x_i}^{(0)}(x) = \lim_{n \rightarrow \infty} u_{x_i}^{(n)}(x) \leq t_i$. If $t_i = 1$, the inequality holds since $u^{(0)} \in \mathcal{U}_{\text{Lip},1}$. Suppose that $t_i < 1$ and $u_{x_i}^{(0)}(x) > t_i$. Then for all large enough n we have $u_{x_i}^{(n)}(x) \geq t_i$ and $t_{n,i} \geq t_i$. In this case,

$$\varphi_i^{(n)}(u_{x_i}^{(n)}(x)) \geq \varphi_i^{(n)}(t_i) = \varphi_i(t_i) = 2I + 1.$$

On the other hand, by Proposition 2, the inequality

$$\varphi_i^{(n)}(u_{x_i}^{(n)}(x)) \leq \langle x, \nabla u^{(n)}(x) \rangle$$

holds for almost all $x \in X$. For all $n \geq 4I$ we have $u^{(n)} \in \mathcal{U}_{\text{Lip},2}$; therefore, for almost all x we have

$$\varphi_i^{(n)}(u_{x_i}^{(n)}(x)) \leq 2I,$$

which contradicts the previous inequality.

Thus

$$\begin{aligned} \liminf_{n \rightarrow \infty} \varphi_i^{(n)}(u_{x_i}^{(n)}(x)) &\geq \liminf_{n \rightarrow \infty} \varphi_i^{(n)}(\min(u_{x_i}^{(n)}(x), t_i)) \geq \\ &\geq \liminf_{n \rightarrow \infty} \varphi_i(\min(u_{x_i}^{(n)}(x), t_i)) = \\ &= \varphi_i\left(\liminf_{n \rightarrow \infty} \min(u_{x_i}^{(n)}(x), t_i)\right) = \varphi_i(u_{x_i}^{(0)}(x)). \end{aligned}$$

Therefore, it follows from Fatou's lemma that

$$\liminf_{n \rightarrow \infty} \int \varphi_i^{(n)}(u_{x_i}^{(n)}(x)) \, d\mu \geq \int \liminf_{n \rightarrow \infty} \varphi_i^{(n)}(u_{x_i}^{(n)}(x)) \, d\mu \geq \int \varphi_i(u_{x_i}^{(0)}(x)) \, d\mu.$$

So, combining it with (61) we conclude that

$$\Phi(u^{(0)}, \varphi_i) \geq \limsup_{n \rightarrow \infty} \Phi(u^{(n)}, \varphi_i^{(n)}) = \lim_{n \rightarrow \infty} \max_{u \in \mathcal{U}^1} \Phi(u, \varphi_i^{(n)}).$$

Comparing it to (60), we conclude that the equality holds and this completes the proof of the statement. \square

Proof of Proposition 4. By the standard argument,

$$\max_{u \in \mathcal{U}_{\text{Lip},1}, \nu_i \preceq \eta_i} \int (\langle x, \nabla u \rangle - u(x)) d\mu \leq \inf_{\varphi_i \in \mathcal{Q}} \max_{u \in \mathcal{U}^1} \left[\Phi(u, \varphi_i) + \sum_{i \in \mathcal{I}} \int_0^1 \varphi_i(x) d\eta_i(x) \right].$$

Let $(\varphi_i)_{i \in \mathcal{I}} \subset \mathcal{U}_{\eta_i, +\infty}$ be family of functions on which the minimum in the right-hand side of (48) is reached. By Lemma 8, there exist increasing sequences $\{\varphi_i^{(n)}\}_n \subset \mathcal{Q}$, $1 \leq i \leq I$, such that the sequence $\{\varphi_i^{(n)}\}_n$ converges to φ_i pointwise on $[0, 1]$ and that

$$\max_{u \in \mathcal{U}_{\text{Lip},1}} \Phi(u, \varphi_i) = \lim_{n \rightarrow \infty} \max_{u \in \mathcal{U}^1} \Phi(u, \varphi_i^{(n)}).$$

By the monotone convergence theorem, we additionally have $\lim_{n \rightarrow \infty} \int \varphi_i^{(n)} d\eta_i = \int \varphi_i d\eta_i$. Thus

$$\begin{aligned} \max_{u \in \mathcal{U}_{\text{Lip},1}, \nu_i \preceq \eta_i} \int (\langle x, \nabla u \rangle - u(x)) d\mu &= \\ &= \max_{u \in \mathcal{U}_{\text{Lip},1}} \left[\Phi(u, \varphi_i) + \sum_{i \in \mathcal{I}} \int_0^1 \varphi_i(x) d\eta_i(x) \right] \\ &= \lim_{n \rightarrow \infty} \max_{u \in \mathcal{U}_{\text{Lip},1}} \left[\Phi(u, \varphi_i^{(n)}) + \sum_{i \in \mathcal{I}} \int_0^1 \varphi_i^{(n)}(x) d\eta_i(x) \right]. \end{aligned}$$

□

C.4 Complete duality

Relying on the partial duality established in Sections C.1 and the a priori estimate from Section C.3, we are ready to prove complete duality for the monopolist's problem with general majorization (44) extending Theorems 1 and 2.

We will rely on notation introduced in Section C.2. Denote by \mathcal{C} the set of smooth nonnegative (coordinate-wise) vector fields $c = (c_i)_{i \in \mathcal{I}}$ such that

$$\int (\langle x, \nabla u(x) \rangle - u(x)) d\mu \leq \int \langle c(x), \nabla u(x) \rangle d\mu \quad (62)$$

for all $u \in \mathcal{U}^1$. If $d\mu(x) = \rho(x) dx$ where ρ is continuously differentiable, the condition above is equivalent to the majorization constraint $-\text{div}_\rho[c] \succeq m$, where m is the transform measure. Note, in particular, that $x \in \mathcal{C}$.

Theorem 5. Let η_i be probability measures on $[0, 1]$ such that $\eta_i([t, 1]) > 0$ for any $t > 0$. Then the value¹³ of the Rochet-Choné problem with general majorization (44) is equal to

$$\inf_{\varphi_i \in \mathcal{Q}, c \in \mathcal{C}} \left(\sum_{i \in \mathcal{I}} \int \varphi_i^*(c_i) d\mu + \sum_{i \in \mathcal{I}} \int_0^1 \varphi_i(x) d\eta_i(x) \right). \quad (63)$$

Let us check that Theorem 1 is a corollary of Theorem 5.

Proof of Theorem 1. Fix all η_i to coincide with the distribution of ξ^{B-1} where ξ is uniform on $[0, 1]$. Then the value of the Rochet-Choné problem with general majorization is $\frac{1}{B}$ fraction of the value of the corresponding auctioneer's problem with B bidders (Proposition 1).

In Section 4, we already demonstrated weak duality for the auctioneer's problem, i.e., we checked that $\frac{1}{B}$ fraction of the auctioneer's revenue cannot exceed

$$\inf_{\substack{(\varphi_i)_{i \in \mathcal{I}}, \\ \pi \succeq m}} \left[\text{Beck}_\rho(\pi, \Phi) + \sum_{i \in \mathcal{I}} \int_0^1 \varphi_i(z^{B-1}) dz \right], \quad (64)$$

where $\Phi(c) = \sum_i \varphi_i^*(|c_i|)$. Hence, to prove that this expression coincides with $\frac{1}{B}$ fraction of the optimal revenue, it is enough to demonstrate that it is upper-bounded by (63).

Comparing (62) to the definition of divergence (15) and that of the transform measure (18), we see that \mathcal{C} consists of vector fields c with non-negative components such that $-\text{div}_\rho[c] \succeq m$ or, equivalently, there exists $\pi \succeq m$ such that $\text{div}_\rho[c] + \pi = 0$. Let \mathcal{C}_\pm be the superset of \mathcal{C} obtained by dropping the non-negativity condition. We get

$$\inf_{c \in \mathcal{C}_\pm} \sum_{i \in \mathcal{I}} \int \varphi_i^*(|c_i|) d\mu = \text{Beck}_\rho(\pi, \Phi).$$

Since $\mathcal{C} \subset \mathcal{C}_\pm$, we conclude that

$$\inf_{c \in \mathcal{C}} \sum_{i \in \mathcal{I}} \int \varphi_i^*(c_i) d\mu = \inf_{c \in \mathcal{C}} \sum_{i \in \mathcal{I}} \int \varphi_i^*(|c_i|) d\mu \geq \inf_{\substack{(\varphi_i)_{i \in \mathcal{I}}, \\ \pi \succeq m}} \text{Beck}_\rho(\pi, \Phi).$$

Hence, (63) is an upper bound on (64). Thus (64) is equal to $\frac{1}{B}$ fraction of the auctioneer's optimal revenue. \square

¹³Recall that this value is defined by $\max_{u \in \mathcal{U}_{\text{Lip}, 1}, \nu_i \preceq \eta_i} \int (\langle x, \nabla u \rangle - u(x)) d\mu$.

As a preliminary step to proving Theorem 5, we prove a complete duality result for the monopolist's problem with fixed production costs. Denote by $\bar{\mathcal{C}}$ the set of bounded nonnegative vector fields $c = (c_i)_{i \in \mathcal{I}}$, not necessary smooth, such that

$$\int (\langle x, \nabla u(x) \rangle - u(x)) d\mu \leq \int \langle c(x), \nabla u(x) \rangle d\mu$$

for all $u \in \mathcal{U}^1$. Note that $\mathcal{C} \subset \bar{\mathcal{C}}$. Recall that \mathcal{Q} is defined in (59).

Proposition 5. *For any family of functions $(\varphi_i)_{i \in \mathcal{I}} \subset \mathcal{Q}$, the following relation holds*

$$\max_{u \in \mathcal{U}^1} \Phi(u, \varphi_i) = \min_{c \in \bar{\mathcal{C}}} \int \sum_{i \in \mathcal{I}} \varphi_i^*(c_i) d\mu.$$

Moreover, if all the functions φ_i are continuously differentiable, then the vector field $c_i = \varphi'_i(\bar{u}_{x_i})$ solves the dual problem, where \bar{u} is an optimal solution to the problem $\max_{u \in \mathcal{U}^1} \Phi(u, \varphi_i)$.

Proof. For every $\varphi_i \in \mathcal{Q}$, one has

$$\langle x, \nabla u \rangle - u(x) - \sum_{i \in \mathcal{I}} \varphi_i(u_{x_i}) = \min_{c(x) \geq 0} \left(\langle x - c(x), \nabla u \rangle - u(x) + \sum_{i \in \mathcal{I}} \varphi_i^*(c_i) \right).$$

The minimum is taken among of all nonnegative vector fields and it is attained at $c_i = \varphi'_i(u_{x_i})$. In particular, $0 \leq c_i \leq \sup \varphi'_i$ (we apply here that $\varphi_i \in \mathcal{Q}$, hence the derivatives φ'_i are uniformly bounded).

Thus for every $\varphi_i \in \mathcal{Q}$ we get

$$\langle x, \nabla u \rangle - u(x) - \sum_{i \in \mathcal{I}} \varphi_i(u_{x_i}) = \min_{B(\varphi_i)} \left(\langle x - c(x), \nabla u \rangle - u(x) + \sum_{i \in \mathcal{I}} \varphi_i^*(c_i) \right),$$

where $B(\varphi_i)$ is the set of non-negative vector fields $c = (c_i)$ satisfying $c_i(x) \leq \sup \varphi'_i(x)$ for μ -a.e. x . Hence,

$$\max_{u \in \mathcal{U}^1} \Phi(u, \varphi_i) = \sup_{u \in \mathcal{U}^1} \min_{B(\varphi_i)} \left(\int (\langle x - c(x), \nabla u \rangle - u(x) + \sum_{i \in \mathcal{I}} \varphi_i^*(c_i)) d\mu \right),$$

We apply the minimax principle and the fact that $B(\varphi_i)$ is a closed subset of a ball in $L^\infty(\mu)$, endowed with the *-weak topology. The Banach-Alaoglu theorem implies that $B(\varphi_i)$ is compact. Hence,

$$\begin{aligned} & \sup_{u \in \mathcal{U}^1} \min_{c \in B(\varphi_i)} \left(\int (\langle x - c(x), \nabla u \rangle - u(x) + \sum_{i \in \mathcal{I}} \varphi_i^*(c_i)) d\mu \right) = \\ & = \min_{c \in B(\varphi_i)} \sup_{u \in \mathcal{U}^1} \left(\int (\langle x - c(x), \nabla u \rangle - u(x) + \sum_{i \in \mathcal{I}} \varphi_i^*(c_i)) d\mu \right) \end{aligned}$$

Let us check that the minimax principle is applicable. Indeed, the convexity of the functional on $B(\varphi_i)$ is obvious, it is sufficient to check the lower semicontinuity. Let us consider a sequence $c^{(n)} \in B(\varphi_i)$ such that $c^{(n)} \rightarrow c$ $*$ -weakly in $L^\infty(\mu)$ (in particular, weakly in $L^2(\mu)$). It is sufficient to show that $\liminf_n \int \varphi_i^*(c_i^{(n)}) d\mu \geq \int \varphi_i^*(c_i) d\mu$.

Passing to a subsequence (if necessary), which we denote again by $c_i^{(n)}$, one can assume without loss of generality that $\int \varphi_i^*(c_i^{(n)}) d\mu$ has a limit and the sequence of $\frac{1}{N} \sum_{n=1}^N c_i^{(n)}$ converges in $L^2(\mu)$ and μ -a.e. Applying convexity of φ_i , one gets

$$\begin{aligned} \lim_n \int \varphi_i^*(c_i^{(n)}) d\mu &= \lim_N \frac{1}{N} \sum_{n=1}^N \int \varphi_i^*(c_i^{(n)}) d\mu \geq \\ &\geq \lim_N \int \varphi_i^* \left(\frac{1}{N} \sum_{n=1}^N c_i^{(n)} \right) d\mu \geq \int \varphi_i^*(c_i) d\mu. \end{aligned}$$

In the last inequality we use convergence almost everywhere and the Fatou lemma.

The next step is obvious:

$$\begin{aligned} \min_{c \in B(\varphi_i)} \sup_{u \in \mathcal{U}^1} \left(\int (\langle x - c(x), \nabla u \rangle - u(x) + \sum_{i \in \mathcal{I}} \varphi_i^*(c_i)) d\mu \right) = \\ = \min_{c \in B(\varphi_i) \cap \bar{\mathcal{C}}} \sum_{i \in \mathcal{I}} \int \varphi_i^*(c_i) d\mu. \end{aligned}$$

Hence,

$$\max_{u \in \mathcal{U}^1} \Phi(u, \varphi_i) = \min_{c \in B(\varphi_i) \cap \bar{\mathcal{C}}} \left(\sum_{i \in \mathcal{I}} \int \varphi_i^*(c_i) d\mu \right).$$

Clearly, $\min_{c \in B(\varphi_i) \cap \bar{\mathcal{C}}} \sum_{i \in \mathcal{I}} \int \varphi_i^*(c_i) d\mu$ can be replaced with $\min_{c \in \bar{\mathcal{C}}} \sum_{i \in \mathcal{I}} \int \varphi_i^*(c_i) d\mu$, since, by the standard arguments, $\max_{u \in \mathcal{U}^1} \Phi(u, \varphi_i) \leq \min_{c \in \bar{\mathcal{C}}} \sum_{i \in \mathcal{I}} \int \varphi_i^*(c_i) d\mu$.

Now, assume that all the functions φ_i are continuously differentiable. Let $c^{(0)}$ be an optimal solution to the dual problem $\min_{c \in \bar{\mathcal{C}}} \int \sum_{i \in \mathcal{I}} \varphi_i^*(c_i) d\mu$, and let \bar{u} be an optimal solution to the problem $\max_{u \in \mathcal{U}^1} \Phi(u, \varphi_i)$. The following sequence of inequalities holds:

$$\begin{aligned} &\int \left(\langle \nabla \bar{u}(x), x \rangle - \bar{u}(x) - \sum_{i \in \mathcal{I}} \varphi_i(\bar{u}_{x_i}(x)) \right) d\mu \leq \\ &\leq \int \left(\langle \nabla \bar{u}(x), x \rangle - \bar{u}(x) - \sum_{i \in \mathcal{I}} \bar{u}_{x_i}(x) \cdot c_i^{(0)}(x) \right) d\mu + \int \sum_{i \in \mathcal{I}} \varphi_i^*(c_i^{(0)}) d\mu \\ &\leq \int \sum_{i \in \mathcal{I}} \varphi_i^*(c_i^{(0)}) d\mu. \end{aligned}$$

The left-hand and right-hand sides of this inequality are equal. Therefore,

$$\int \left(\langle \nabla \bar{u}(x), x \rangle - \bar{u}(x) - \sum_{i \in I} \bar{u}_{x_i}(x) \cdot c_i^{(0)}(x) \right) d\mu = 0$$

and

$$\sum_{i \in I} \int \left(\varphi_i(\bar{u}_{x_i}) + \varphi_i^*(c_i^{(0)}) - \bar{u}_{x_i} \cdot c_i^{(0)} \right) d\mu = 0.$$

Thus $c_i^{(0)}(x) \in \partial\varphi_i(\bar{u}_{x_i}(x))$ for μ -almost all x . Since the functions φ_i are continuously differentiable, we conclude that $\partial\varphi_i(x_0) = \{\varphi'_i(x_0)\}$ for all $x_0 > 0$ and $\partial\varphi_i(x_0) = (-\infty, \varphi'_i(x_0))$ at the point $x_0 = 0$. Thus $c_i^{(0)}(x) = \varphi'_i(\bar{u}_{x_i}(x))$ for μ -almost all such points x that $\bar{u}_{x_i}(x) > 0$.

Consider a vector field defined by the equation $\bar{c}_i(x) = \varphi'_i(\bar{u}_{x_i}(x))$ for all x . Since $\bar{c}_i(x) \geq c_i^{(0)}(x)$ for μ -almost all x , we conclude easily that $\bar{c} \in \bar{\mathcal{C}}$. Since $\bar{c}_i(x) \in \partial\varphi_i(\bar{u}_{x_i}(x))$ for μ -almost all x ,

$$\sum_{i \in I} \int \left(\varphi_i(\bar{u}_{x_i}) + \varphi_i^*(\bar{c}_i) - \bar{u}_{x_i} \cdot \bar{c}_i \right) d\mu = 0.$$

In addition, $\bar{u}_{x_i} \cdot c_i^{(0)} = \bar{u}_{x_i} \cdot \bar{c}_i$ for μ -almost all x ; therefore,

$$\begin{aligned} \int \left(\langle \nabla \bar{u}(x), x \rangle - \bar{u}(x) - \sum_{i \in I} \bar{u}_{x_i}(x) \cdot \bar{c}_i(x) \right) d\mu &= \\ &= \int \left(\langle \nabla \bar{u}(x), x \rangle - \bar{u}(x) - \sum_{i \in I} \bar{u}_{x_i}(x) \cdot c_i^{(0)}(x) \right) d\mu = 0. \end{aligned}$$

Thus we finally conclude that

$$\begin{aligned} \int \sum_{i \in I} \varphi_i^*(\bar{c}_i) d\mu &= \int \left(\langle \nabla \bar{u}(x), x \rangle - \bar{u}(x) - \sum_{i \in I} \varphi_i(\bar{u}_{x_i}) \right) d\mu + \\ &+ \sum_{i \in I} \int \left(\varphi_i(\bar{u}_{x_i}) + \varphi_i^*(\bar{c}_i) - \bar{u}_{x_i} \cdot \bar{c}_i \right) d\mu - \\ &- \int \left(\langle \nabla \bar{u}(x), x \rangle - \bar{u}(x) - \sum_{i \in I} \bar{u}_{x_i}(x) \cdot \bar{c}_i(x) \right) d\mu = \\ &= \int \left(\langle \nabla \bar{u}(x), x \rangle - \bar{u}(x) - \sum_{i \in I} \varphi_i(\bar{u}_{x_i}) \right) d\mu. \end{aligned}$$

This equality means that $\bar{c} \in \arg \min_{c \in \bar{\mathcal{C}}} \int \sum_{i \in I} \varphi_i^*(c_i) d\mu$. □

Next, we extend the previous result to smooth vector fields.

Proposition 6. *For any family of strictly convex continuously differentiable functions $(\varphi_i)_{i \in \mathcal{I}} \subset \mathcal{Q}$, the following relation holds*

$$\max_{u \in \mathcal{U}^1} \Phi(u, \varphi_i) = \inf_{c \in \mathcal{C}} \int \sum_{i \in \mathcal{I}} \varphi_i^*(c_i) d\mu.$$

Proof. Let $\bar{u} \in \arg \max_{u \in \mathcal{U}^1} \Phi(u, \varphi_i)$, and let $(\bar{c}_i)_{i \in \mathcal{I}}$ be the vector field defined by the formula $\bar{c}_i(x) = \varphi_i'(\bar{u}_{x_i}(x))$. By Proposition 5, the vector field \bar{c} is an optimal solution to the dual problem

$$\max_{c \in \bar{\mathcal{C}}} \int \sum_{i \in \mathcal{I}} \varphi_i^*(c_i) d\mu.$$

We claim that the function \bar{u}_{x_i} is continuous at all the points x_0 such that \bar{u} is differentiable at x_0 . Indeed, consider any sequence of points $\{x_n\}_n$ converging to x_0 . One can verify easily that if $p_n \in \partial \bar{u}(x_n)$ and if $\partial \bar{u}(x_0)$ contains only one point $\nabla \bar{u}(x_0)$, then the sequence $\{p_n\}_n$ converges to $\nabla \bar{u}(x_0)$. Hence, the sequence $\{\bar{u}_{x_i}(x_n)\}_n$ converges to $\bar{u}_{x_i}(x_0)$, and this implies the continuity of \bar{u}_{x_i} at the point x_0 .

In particular, this means that the function $\bar{c}_i = \varphi_i'(\bar{u}_{x_i})$ is continuous almost everywhere; therefore, the function

$$\bar{c}_i^{\text{up}}(x_0) = \limsup_{x \rightarrow x_0} \bar{c}_i(x)$$

is upper semi-continuous and $\bar{c}_i = \bar{c}_i^{\text{up}}$ almost everywhere. Thus the vector field $\bar{c}^{\text{up}} = (\bar{c}_i^{\text{up}})_{i \in \mathcal{I}}$ belongs to $\bar{\mathcal{C}}$ and is an optimal solution to the dual problem.

Denote $L_i = \sup_x \varphi_i'(\bar{u}_{x_i}(x)) = \sup_x \bar{c}_i^{\text{up}}(x)$. Since the function \bar{u}_{x_i} is bounded by Proposition 3 and the function φ_i' is strictly increasing, we have $L_i < \sup_x \varphi_i'(x) = \widehat{L}_i$. Since the function \bar{c}_i^{up} is upper semi-continuous, it can be written as the pointwise limit of a non-increasing family $\{c_i^{(n)}\}$ of smooth functions; moreover, we can require that $\sup_x c_i^{(n)} \leq (L_i + \widehat{L}_i)/2$ for all n .

Since $c_i^{(n)} \geq \bar{c}_i^{\text{up}}$, we clearly have $c^{(n)} = (c_i^{(n)})_{i \in \mathcal{I}} \in \mathcal{C}$. Let us check that

$$\lim_{n \rightarrow \infty} \int \sum_{i \in \mathcal{I}} \varphi_i^*(c_i^{(n)}) d\mu = \int \sum_{i \in \mathcal{I}} \varphi_i^*(\bar{c}_i^{\text{up}}) d\mu. \quad (65)$$

Indeed, the function φ_i^* is continuous and non-decreasing on the interval $[0, (L_i + \widehat{L}_i)/2]$; therefore, the sequence of functions $\varphi_i^*(c_i^{(n)})$ is a non-increasing family that converges to $\varphi_i^*(\bar{c}_i^{\text{up}})$ pointwise. Then Beppo Levi's lemma implies (65).

This implies that

$$\inf_{c \in \mathcal{C}} \int \sum_{i \in \mathcal{I}} \varphi_i^*(c_i) d\mu \leq \lim_{n \rightarrow \infty} \int \sum_{i \in \mathcal{I}} \varphi_i^*(c_i^{(n)}) d\mu = \int \sum_{i \in \mathcal{I}} \varphi_i^*(\bar{c}_i^{\text{up}}) d\mu = \max_{u \in \mathcal{U}^1} \Phi(u, \varphi_i).$$

On the other hand, since $\mathcal{C} \subset \bar{\mathcal{C}}$, we have

$$\inf_{c \in \mathcal{C}} \int \sum_{i \in \mathcal{I}} \varphi_i^*(c_i) d\mu \geq \min_{c \in \bar{\mathcal{C}}} \int \sum_{i \in \mathcal{I}} \varphi_i^*(c_i) d\mu = \max_{u \in \mathcal{U}^1} \Phi(u, \varphi_i).$$

This implies the desired duality relation. \square

Now, we can prove Theorem 5.

Proof of Theorem 5. Denote by \mathcal{Q}_{st} the subset of functions $\varphi \in \mathcal{Q}$ such that φ is continuously differentiable and strictly convex. We claim that

$$\begin{aligned} \inf_{\varphi_i \in \mathcal{Q}} \max_{u \in \mathcal{U}^1} \left[\Phi(u, \varphi_i) + \sum_{i \in \mathcal{I}} \int_0^1 \varphi_i(x) d\eta_i(x) \right] &= \\ &= \inf_{\varphi_i \in \mathcal{Q}_{\text{st}}} \max_{u \in \mathcal{U}^1} \left[\Phi(u, \varphi_i) + \sum_{i \in \mathcal{I}} \int_0^1 \varphi_i(x) d\eta_i(x) \right]. \end{aligned} \quad (66)$$

Since $\mathcal{Q}_{\text{st}} \subset \mathcal{Q}$, we conclude that the left-hand side is not greater than the right-hand side. Let us check the opposite inequality.

Consider any function $\varphi \in \mathcal{Q}$. Let $\{p^{(n)}(x)\}$ be a sequence of smooth non-negative kernel functions such that each function $p^{(n)}(x)$ is supported on the interval $[-1/n, 0]$ and $\int_{-\infty}^{+\infty} p^{(n)}(x) dx = 1$. Consider the function

$$\varphi^{(n)}(x) = (\varphi * p^{(n)})(x) = \int_0^{+\infty} \varphi(t) p^{(n)}(x - t) dt = \int_0^{\frac{1}{n}} \varphi(x + t) p^{(n)}(-t) dt.$$

One can easily check that the function $\varphi^{(n)}(x)$ is smooth, convex, non-negative, and non-decreasing for $x \geq 0$. Moreover, if $\varphi'(x) \leq L$ for all x , then $(\varphi^{(n)})'(x) \leq L$ for all x , so $\varphi^{(n)}(x) - \varphi^{(n)}(0) \in \mathcal{Q}$. Finally, $\varphi(x) \leq \varphi^{(n)}(x) \leq \varphi(x + \frac{1}{n}) \leq \varphi(x) + \frac{L}{n}$ for all $x \geq 0$. Denoting

$$\widehat{\varphi}^{(n)}(x) = (\varphi^{(n)}(x) - \varphi^{(n)}(0)) + \frac{1}{n} (x + \exp(-x) - 1),$$

we conclude that $\widehat{\varphi}^{(n)} \in \mathcal{Q}_{\text{st}}$, that $\widehat{\varphi}^{(n)}(x) \geq \varphi(x) - \varphi^{(n)}(0) \geq \varphi(x) - \frac{L}{n}$ for all $x \geq 0$, and that

$$\widehat{\varphi}^{(n)}(x) - \varphi(x) \leq \frac{L + 1}{n} \quad \text{for all } x \in [0, 1].$$

Consider any family of functions $(\varphi_i)_{i \in I} \subset \mathcal{Q}$. For each $i \in I$, let $\{\varphi_i^{(n)}\}_n \subset \mathcal{Q}_{\text{st}}$ be a sequence of functions such that $\varphi_i^{(n)}(x) - \varphi_i(x) \geq -\frac{1}{n}$ for all $x \geq 0$ and that $\varphi_i^{(n)}(x) - \varphi_i(x) \leq \frac{1}{n}$ for all $x \in [0, 1]$. Since $\varphi_i^{(n)}(x) \geq \varphi_i(x) - \frac{1}{n}$ for all $x \geq 0$, we have

$$\Phi(u, \varphi_i^{(n)}) \leq \Phi(u, \varphi_i) + \frac{I}{n} \quad \Rightarrow \quad \max_{u \in \mathcal{U}^1} \Phi(u, \varphi_i^{(n)}) \leq \max_{u \in \mathcal{U}^1} \Phi(u, \varphi_i) + \frac{I}{n}.$$

In addition,

$$\int_0^1 \varphi_i^{(n)}(x) d\eta_i(x) \leq \int_0^1 \varphi_i(x) d\eta_i(x) + \frac{1}{n}.$$

Thus

$$\begin{aligned} \max_{u \in \mathcal{U}^1} \left[\Phi(u, \varphi_i^{(n)}) + \sum_{i \in \mathcal{I}} \int_0^1 \varphi_i^{(n)}(x) d\eta_i(x) \right] &\leq \\ &\leq \max_{u \in \mathcal{U}^1} \left[\Phi(u, \varphi_i) + \sum_{i \in \mathcal{I}} \int_0^1 \varphi_i(x) d\eta_i(x) \right] + \frac{2I}{n}. \end{aligned}$$

In particular, for any family of functions $(\varphi_i)_{i \in I} \subset \mathcal{Q}$ we have

$$\begin{aligned} \inf_{\varphi_i \in \mathcal{Q}_{\text{st}}} \max_{u \in \mathcal{U}^1} \left[\Phi(u, \varphi_i) + \sum_{i \in \mathcal{I}} \int_0^1 \varphi_i(x) d\eta_i(x) \right] &\leq \\ &\leq \liminf_{n \rightarrow \infty} \max_{u \in \mathcal{U}^1} \left[\Phi(u, \varphi_i^{(n)}) + \sum_{i \in \mathcal{I}} \int_0^1 \varphi_i^{(n)}(x) d\eta_i(x) \right] \\ &\leq \max_{u \in \mathcal{U}^1} \left[\Phi(u, \varphi_i) + \sum_{i \in \mathcal{I}} \int_0^1 \varphi_i(x) d\eta_i(x) \right]. \end{aligned}$$

This implies the equality (66).

Now, the desired result follows from the combination of Propositions 4 and 6. \square

Next, we prove the duality theorem in the strong form $\max = \min$. To do it, we need an extension of the set of feasible vector fields \mathcal{C} . We denote by \mathcal{C}^{mes} the set of tuples of non-negative measures $(\varsigma_i)_{i \in \mathcal{I}}$ satisfying

$$\int (\langle \nabla u(x), x \rangle - u(x)) d\mu \leq \sum_{i \in \mathcal{I}} \int u_{x_i} d\varsigma_i$$

for every smooth $u \in \mathcal{U}^1$.

Theorem 6. *Under the assumptions of Theorem 5, the following identity holds:*

$$\begin{aligned} \max_{u \in \mathcal{U}^1, \nu_i \leq \eta_i} \int (\langle x, \nabla u \rangle - u(x)) d\mu &= \\ &= \min_{\substack{\varsigma \in \mathcal{C}^{mes}, \\ \varphi_i \in \mathcal{U}_{\eta_i, +\infty}}} \sum_{i \in \mathcal{I}} \left(\varsigma_i^{\text{sing}}(X) + \int_0^1 \varphi_i(x) d\eta_i + \int \varphi_i^*(\varsigma_i^a(x)) d\mu \right). \end{aligned}$$

Note that Theorem 2 is a particular case of Theorem 6 for all η_i equal to the distribution of ξ^{B-1} with ξ uniform on $[0, 1]$.

To prove Theorem 6 we need several auxiliary results.

Lemma 9. *The set \mathcal{C}^{mes} is closed in the weak*-topology.*

Proof. Trivial, as we can additionally require the test function u to be smooth. \square

Lemma 10. *Let $(\varsigma_i)_{i \in \mathcal{I}} \in \mathcal{C}^{mes}$, and let $\varsigma_i = \varsigma_i^a(x) d\mu(x) + \varsigma_i^{\text{sing}}$ be a decomposition of the component ς_i into an absolutely continuous and a singular part w.r.t. μ . Then for any $u \in \mathcal{U}_{\text{Lip}, 1}$, the following inequality holds:*

$$\int (\langle \nabla u(x), x \rangle - u(x)) d\mu \leq \sum_{i \in \mathcal{I}} \left(\varsigma_i^{\text{sing}}(X) + \int u_{x_i}(x) \cdot \varsigma_i^a(x) d\mu \right).$$

Proof. Let \bar{u} be any non-negative convex function defined on the whole $\mathbb{R}^{\mathcal{I}}$ such that $0 \leq \bar{u}_{x_i}(x) \leq 1$ for all $x \in \mathbb{R}^{\mathcal{I}}$ and that $\bar{u}|_X = u$. It can be defined, for instance, in the following way: $\bar{u}(x) = \sup_{\alpha} l_{\alpha}$, where $\{l_{\alpha}\}$ is the set of affine functions satisfying $l_{\alpha}|_X \leq u$.

Let $\{p_n\}$ be a sequence of Gaussian kernels converging to $\delta(0)$, and denote by $\bar{u}^{(n)}$ the convolution $\bar{u} * p_n$. One can easily check that $u^{(n)}(x)$ is a smooth non-negative convex function such that $0 \leq \bar{u}_{x_i}^{(n)}(x) \leq 1$ for all $x \in \mathbb{R}^{\mathcal{I}}$. Moreover, the sequence $\{\bar{u}^{(n)}\}$ converges uniformly to u on X . Thus, denoting by $u^{(n)}(x)$ the function $\bar{u}^{(n)}(x) - \bar{u}^{(n)}(0)$, we conclude that $u^{(n)}$ is smooth, $u^{(n)} \in \mathcal{U}_{\text{Lip}, 1}$, and the sequence $\{u^{(n)}\}$ converges uniformly to u on X .

Since $u_{x_i}^{(n)}(x) \leq 1$ for all $x \in X$, the following inequality holds:

$$\begin{aligned} \int (\langle \nabla u^{(n)}(x), x \rangle - u^{(n)}(x)) d\mu &\leq \sum_{i \in \mathcal{I}} \left(\int u_{x_i}^{(n)}(x) \cdot \varsigma_i^a(x) d\mu + \int u_{x_i}^{(n)}(x) d\varsigma_i^{\text{sing}} \right) \\ &\leq \sum_{i \in \mathcal{I}} \left(\int u_{x_i}^{(n)}(x) \cdot \varsigma_i^a(x) d\mu + \varsigma_i^{\text{sing}}(X) \right). \end{aligned} \tag{67}$$

Since $\{u^{(n)}\}$ converges to u uniformly on X , the sequence $\{\nabla u^{(n)}(x)\}$ converges to $\nabla u(x)$ for μ -almost all x . Therefore, by the Lebesgue's dominated convergence theorem

$$\begin{aligned}\lim_{n \rightarrow \infty} \int (\langle \nabla u^{(n)}(x), x \rangle - u^{(n)}(x)) d\mu &= \int (\langle \nabla u(x), x \rangle - u(x)) d\mu, \\ \lim_{n \rightarrow \infty} \int u_{x_i}^{(n)}(x) \cdot \varsigma_i^a(x) d\mu &= \int u_{x_i}(x) \cdot \varsigma_i^a(x) d\mu.\end{aligned}$$

Thus, passing to the limits in (67), we obtain the desired inequality. \square

The following proposition extends the complete duality result for the monopolist's problem with fixed costs (Proposition 5) so that the minimum in the dual is attained.

Proposition 7. *For any given family of functions $(\varphi_i)_{i \in \mathcal{I}} \subset \mathcal{U}_{[0,1]}^{+\infty}$ and an absolutely continuous measure μ on X , the following duality relation holds:*

$$\max_{u \in \mathcal{U}_{\text{Lip},1}} \Phi(u, \varphi_i) = \min_{\varsigma \in \mathcal{C}^{mes}} \sum_{i \in \mathcal{I}} \left(\varsigma_i^{\text{sing}}(X) + \int \varphi_i^*(\varsigma_i^a(x)) d\mu \right),$$

where $\varsigma_i = \varsigma_i^a(x) d\mu + \varsigma_i^{\text{sing}}$ is a decomposition of the component ς_i into an absolutely continuous and a singular part w.r.t. μ .

Proof. By Lemma 10, for any $u \in \mathcal{U}_{\text{Lip},1}$ and $\varsigma \in \mathcal{C}^{mes}$ we have

$$\int (\langle \nabla u(x), x \rangle - u(x)) d\mu \leq \sum_{i \in \mathcal{I}} \left(\varsigma_i^{\text{sing}}(X) + \int u_{x_i}(x) \cdot \varsigma_i^a(x) d\mu \right).$$

Therefore,

$$\begin{aligned}\Phi(u, \varphi_i) &\leq \sum_{i \in \mathcal{I}} \left(\varsigma_i^{\text{sing}}(X) + \int [u_{x_i}(x) \cdot \varsigma_i^a(x) - \varphi_i(u_{x_i}(x))] d\mu \right) \leq \\ &\leq \sum_{i \in \mathcal{I}} \left(\varsigma_i^{\text{sing}}(X) + \int \varphi_i^*(\varsigma_i^a(x)) d\mu \right),\end{aligned}$$

where the last part of the inequality follows from the inequality $\varphi_i(u_{x_i}(x)) + \varphi_i^*(\varsigma_i^a(x)) \geq u_{x_i}(x) \cdot \varsigma_i^a(x)$, which holds for all x . Thus we conclude that

$$\max_{u \in \mathcal{U}_{\text{Lip},1}} \Phi(u, \varphi_i) \leq \min_{\varsigma \in \mathcal{C}^{mes}} \sum_{i \in \mathcal{I}} \left(\varsigma_i^{\text{sing}}(X) + \int \varphi_i^*(\varsigma_i^a(x)) d\mu \right). \quad (68)$$

By Lemma 8, there exist increasing sequences $\{\varphi_i^{(n)}\}_n \subset \mathcal{Q}$, $1 \leq i \leq I$ that converge to φ_i pointwise on $[0, 1]$ and that

$$\max_{u \in \mathcal{U}_{\text{Lip},1}} \Phi(u, \varphi_i) = \lim_{n \rightarrow \infty} \max_{u \in \mathcal{U}^1} \Phi(u, \varphi_i^{(n)}). \quad (69)$$

Denote by M the maximal value of $\Phi(u, \varphi_i)$. We may assume that for all n we have

$$2M \geq \max_{u \in \mathcal{U}^1} \Phi(u, \varphi_i^{(n)}).$$

By Proposition 5, for each n there exists a tuple of functions $\{c_i^{(n)}\}_{i \in \mathcal{I}} \subset \overline{\mathcal{C}}$ such that

$$\max_{u \in \mathcal{U}^1} \Phi(u, \varphi_i^{(n)}) = \sum_{i \in \mathcal{I}} \int \left(\varphi_i^{(n)} \right)^* (c_i^{(n)}(x)) d\mu. \quad (70)$$

Denote by $\varsigma_i^{(n)}$ the measure $c_i^{(n)}(x) d\mu$. By the definition of $(\varphi_i^{(n)})^*$, for all $x \geq 0$ and for every $t \in [0, 1]$ we have

$$(\varphi_i^{(n)})^*(x) \geq t \cdot x - \varphi_i^{(n)}(t) \geq t \cdot x - \varphi_i(t); \quad (71)$$

in the last inequality, we use that $\{\varphi_i^{(n)}\}_n$ is an increasing sequence of functions. So, for each i and for every $t \in [0, 1]$ the following inequality holds:

$$\begin{aligned} 2M &\geq \sum_{i \in \mathcal{I}} \int \left(\varphi_i^{(n)} \right)^* (c_i^{(n)}(x)) d\mu \geq \\ &\geq \int \left(\varphi_i^{(n)} \right)^* (c_i^{(n)}(x)) d\mu \geq t \int c_i^{(n)}(x) d\mu - \varphi_i(t). \end{aligned}$$

This means that the sequence $\varsigma_i^{(n)}(X) = \int c_i^{(n)}(x) d\mu$ is bounded from above by $(2M + \varphi_i(t))/t$. Applying the Prokhorov theorem and passing to a subsequence, we may assume that the sequence of measures $\{\varsigma_i^{(n)}\}_n$ converges weakly to some non-negative measure ς_i . Also, applying the Komlos theorem and passing to a subsequence, we may assume that

$$\frac{1}{n} \sum_{i=1}^n c_i^{(n)} \xrightarrow{n \rightarrow \infty} c_i$$

for some $c_i \in L^1(\mu)$ almost everywhere.

Since $\varsigma_i^{(n)}$ converges weakly to ς_i , one has

$$\lim_{n \rightarrow \infty} \int c_i^{(n)}(x) d\mu = \lim_{n \rightarrow \infty} \varsigma_i^{(n)}(X) = \varsigma_i(X).$$

So, combining equations (69) and (70), we conclude that for every $t \in [0, 1)$ the following equality holds:

$$\begin{aligned} \max_{u \in \mathcal{H}_{\text{Lip}, 1}} \Phi(u, \varphi_i) &= \lim_{n \rightarrow \infty} \sum_{i \in \mathcal{I}} \int (\varphi_i^{(n)})^* (c_i^{(n)}(x)) d\mu = \\ &= \sum_{i \in \mathcal{I}} t \cdot \varsigma_i(X) + \lim_{n \rightarrow \infty} \sum_{i \in \mathcal{I}} \int \left[(\varphi_i^{(n)})^* (c_i^{(n)}(x)) - t \cdot c_i^{(n)}(x) \right] d\mu. \end{aligned} \quad (72)$$

Consider the last item of the previous expression's right-hand side. By Cesaro means,

$$\begin{aligned} \lim_{n \rightarrow \infty} \sum_{i \in \mathcal{I}} \int \left[(\varphi_i^{(n)})^* (c_i^{(n)}(x)) - t \cdot c_i^{(n)}(x) \right] d\mu \\ = \lim_{n \rightarrow \infty} \int \sum_{i \in \mathcal{I}} \frac{1}{n} \sum_{k=1}^n \left[(\varphi_i^{(k)})^* (c_i^{(k)}(x)) - t \cdot c_i^{(k)}(x) \right] d\mu \\ \geq \sum_{i \in \mathcal{I}} \liminf_{n \rightarrow \infty} \left(\int \frac{1}{n} \sum_{k=1}^n \left[(\varphi_i^{(k)})^* (c_i^{(k)}(x)) - t \cdot c_i^{(k)}(x) \right] d\mu \right). \end{aligned} \quad (73)$$

Denote by $\psi_i^{(n)}(x)$ the function $(\varphi_i^{(n)})^* (c_i^{(n)}(x)) - t \cdot c_i^{(n)}(x)$. Inequality (71) implies that each function $\psi_i^{(n)}(x)$ is bounded from below by $-\varphi_i(t)$, so the function $(\psi_i^{(1)}(x) + \dots + \psi_i^{(n)}(x))/n$ is also bounded from below by $-\varphi_i(t)$. Therefore, it follows from the Fatou lemma that

$$\begin{aligned} \liminf_{n \rightarrow \infty} \left(\int \frac{1}{n} \sum_{k=1}^n \left[(\varphi_i^{(k)})^* (c_i^{(k)}(x)) - t \cdot c_i^{(k)}(x) \right] d\mu \right) &\geq \\ &\geq \int \liminf_{n \rightarrow \infty} \left(\frac{1}{n} \sum_{k=1}^n \left[(\varphi_i^{(k)})^* (c_i^{(k)}(x)) - t \cdot c_i^{(k)}(x) \right] \right) d\mu \\ &= \int \liminf_{n \rightarrow \infty} \left(\frac{1}{n} \sum_{k=1}^n (\varphi_i^{(k)})^* (c_i^{(k)}(x)) \right) d\mu - t \cdot \int c_i(x) d\mu, \end{aligned} \quad (74)$$

where the last equation follows from the fact that $\frac{1}{n} \sum_{i=1}^n c_i^{(k)}$ converges to $c_i(x)$ for μ -almost every x .

Finally, for every $x \geq 0$ and $y \in [0, 1)$ the following inequality holds:

$$(\varphi_i^{(n)})^* (c_i^{(n)}(x)) + \varphi_i^{(n)}(y) \geq y \cdot c_i^{(n)}(x);$$

therefore,

$$\frac{1}{n} \sum_{k=1}^n (\varphi_i^{(k)})^*(c_i^{(k)}(x)) + \frac{1}{n} \sum_{k=1}^n \varphi_i^{(k)}(y) \geq y \cdot \frac{1}{n} \sum_{k=1}^n c_i^{(k)}(x).$$

Passing to the limits, we conclude that

$$\liminf_{n \rightarrow \infty} \left(\frac{1}{n} \sum_{k=1}^n (\varphi_i^{(k)})^*(c_i^{(k)}(x)) \right) + \varphi_i(y) \geq y \cdot c_i(x)$$

for all $y \in [0, 1]$. Thus

$$\liminf_{n \rightarrow \infty} \left(\frac{1}{n} \sum_{k=1}^n (\varphi_i^{(k)})^*(c_i^{(k)}(x)) \right) \geq \max_{y \in [0, 1]} [y \cdot c_i(x) - \varphi_i(y)] = \varphi_i^*(c_i(x)). \quad (75)$$

Combining inequalities (72)–(75), we conclude that for all $t \in [0, 1]$ the following inequality holds:

$$\max_{u \in \mathcal{U}_{\text{Lip}, 1}} \Phi(u, \varphi_i) \geq \sum_{i \in \mathcal{I}} \left(t \cdot \varsigma_i(X) + \int \varphi_i^*(c_i(x)) d\mu - t \cdot \int c_i(x) d\mu \right).$$

Letting t tend to 1, we obtain the following inequality:

$$\max_{u \in \mathcal{U}_{\text{Lip}, 1}} \Phi(u, \varphi_i) \geq \sum_{i \in \mathcal{I}} \left(\varsigma_i(X) + \int \varphi_i^*(c_i(x)) d\mu - \int c_i(x) d\mu \right).$$

Next, we check that $c_i(x) d\mu \leq \varsigma_i$ for all $i \in \mathcal{I}$. By Cesaro means, the sequence of measures $(c_i^{(1)}(x) d\mu + \dots + c_i^{(n)}(x) d\mu)/n$ converges weakly to ς_i . Then by the well-known property of the weak convergence for any closed subset A of X we have

$$\limsup \int_A \frac{1}{n} \sum_{k \in \mathcal{I}} c_i^{(k)}(x) d\mu \leq \varsigma_i(A).$$

By the Fatou lemma,

$$\limsup \int_A \frac{1}{n} \sum_{k \in \mathcal{I}} c_i^{(k)}(x) d\mu \geq \liminf \int_A \frac{1}{n} \sum_{k=1}^{\mathcal{I}} c_i^{(k)}(x) d\mu \geq \int_A c_i(x) d\mu.$$

Thus $\int_A c_i(x) d\mu \leq \varsigma_i(A)$ for all closed subsets A of X ; therefore, $c_i(x) d\mu \leq \varsigma_i$.

Let $\varsigma_i = \varsigma_i^a(x) d\mu + \varsigma_i^{\text{sing}}$ be a decomposition of the component ς_i into an absolutely continuous and a singular part w.r.t. μ . Since $c_i(x) d\mu \leq \varsigma_i$, we conclude that $c_i(x) \leq \varsigma_i^a(x)$ for μ -almost every x . Therefore,

$$\begin{aligned} \varsigma_i(X) + \int \varphi_i^*(c_i(x)) d\mu - \int c_i(x) d\mu &= \varsigma_i^{\text{sing}}(X) + \int (\varphi_i^*(c_i(x)) - c_i(x) + \varsigma_i^a(x)) d\mu \\ &\geq \varsigma_i^{\text{sing}}(X) + \int \varphi_i^*(\varsigma_i^a(x)) d\mu, \end{aligned}$$

where the last inequality follows from the fact that φ_i^* is a 1-Lipschitz function. Finally,

$$\max_{u \in \mathcal{U}_{\text{Lip},1}} \Phi(u, \varphi_i) \geq \sum_{i \in \mathcal{I}} \left(\int_0^1 \varphi_i(x) d\eta_i + \varsigma_i^{\text{sing}}(X) + \int \varphi_i^*(\varsigma_i^a(x)) d\mu \right).$$

Comparing this inequality to (68), we conclude that the equality holds and thus complete the proof of Proposition 7. \square

Proof of Theorem 6. The theorem is a combination of Proposition 7 and Theorem 3. \square

D Examples

We show how vector fields solving the dual problem from Theorem 2 can be constructed explicitly. As we will see, the solutions may be non-unique and singular.

We illustrate this points for the simplest benchmark problem with two items and one bidder whose values are uniformly distributed on $[0, 1]^2$. The solution was first obtained by (Manelli and Vincent, 2006). Daskalakis et al. (2017) demonstrated how it can be constructed via the the associated Monge-Kantorovich transportation problem (34). The optimal function $u = u^{\text{opt}}$ is given by

$$u^{\text{opt}}(x, y) = \begin{cases} 0 & (x, y) \in \mathcal{Z} \\ x - \frac{2}{3} & (x, y) \in \mathcal{A} \\ y - \frac{2}{3} & (x, y) \in \mathcal{B} \\ x + y - \frac{4-\sqrt{2}}{3} & (x, y) \in \mathcal{W} \end{cases}, \quad (76)$$

where the sets $\mathcal{Z}, \mathcal{A}, \mathcal{B}$ and \mathcal{W} are depicted in Figure 4 borrowing the notation from Daskalakis et al. (2017). The answers for the transform measure m defined by (18) and for the optimal “imbalance” $\pi = \pi^{\text{opt}}$ majorizing m are as follows:

$$m = \delta_0 + \lambda_1|_{[0,1] \times \{0\}} + \lambda_1|_{\{0\} \times [0,1]} - 3\lambda_2|_{[0,1]^2}, \quad (77)$$

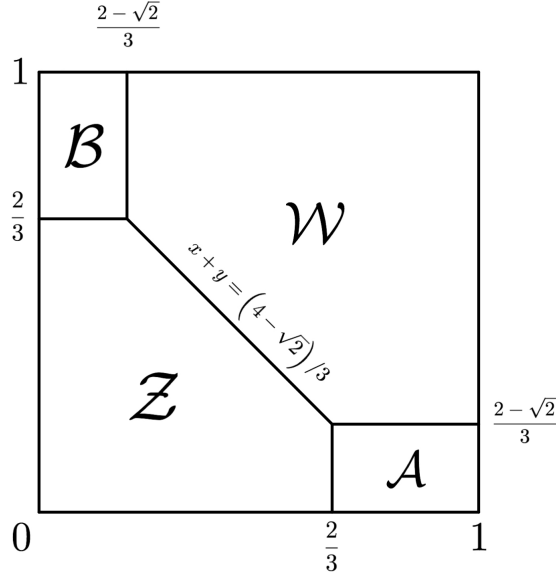


Figure 4: Partition of the square with respect to the optimal u .

$$\pi^{\text{opt}} = \lambda_1|_{[0,1] \times \{0\}} + \lambda_1|_{\{0\} \times [0,1]} - 3\lambda_2|_{[0,1]^2 \setminus \mathcal{Z}}, \quad (78)$$

where λ_2, λ_1 are the two- and one-dimensional Lebesgue measures, respectively.

Recall that for $B = 1$ bidder, the dual problem from Theorem 2 can be simplified (Corollary 3). It takes the following form:

$$\text{minimize: } \int_{[0,1]^2} (d\varsigma_1 + d\varsigma_2) \quad (79)$$

over vector measures $\varsigma = (\varsigma_1, \varsigma_2)$ satisfying

$$\int u \, dm \leq \int (u_x d\varsigma_1 + u_y d\varsigma_2) \quad (80)$$

for all convex non-decreasing u with $u(0) = 0$.

First, let us construct an absolutely continuous solution, i.e., such that $d\varsigma_i = c^i dx dy$. We will need the following lemma.

Lemma 11. *Assume that a couple of nonnegative functions c^1, c^2 satisfy*

1.

$$c^1|_{\mathcal{Z} \cup \mathcal{B}} = 0, \quad c^2|_{\mathcal{Z} \cup \mathcal{A}} = 0, \quad (81)$$

2. c^1 is weakly differentiable along x and satisfies the following integration by parts identity for every smooth φ

$$\int_{[0,1]^2} \varphi_x \cdot c^1 \, dx \, dy = - \int_{[0,1]^2} \varphi \cdot c_x^1 \, dx \, dy + \int_0^1 \varphi(1, y) \, dy.$$

Similarly c^2 is weakly differentiable along y and satisfies the following integration by parts identity for every smooth φ

$$\int_{[0,1]^2} \varphi_y \cdot c^2 \, dx \, dy = - \int_{[0,1]^2} \varphi \cdot c_y^2 \, dx \, dy + \int_0^1 \varphi(x, 1) \, dx.$$

3.

$$c_x^1 + c_y^2 = 3, \tag{82}$$

on $[0, 1]^1 \setminus \mathcal{Z}$ almost everywhere.

Then the vector field $c = (c^1, c^2)$ satisfies

$$\operatorname{div}[c] + \pi^{\text{opt}} = 0 \tag{83}$$

and $\varsigma = (\varsigma_1, \varsigma_2)$ such that $d\varsigma_i = c^i \, dx \, dy$ is a solution to the dual problem (79).

Proof. Take any convex non-decreasing 1-Lipschitz function u with $u(0) = 0$. Then

$$\int (xu_x + yu_y - u) \, dx \, dy = \int u \, dm \leq \int u \, d\pi^{\text{opt}}.$$

The first equality is the definition of the transform measure m , the second one holds since $m \preceq \pi^{\text{opt}}$. Since π^{opt} is given by an explicit formula (78), the identity for the divergence (83) follows from an elementary computation. Using (83) and the definition of divergence, we obtain

$$\begin{aligned} \int u \, d\pi^{\text{opt}} &= - \int u \operatorname{div}[c] \, dx \, dy = \int \langle \nabla u, c \rangle \, dx \, dy \leq \\ &\leq \int (c^1 + c^2) \, dx \, dy = \int (d\varsigma_1 + d\varsigma_2), \end{aligned}$$

where, we used that $0 \leq u_x \leq 1$ and $0 \leq u_y \leq 1$ to get the inequality. Substituting $u = u^{\text{opt}}$ given by (76), we see that the two inequalities become equalities because $u|_{\mathcal{Z}} = 0$ and $c|_{\mathcal{Z}} = 0$. We conclude that the objective in the dual problem on ς coincides with the optimal value of the primal problem $\int u^{\text{opt}} \, dm$. Thus ς is the optimal solution of the dual. \square

Example 1 (Absolutely continuous solution). Consider the following vector field:

$$c(x, y) = (c^1(x, y), c^2(x, y)),$$

where

$$c^2(x, y) = c^1(y, x)$$

and

$$c^1(x, y) = \begin{cases} 0 & (x, y) \in \mathcal{Z} \cup \mathcal{B} \\ 3x - 2 & (x, y) \in \mathcal{A} \\ \frac{3}{2}(x + y - \frac{4-\sqrt{2}}{3}) & (x, y) \in \mathcal{W}, x \leq \frac{2}{3}, y \leq \frac{2}{3} \\ \frac{9}{2}(1-x)(y - \frac{2-\sqrt{2}}{3}) & \\ \quad + 3(x - \frac{2}{3}) & (x, y) \in \mathcal{W}, \frac{2}{3} \leq x \leq 1, \frac{2-\sqrt{2}}{3} \leq y \leq \frac{2}{3} \\ \frac{9}{4}(x - \frac{2-\sqrt{2}}{3})^2 & (x, y) \in \mathcal{W}, \frac{2-\sqrt{2}}{3} \leq x \leq \frac{2}{3}, \frac{2}{3} \leq y \leq 1 \\ \frac{1}{2} + \frac{3}{2}(x - \frac{2}{3}) & (x, y) \in \mathcal{W}, \frac{2}{3} \leq x \leq 1, \frac{2}{3} \leq y \leq 1 \end{cases}$$

One can check that the vector field $c = (c^1, c^2)$ satisfies the assumptions of Lemma 11. We conclude that c solves the dual problem.

Remark 1 (Non-uniqueness). It turns out that there are many solutions to the dual problem. However, the reader should be not confused by existing results on uniqueness of the optimal vector field c in Beckmann's problem; see, e.g., [Santambrogio \(2015\)](#). Unlike most of the works on Beckmann's problem, our cost function is given by the l^1 -norm $\sum_i |c_i|$ instead of the Euclidean l^2 -norm $\sqrt{\sum_i |c_i|^2}$. The l^1 -norm constitutes a degenerate case. Indeed, if c is a solution and φ is a smooth function, supported on a small neighbourhood of a point $(x_0, y_0) \in \text{int}(\mathcal{W})$, where $c^i(x_0, y_0) > 0$, then for sufficiently small ε the vector field

$$c_\varepsilon = (c^1 + \varepsilon\varphi_y, c^2 - \varepsilon\varphi_x)$$

satisfies all the assumptions. Integrating by parts one gets

$$\int \sum_{i=1}^2 c_\varepsilon^i \, dx \, dy = \int \sum_{i=1}^2 c^i \, dx \, dy.$$

Thus c_ε is also a solution.

Moreover, one can easily find solutions which are not weakly differentiable. Let a, δ be numbers and Q be the square with the center (a, a) and vertices

$$q_{-,a} = (a - \delta, a), \quad q_{a,-} = (a, a - \delta), \quad q_{+,a} = (a + \delta, a), \quad q_{a,+} = (a, a + \delta).$$

Define

$$\begin{aligned} \psi_1(x, y) &= I_Q(x, y)(-I_{a-\delta \leq x \leq a}(x) + I_{a \leq x \leq a+\delta}(x)) \\ \psi_2(x, y) &= I_Q(x, y)(-I_{a-\delta \leq y \leq a}(y) + I_{a \leq y \leq a+\delta}(y)) \end{aligned}$$

It is easy to verify that

$$\frac{\partial \psi_1(x, y)}{\partial y} = \frac{\partial \psi_2(x, y)}{\partial x} = \frac{1}{\sqrt{2}} \left[-\lambda_1|_{[q_{-,a}, q_{a,-}]} + \lambda_1|_{[q_{-,a}, q_{a,+}]} - \lambda_1|_{[q_{a,+}, q_{+,a}]} + \lambda_1|_{[q_{a,-}, q_{+,a}]} \right] \quad (84)$$

in the weak sense, where $[a, b]$ denotes the segment joining a and b . Clearly,

$$\int \psi_1 \, dx \, dy = \int \psi_2 \, dx \, dy = 0$$

and (84) implies that

$$\operatorname{div}[(\psi_2, -\psi_1)] = 0.$$

Thus for any solution c to the dual problem, strictly positive in some neighbourhood U of a point $(a, a) \in \operatorname{int}(\mathcal{W})$, the vector field

$$c + (\psi_2, -\psi_1)$$

is a solution to the dual problem for sufficiently small δ .

D.0.1 Singular solutions

It may seem intuitive — at least for our toy example — that vector fields solving the dual problem must be integrable functions. Surprisingly, there exist singular solutions. We construct a measure-valued solution ς with the following properties:

- The vector field ς is singular, i.e., its components are not absolutely continuous measures: namely, ς has an atom at $(1, 1)$.
- Because of this atom, the divergence of ς is not a measure and can only be defined in the space of generalized functions.

Example 2 (Singular solution). Define a couple of measures $(\varsigma_1, \varsigma_2)$ as follows:

$$\begin{aligned} \varsigma_1 &= \mathbb{1}((x, y) \in \mathcal{A}) \cdot 3 \left(x - \frac{2}{3} \right) \, dx \, dy + \\ &\quad + \mathbb{1} \left((x, y) \in \mathcal{W} \cap \left\{ y \leq \frac{2}{3} \right\} \right) \cdot 3 \left(x + y - \frac{4 - \sqrt{2}}{3} \right) \, dx \, dy, \\ \varsigma_2 &= \mathbb{1} \left(y \geq \frac{2}{3} \right) \cdot 3 \left(y - \frac{2}{3} \right) \, dx \, dy + C \cdot \delta(x = 1, y = 1), \end{aligned}$$

where $C = \frac{1}{18} + \frac{\sqrt{2}}{27}$; see Fig. 4.

Let us show that the vector-measure $\varsigma = (\varsigma_1, \varsigma_2)$ is a solution to the dual problem. We need to demonstrate that ς satisfies the majorization condition (80) and minimizes the dual objective (79) over such vector measures.

First, we check that ς satisfies (80). Integrating by parts, we conclude that for any smooth u defined on $[0, 1]^2$,

$$\begin{aligned} \int \frac{\partial u}{\partial x} d\varsigma_1 &= \int_0^{\frac{2-\sqrt{2}}{3}} v(y) dy + \int_{\frac{2-\sqrt{2}}{3}}^{\frac{2}{3}} v(y) \cdot 3 \left(y + \frac{\sqrt{2}-1}{3} \right) dy - \\ &\quad - 3 \int_{(\mathcal{W} \cup \mathcal{A}) \cap \{y \leq \frac{2}{3}\}} u(x, y) dx dy, \\ \int \frac{\partial u}{\partial y} d\varsigma_2 &= \int_0^1 u(x, 1) dx + C \cdot v'(1) - 3 \int_{y \geq \frac{2}{3}} u(x, y) dx dy, \end{aligned} \quad (85)$$

where $v(y) := u(1, y)$.

Let us prove that for any smooth convex non-decreasing u with $u(0) = 0$,

$$\int_0^1 \int_0^1 \left(x \cdot \frac{\partial u}{\partial x} + y \cdot \frac{\partial u}{\partial y} - u \right) dx dy \leq \int \frac{\partial u}{\partial x} d\varsigma_1 + \int \frac{\partial u}{\partial y} d\varsigma_2.$$

Integrating by parts,

$$\begin{aligned} \int_0^1 \int_0^1 \left(x \cdot \frac{\partial u}{\partial x} + y \cdot \frac{\partial u}{\partial y} - u \right) dx dy &= \\ &= \int_0^1 u(x, 1) dx + \int_0^1 v(y) dy - 3 \int_0^1 \int_0^1 u(x, y) dx dy. \end{aligned}$$

Comparing it to (85), we conclude that the inequality above is equivalent to the following one:

$$C \cdot v'(1) - \int_{\frac{2}{3}}^1 v(y) dy + \int_{\frac{2-\sqrt{2}}{3}}^{\frac{2}{3}} v(y) \cdot 3 \left(y - \frac{2-\sqrt{2}}{3} \right) dy \geq -3 \int_{\mathcal{Z}} u(x, y) dx dy.$$

The right-hand side is non-positive, so it is sufficient to prove that the left-hand side is non-negative for any smooth convex non-decreasing u with $u(0) = 0$. For any such u , the function $v(y) = u(1, y)$ is a non-decreasing convex function defined on $[0, 1]$. So, for any $y \in [\frac{2}{3}, 1]$, we have $v(y) \leq v(2/3) + (y - 2/3) \cdot v'(1)$. Therefore,

$$\int_{\frac{2}{3}}^1 v(y) dy \leq \frac{1}{3} \cdot v\left(\frac{2}{3}\right) + \frac{1}{18} \cdot v'(1).$$

In addition, for any $y \in [\frac{2-\sqrt{2}}{3}, \frac{2}{3}]$, we have

$$v\left(\frac{2}{3}\right) - v(y) \leq v'\left(\frac{2}{3}\right) \cdot \left(\frac{2}{3} - y\right) \leq v'(1) \cdot \left(\frac{2}{3} - y\right).$$

Therefore,

$$\begin{aligned}
& \int_{\frac{2-\sqrt{2}}{3}}^{\frac{2}{3}} v(y) \cdot 3 \left(y - \frac{2-\sqrt{2}}{3} \right) dy \geq \\
& \geq v\left(\frac{2}{3}\right) \cdot \int_{\frac{2-\sqrt{2}}{3}}^{\frac{2}{3}} 3 \left(y - \frac{2-\sqrt{2}}{3} \right) dy - v'(1) \cdot \int_{\frac{2-\sqrt{2}}{3}}^{\frac{2}{3}} 3 \left(y - \frac{2-\sqrt{2}}{3} \right) \left(\frac{2}{3} - y \right) dy \\
& = \frac{1}{3} \cdot v\left(\frac{2}{3}\right) - \frac{\sqrt{2}}{27} \cdot v'(1). \quad (86)
\end{aligned}$$

Finally,

$$\begin{aligned}
& C \cdot v'(1) - \int_{\frac{2}{3}}^1 v(y) dy + \int_{\frac{2-\sqrt{2}}{3}}^{\frac{2}{3}} v(y) \cdot 3 \left(y - \frac{2-\sqrt{2}}{3} \right) dy \geq \\
& \geq C \cdot v'(1) - \frac{1}{3} \cdot v\left(\frac{2}{3}\right) - \frac{1}{18} \cdot v'(1) + \frac{1}{3} \cdot v\left(\frac{2}{3}\right) - \frac{\sqrt{2}}{27} \cdot v'(1) = 0. \quad (87)
\end{aligned}$$

We conclude that $(\varsigma_1, \varsigma_2)$ satisfies (80).

Let us verify the optimality of $(\varsigma_1, \varsigma_2)$. It is enough to check that the value of the dual objective (79) on ς coincides with the optimal value of the primal problem. Recall that u^{opt} denotes the optimal function in the primal problem. Hence,

$$\int (d\varsigma_1 + d\varsigma_2) = \int \frac{\partial u^{\text{opt}}}{\partial x} d\varsigma_1 + \int \frac{\partial u^{\text{opt}}}{\partial y} d\varsigma_2.$$

Thus it is enough to check that the right-hand side is equal to the value of the primal problem:

$$\int_0^1 \int_0^1 \left(x \cdot \frac{\partial u^{\text{opt}}}{\partial x} + y \cdot \frac{\partial u^{\text{opt}}}{\partial y} - u \right) dx dy = \int \frac{\partial u^{\text{opt}}}{\partial x} d\varsigma_1 + \int \frac{\partial u^{\text{opt}}}{\partial y} d\varsigma_2.$$

Equivalently,

$$\begin{aligned}
& C \cdot (v^{\text{opt}})'(1) - \int_{\frac{2}{3}}^1 v^{\text{opt}}(y) dy + \int_{\frac{2-\sqrt{2}}{3}}^{\frac{2}{3}} v^{\text{opt}}(y) \cdot 3 \left(y - \frac{2-\sqrt{2}}{3} \right) dy = \\
& = -3 \int_{\mathcal{Z}} u^{\text{opt}}(x, y) dx dy.
\end{aligned}$$

By (76), the right-hand side is equal to 0. The function v^{opt} is linear on $\left[\frac{2-\sqrt{2}}{3}, 1\right]$ and, hence, both inequalities (86) and (87) hold as equalities. Therefore, the left-hand side is also 0. Thus ς is an optimal solution to the dual problem as its objective (79) on ς is equal to the optimal value of the primal problem.

D.1 Upper bound on auctioneer's revenue

In Section 4, we showed that the auctioneer's revenue is upper-bounded by

$$B \cdot \inf_{(\varphi_i)_{i \in \mathcal{I}}} \sum_{i \in \mathcal{I}} \left(\int_X \varphi_i^*(x_i) \rho(x) dx + \int_0^1 \varphi_i(z^{B-1}) dz \right). \quad (88)$$

for any number B of bidders, I of items, and any density ρ ; see formula (22). Here we show that this upper bound corresponds to full surplus extraction.

Our goal is to show that the expression (88) equals to the full surplus defined by

$$\sum_{y \in \mathcal{I}} \mathbb{E} \left[\max_{b \in \mathcal{B}} \chi_{b,i} \right], \quad (89)$$

where $\chi_b \in X$ are i.i.d. random vectors distributed with density ρ .

Let ρ_i be the one-dimensional marginals of ρ onto the i -th coordinate. Then (88) equals to $B \sum_{i \in \mathcal{I}} DMK_i$, where

$$DMK_i = \inf_{(\varphi_i)_{i \in \mathcal{I}}} \left(\int_0^1 \varphi_i^*(x_i) \rho_i(x_i) dx_i + \frac{1}{B-1} \int_0^1 \varphi_i(y_i) y_i^{\frac{2-B}{B-1}} dy_i \right).$$

The value DMK_i is nothing else but the value of the dual Monge–Kantorovich problem for the cost function $-x_i y_i$. Adding the terms $\frac{1}{2} \int_0^1 x_i^2 \rho_i(x_i) dx_i$ and $\frac{1}{2(B-1)} \int_0^1 y_i^{2+\frac{2-B}{B-1}} dy_i$ with known value, the reader can easily verify that this problem is equivalent to the transportation problem with the standard cost $\frac{1}{2} |x_i - y_i|^2$. Thus, according to the one-dimensional version of the Brenier theorem, the solution is concentrated on the graph of the mapping T_i given by

$$\int_0^{t_i} \rho_i dx_i = \frac{1}{B-1} \int_0^{T_i(t_i)} y_i^{\frac{2-B}{B-1}} dy_i = T_i^{\frac{1}{B-1}}(t_i)$$

and the cost equals $\int_0^1 x_i T_i(x_i) \rho_i dx_i$. Let F_i be the cumulative distribution function of x_i . Finally, we get

$$DMK_i = \int_0^1 x_i F_i^{B-1}(x_i) \rho_i(x_i) dx_i = \frac{1}{B} \int_0^1 x_i \rho_i(x_i) dF_i^B(x_i) = \frac{1}{B} \mathbb{E} \left[\max_{b \in \mathcal{B}} \chi_{b,i} \right]$$

and conclude that (88) is equal to (89).

E Numerical approach

This section is devoted to computing the auctioneer's optimal revenue and an optimal reduced-form mechanism. By Proposition 1, the auctioneer's problem is

equivalent to the multi-bidder Rochet-Choné problem (11). We describe a numerical approximation scheme for this problem and prove convergence results.

Recall that in the multi-bidder Rochet-Choné problem, we are given the number B of bidders, the set \mathcal{I} of $|\mathcal{I}| = I$ items, and a distribution μ on the set of types $X = [0, 1]^{\mathcal{I}}$ with density ρ . Let η be the majorizing measure equal to the distribution of ξ^{B-1} , where ξ is uniform on $[0, 1]$. The goal is to maximize

$$B \cdot \int_X \left(\langle \nabla u(x), x \rangle - u(x) \right) d\mu(x)$$

over functions $u \in \mathcal{U}_{\text{Lip},1}$ satisfying the majorization constraint

$$\text{law}(u_{x_i}) \preceq \eta,$$

for all $i \in \mathcal{I}$. Recall that $\mathcal{U}_{\text{Lip},1}$ is the set of 1-Lipshitz convex non-decreasing functions $u : X \rightarrow \mathbb{R}_+$, $\text{law}(\xi)$ denotes the distribution of a random variable ξ , and the partial derivative $u_{x_i} = u_{x_i}(\chi)$ is treated as a random variable assuming that its argument χ is distributed according to μ .

We will assume that the distribution μ satisfies the following assumption:

Assumption 1. The density ρ is a continuously differentiable function, and there exist constants $0 < c < C$ such that $c \leq \rho(x) \leq C$ for all $x \in X$.

Outline of the results. The multi-bidder Rochet-Choné problem is a well-defined optimization problem, however, converting it into an algorithm approximating the solution — a numerical approximation scheme — is not straightforward. The first obstacle is that the solution u as well as the input data μ and η are continuous objects. Hence, the problem is to be discretized in a way that solutions of the discrete problems approximate those of the continuous one. The second obstacle is that the majorization constraint, in addition to requiring discretization, is non-linear.

We demonstrate that the majorization constraint is equivalent to a linear constraint suggested by a connection between majorization and martingales and construct provably convergent approximations. As a result, we obtain a finite-dimensional linear program approximating the original Rochet-Choné problem. To summarize, the approach consists of three steps:

1. discretize the set of types $X = [0, 1]^{\mathcal{I}}$ and the distribution μ ;
2. approximate the gradient ∇u and the convexity constraint $u \in \mathcal{U}_{\text{Lip},1}$;
3. linearize and approximate the majorization constraint $\text{law}(u_{x_i}) \preceq \eta$;
4. use an LP solver to find a solution to the resulting linear program.

Section E.1 describes the second step, the third step is discussed in Section E.2, and Section E.3 contains provable heuristics improving the run time. Here we provide a high-level overview.

To discretize the domain $X = [0, 1]^I$, we consider the uniform partition of X into n^I equal cubes and replace the probability distribution μ with the associated sum of point masses. After that, we approximate the initial auction design problem with the corresponding discrete version. The convexity constraint can be written as follows:

$$u(\theta_i) - u(\theta_j) \geq \langle \nabla u(\theta_j), \theta_i - \theta_j \rangle \quad \text{for all } \theta_i, \theta_j \text{ from the discrete lattice.} \quad (90)$$

This approach is based on the algorithm described by [Ekeland and Moreno-Bromberg \(2010\)](#). To approximate the majorization constraint, we use a generalization of Strassen’s theorem [Strassen \(1965\)](#) reducing the constraint to the existence of the supermartingale with the given marginals. To get a finite-dimensional linear program, we discretize the distribution η . The convergence of the discretization is demonstrated in [Theorem 7](#), [Theorem 9](#), and [Corollary 6](#).

In practice, the computation can be sped up by reducing the size of the linear program, which can be achieved via heuristics identifying redundant constraints. The approach of “directional convexity” by [Oberman \(2013\)](#) allows us to reduce the number of convexity constraints in (90) and results in the substantial improvement in computation time.

For simplicity, we focus on the case of a common majorizing measure $\eta = \text{law}(\xi^{B-1})$ with $\xi \sim \text{Uniform}[0, 1]$. The results can be easily extended to the Rochet-Choné problem with general majorization (44) and distinct majorizing measures η_i .

E.1 Convexity constraint approximation

With the continuous problem, we associate its discrete version as described in ([Ekeland and Moreno-Bromberg, 2010](#), Section 3).

Fix a positive integer n . We partition the domain $X = [0, 1]^I$ into n^I equal cubes with the edge length $1/n$. The elements of the partition will be denoted by $\sigma_j^{(n)}$, $1 \leq j \leq n^I$. Denote

$$\Theta_n = \{\theta_j^{(n)} : 1 \leq j \leq n^I\},$$

where $\theta_j^{(n)}$ is the center of the cube $\sigma_j^{(n)}$. Finally, we denote

$$\mu_j^{(n)} = \frac{1}{n^I} \cdot \min\{\rho(x) : x \in \sigma_j^{(n)}\}.$$

Note that the weight sum $\sum_j \mu_j^{(n)}$ is not necessary equal to 1; therefore, we define $\mu_0^{(n)} = 1 - \sum_{j=1}^{n^I} \mu_j^{(n)}$.

For every $1 \leq j \leq n^I$, we associate with the cube $\sigma_j^{(n)}$ the scalar variable $u_j^{(n)}$ that corresponds to the value of the utility function u at $\theta_j^{(n)}$, and the vector variable $p_j^{(n)} = (p_{j,1}^{(n)}, \dots, p_{j,I}^{(n)}) \in \mathbb{R}^{\mathcal{I}}$ that corresponds to the value of ∇u at $\theta_j^{(n)}$. After that, we define the following non-linear program \mathcal{D}_n :

$$\text{maximize: } \sum_j \mu_j^{(n)} \cdot \left(\langle \theta_j^{(n)}, p_j^{(n)} \rangle - u_j^{(n)} \right) \quad (\mathcal{D}_n)$$

subject to:

$$\begin{aligned} (\mathbf{ir}) \quad & u_j^{(n)} \geq 0 && \text{for all } 1 \leq j \leq n^I; \\ (\mathbf{fs}) \quad & 0 \leq p_{j,k}^{(n)} \leq 1 && \text{for all } 1 \leq j \leq n^I, k \in \mathcal{I}; \\ (\mathbf{ic}) \quad & u_i^{(n)} - u_j^{(n)} \geq \langle \theta_i^{(n)} - \theta_j^{(n)}, p_j^{(n)} \rangle && \text{for all } 1 \leq i, j \leq n^I; \\ (\mathbf{mj}) \quad & \mu_0^{(n)} \cdot \delta(x=0) + \sum_j \mu_j^{(n)} \cdot \delta(x=p_{j,k}^{(n)}) \preceq \eta && \text{for all } k \in \mathcal{I}. \end{aligned}$$

Here, the shortcuts **(ir)**, **(fs)**, **(ic)**, and **(mj)** correspond to the individual rationality, feasibility, incentive compatibility, and majorization, respectively. The only non-linear constraint in this program is **(mj)** discussed in the next section.

Given a solution $(\bar{u}^{(n)}, \bar{p}^{(n)})$, we define a function

$$\bar{u}^{(n)}(x) = \max \left\{ 0, \max_j \left\{ u_j^{(n)} + \langle x - \theta_j^{(n)}, p_j^{(n)} \rangle \right\} \right\}.$$

One can easily check that $\bar{u}^{(n)} \in \mathcal{U}_{\text{Lip},1}$, $\bar{u}^{(n)}(\theta_j^{(n)}) = \bar{u}_j^{(n)}$, and $\bar{p}_j^{(n)} \in \partial \bar{u}^{(n)}(\theta_j^{(n)})$ for all $1 \leq j \leq n^I$. Unfortunately, it does not necessary true that $\text{law}(\bar{u}_{x_i}^{(n)}) \preceq \eta$; therefore, the function $\bar{u}^{(n)}$ does not necessary correspond to the interim utility function of a feasible auction mechanism. Nevertheless, we prove that the limiting function \bar{u} satisfies the majorization constraint:

Proposition 8. *There exists a subsequence $\{\bar{u}_{n_k}\} \subset \{\bar{u}_n\}$ such that:*

- (a) *the subsequence $\{\bar{u}_{n_k}\}$ converges uniformly to \bar{u} ;*
- (b) *the function $\bar{u} \in \mathcal{U}_{\text{Lip},1}$ and $\text{law}(\bar{u}_{x_i}) \preceq \eta$ for all $i \in \mathcal{I}$;*
- (c) $\lim_{k \rightarrow \infty} \sum_j \mu_j^{(n_k)} \cdot \left(\langle \theta_j^{(n_k)}, \bar{p}_j^{(n_k)} \rangle - \bar{u}_j^{(n_k)} \right) = \int_X (\langle x, \nabla \bar{u}(x) \rangle - \bar{u}(x)) d\mu$

Before proving [Proposition 8](#), we need the following technical result:

Lemma 12. Let Q be a compact subset of \mathbb{R}^I . Consider a function $\phi(\theta, z, p) \in C^1(X \times \mathbb{R} \times Q \rightarrow \mathbb{R})$. Let $f_k: X \rightarrow \mathbb{R}$ be a family of convex functions such that $\partial f_k(\theta) \in Q$ for all $\theta \in X$, whose uniform limit is f . Then

$$\lim_{k \rightarrow \infty} \sum_j \mu_j^{(n)} \cdot \phi\left(\theta_j^{(n)}, f_n(\theta_j^{(n)}), \nabla f_n(\theta_j^{(n)})\right) = \int_X \phi(\theta, f(\theta), \nabla f(\theta)) d\mu.$$

Proof of Lemma 12. The proof is based on the following result.

Lemma ((Ekeland and Moreno-Bromberg, 2010, Lemma A.6)). Under the assumptions of Lemma 12, for every $\varepsilon > 0$ there exists $K \in \mathbb{N}$ such that

$$\left| \frac{1}{n^I} \sum_j \phi\left(\theta_j^{(n)}, f_n(\theta_j^{(n)}), \nabla f_n(\theta_j^{(n)})\right) - \int_X \phi(\theta, f_n(\theta), \nabla f_n(\theta)) d\theta \right| < \varepsilon$$

for all $n > K$.

Consider a function $\psi(\theta, z, p) = \rho(\theta) \cdot \phi(\theta, z, p)$, where ρ is the density of μ . The function ψ is continuously differentiable; therefore, for every $\varepsilon > 0$ there exists $K \in \mathbb{N}$ such that

$$\begin{aligned} & \left| \frac{1}{n^I} \sum_j \psi\left(\theta_j^{(n)}, f_n(\theta_j^{(n)}), \nabla f_n(\theta_j^{(n)})\right) - \int_X \psi(\theta, f_n(\theta), \nabla f_n(\theta)) d\theta \right| = \\ & = \left| \frac{1}{n^I} \sum_j \rho(\theta_j^{(n)}) \cdot \phi\left(\theta_j^{(n)}, f_n(\theta_j^{(n)}), \nabla f_n(\theta_j^{(n)})\right) - \int_X \phi(\theta, f_n(\theta), \nabla f_n(\theta)) d\mu \right| < \varepsilon \end{aligned} \quad (91)$$

for all $n > K$.

Since all f_n are convex and $\{f_n\}$ converges uniformly to f , the sequence of gradients $\{\nabla f_n(\theta)\}$ converges to $\nabla f(\theta)$ for almost all $\theta \in X$. Thus it follows from the continuity of ϕ that

$$\lim_{n \rightarrow \infty} \phi(\theta, f_n(\theta), \nabla f_n(\theta)) = \phi(\theta, f(\theta), \nabla f(\theta))$$

for almost all $\theta \in X$. We may assume that the family of functions f_n is uniformly bounded. Therefore, there exists a constant M such that

$$|\phi(\theta, f_n(\theta), \nabla f_n(\theta))| < M$$

for all $\theta \in X$ and for all n . Thus it follows from Lebesgue's dominated convergence theorem that

$$\lim_{n \rightarrow \infty} \int_X \phi(\theta, f_n(\theta), \nabla f_n(\theta)) d\mu = \int_X \phi(\theta, f(\theta), \nabla f(\theta)) d\mu. \quad (92)$$

Finally, we have

$$\begin{aligned}
& \left| \frac{1}{n^I} \sum_j \rho(\theta_j^{(n)}) \cdot \phi \left(\theta_j^{(n)}, f_n(\theta_j^{(n)}), \nabla f_n(\theta_j^{(n)}) \right) - \right. \\
& \quad \left. - \sum_j \mu_j^{(n)} \cdot \phi \left(\theta_j^{(n)}, f_n(\theta_j^{(n)}), \nabla f_n(\theta_j^{(n)}) \right) \right| \leq \\
& \leq M \cdot \sum_j \left| \frac{\rho(\theta_j^{(n)})}{n^I} - \mu_j^{(n)} \right| \leq M \cdot \sup_{j; x, y \in \sigma_j^{(n)}} |\rho(x) - \rho(y)|,
\end{aligned} \tag{93}$$

and the latter expression tends to 0 as $n \rightarrow \infty$ by the uniform continuity of ρ .

Combining the expressions (91)–(93), we obtain the desired convergence equality. \square

Proof of Proposition 8. Since $\bar{u}^{(n)} \in \mathcal{U}_{\text{Lip},1}$ and this set space is sequentially compact in the uniform convergence topology (Lemma 5), there is a subsequence $\{\bar{u}_{n_k}\}$ that converges uniformly to the function $\bar{u} \in \mathcal{U}_{\text{Lip},1}$.

To prove the majorization condition, it is sufficient to check that for any continuously differentiable non-decreasing convex function φ , we have

$$\int \varphi(\bar{u}_{x_i}) d\mu \leq \int_0^1 \varphi(x) d\eta(x). \tag{94}$$

It follows from Lemma 12 that

$$\left| \int \varphi(\bar{u}_{x_i}) d\mu - \sum_j \mu_j^{(n_k)} \cdot \varphi(\bar{p}_{j,i}^{(n_k)}) \right| \leq \varepsilon(k)$$

and $\varepsilon(k) \rightarrow 0$ as $k \rightarrow +\infty$. It follows from the **(mj)** constraint that

$$\sum_j \mu_j^{(n_k)} \cdot \varphi(\bar{p}_{j,i}^{(n_k)}) \leq \int_0^1 \varphi(x) d\eta(x);$$

therefore, letting k tend to $+\infty$, we conclude that the inequality (94) holds. The point (c) also follows directly from Lemma 12. \square

Let $u^{\text{opt}} \in \mathcal{U}_{\text{Lip},1}$ be the optimum of the multi-bidder Rochet-Choné problem. For each positive integer n , denote

$$\begin{aligned}
u_j^{\text{opt},(n)} &= n^I \cdot \int_{\sigma_j^{(n)}} u^{\text{opt}}(x) dx, \\
p_{j,k}^{\text{opt},(n)} &= n^I \cdot \int_{\sigma_j^{(n)}} u_{x_k}^{\text{opt}}(x) dx.
\end{aligned}$$

Proposition 9. *The variables $u_j^{\text{opt},(n)}$ and $p_j^{\text{opt},(n)}$ satisfy all the constraints of the program \mathcal{D}_n .*

Proof. The constraints **(ir)** and **(fs)** follow from the inequalities $u^{\text{opt}}(x) \geq 0$ and $0 \leq u_{x_k}^{\text{opt}}(x) \leq 1$.

Since $u^{\text{opt}}(x)$ is convex, we have

$$u^{\text{opt}}(x - \theta_j^{(n)} + \theta_i^{(n)}) - u^{\text{opt}}(x) \geq \langle \theta_i^{(n)} - \theta_j^{(n)}, \nabla u^{\text{opt}}(x) \rangle$$

for almost all $x \in \sigma_j^{(n)}$. Integrating this inequality over the cube $\sigma_j^{(n)}$, we conclude that

$$u_i^{\text{opt},(n)} - u_j^{\text{opt},(n)} \geq \langle \theta_i^{(n)} - \theta_j^{(n)}, p_j^{\text{opt},(n)} \rangle.$$

Thus the constraint **(ic)** holds.

Consider any non-decreasing convex function φ . Since $\text{law}(u_{x_i}^{\text{opt}}) \preceq \eta_i$, we conclude that

$$\int_0^1 \varphi(x) \, d\eta(x) \geq \sum_j \int_{\sigma_j^{(n)}} \varphi(u_{x_i}^{\text{opt}}) \, d\mu.$$

Since $\varphi(x)$ is non-decreasing, we have

$$\int_{\sigma_j^{(n)}} \varphi(u_{x_i}^{\text{opt}}) \, d\mu \geq \mu_j^{(n)} \cdot n^I \cdot \int_{\sigma_j^{(n)}} \varphi(u_{x_i}^{\text{opt}}) \, dx + (\mu(\sigma_j^{(n)}) - \mu_j^{(n)}) \cdot \varphi(0).$$

Finally, it follows from Jensen's inequality that

$$n^I \cdot \int_{\sigma_j^{(n)}} \varphi(u_{x_i}^{\text{opt}}) \, dx \geq \varphi \left(n^I \cdot \int_{\sigma_j^{(n)}} u_{x_i}^{\text{opt}}(x) \, dx \right) = \varphi(p_{j,i}^{\text{opt},(n)})$$

Thus

$$\int_0^1 \varphi(x) \, d\eta(x) \geq \sum_j \mu_j^{(n)} \cdot \varphi(p_{j,i}^{\text{opt},(n)}) + \mu_0^{(n)} \cdot \varphi(0).$$

Since this inequality holds for all φ , we conclude that the constraint **(mj)** holds. \square

The next result demonstrates that the optimal revenue in the continuous problem is approximated by its discretization.

Proposition 10. *The following identity holds:*

$$\lim_{n \rightarrow \infty} \sum_j \mu_j^{(n)} \cdot \left(\langle \theta_j^{(n)}, p_j^{\text{opt},(n)} \rangle - u_j^{\text{opt},(n)} \right) = \int_X (\langle x, \nabla u^{\text{opt}}(x) \rangle - u^{\text{opt}}(x)) \, d\mu.$$

Proof. The definition of $u_j^{\text{opt},(n)}$ and $p_j^{\text{opt},(n)}$ implies that

$$\begin{aligned} \int_X (\langle x, \nabla u^{\text{opt}}(x) \rangle - u^{\text{opt}}(x)) \, d\mu - \sum_j \mu_j^{(n)} \cdot (\langle \theta_j^{(n)}, p_j^{\text{opt},(n)} \rangle - u_j^{\text{opt},(n)}) &= \\ &= \sum_j \int_{\sigma_j^{(n)}} (\langle x, \nabla u^{\text{opt}}(x) \rangle - u^{\text{opt}}(x)) \cdot (\rho(x) - n^I \mu_j^{(n)}) \, dx \leq \\ &\leq \sup_{x \in X} |\langle x, \nabla u^{\text{opt}}(x) \rangle - u^{\text{opt}}(x)| \cdot \sup_{j, x \in \sigma_j^{(n)}} |\rho(x) - n^I \mu_j^{(n)}|. \end{aligned}$$

The function $|\langle x, \nabla u^{\text{opt}}(x) \rangle - u^{\text{opt}}(x)|$ is bounded, and the result follows from the uniform continuity of ρ . \square

Putting all the pieces together, we obtain the following convergence result.

Theorem 7.

(a) *The function $\bar{u} = \lim_{k \rightarrow \infty} \bar{u}^{(n_k)}$ is a solution to the multi-bidder Rochet-Choné problem.*

$$(b) \lim_{k \rightarrow \infty} \sum_j \mu_j^{(n_k)} \cdot (\langle \theta_j^{(n_k)}, \bar{p}_j^{(n_k)} \rangle - \bar{u}_j^{(n_k)}) = \max_{\substack{u \in \mathcal{U}_{\text{Lip},1}, \\ \text{law}(u_{x_i}) \preceq \eta}} \int_X (\langle x, \nabla u(x) \rangle - u(x)) \, d\mu.$$

Proof. Recall that $u^{\text{opt}}(x)$ is a solution to the multi-bidder Rochet-Choné problem. Hence, it follows from [Proposition 10](#) that

$$\begin{aligned} \max_{\substack{u \in \mathcal{U}_{\text{Lip},1}, \\ \text{law}(u_{x_i}) \preceq \eta}} \int_X (\langle x, \nabla u(x) \rangle - u(x)) \, d\mu &= \int_X (\langle x, \nabla u^{\text{opt}}(x) \rangle - u^{\text{opt}}(x)) \, d\mu = \\ &= \lim_{n \rightarrow \infty} \sum_j \mu_j^{(n)} \cdot (\langle \theta_j^{(n)}, p_j^{\text{opt},(n)} \rangle - u_j^{\text{opt},(n)}). \end{aligned}$$

Since $(u_j^{\text{opt},(n)}, p_j^{\text{opt},(n)})$ satisfies all the constraints of the problem \mathcal{D}_n and $(\bar{u}_j^{(n)}, \bar{p}_j^{(n)})$ is an optimal solution to this problem, we have

$$\begin{aligned} \sum_j \mu_j^{(n)} \cdot (\langle \theta_j^{(n)}, p_j^{\text{opt},(n)} \rangle - u_j^{\text{opt},(n)}) &\leq \sum_j \mu_j^{(n)} \cdot (\langle \theta_j^{(n)}, \bar{p}_j^{(n)} \rangle - \bar{u}_j^{(n)}) \\ \lim_{n \rightarrow \infty} \sum_j \mu_j^{(n)} \cdot (\langle \theta_j^{(n)}, p_j^{\text{opt},(n)} \rangle - u_j^{\text{opt},(n)}) &\leq \lim_{n \rightarrow \infty} \sum_j \mu_j^{(n)} \cdot (\langle \theta_j^{(n)}, \bar{p}_j^{(n)} \rangle - \bar{u}_j^{(n)}). \end{aligned}$$

By [Proposition 8](#),

$$\lim_{n \rightarrow \infty} \sum_j \mu_j^{(n)} \cdot (\langle \theta_j^{(n)}, \bar{p}_j^{(n)} \rangle - \bar{u}_j^{(n)}) = \int_X (\langle x, \nabla \bar{u}(x) \rangle - \bar{u}(x)) \, d\mu.$$

Thus

$$\int_X (\langle x, \nabla \bar{u}(x) \rangle - \bar{u}(x)) \, d\mu \geq \max_{\substack{u \in \mathcal{U}_{\text{Lip},1}, \\ \text{law}(u_{x_i}) \preceq \eta}} \int_X (\langle x, \nabla u(x) \rangle - u(x)) \, d\mu.$$

At the same time, by [Proposition 8](#), we have $\bar{u} \in \mathcal{U}_{\text{Lip},1}$ and $\text{law}(\bar{u}_{x_i}) \preceq \eta$ for all $i \in \mathcal{I}$. Thus the equality holds; therefore, \bar{u} is a solution the multi-bidder Rochet-Choné problem and

$$\lim_{k \rightarrow \infty} \sum_j \mu_j^{(n)} \cdot (\langle \theta_j^{(n_k)}, \bar{p}_j^{(n_k)} \rangle - \bar{u}_j^{(n_k)}) = \max_{\substack{u \in \mathcal{U}_{\text{Lip},1}, \\ \text{law}(u_{x_i}) \preceq \eta}} \int_X (\langle x, \nabla u(x) \rangle - u(x)) \, d\mu.$$

□

E.2 Approximation of the majorization constraints

The majorization constraint **(mj)** is non-linear. An equivalent linear constraint can be obtained using the following characterization of the majorization order.

Theorem 8 (([Shaked and Shanthikumar, 2007](#), Theorem 4.A.5)). *Two random variables X and Y satisfy $X \preceq Y$ if and only if there exist two random variables \hat{X} and \hat{Y} defined on the same probability space such that*

$$\begin{aligned} \text{law}(X) &= \text{law}(\hat{X}), \\ \text{law}(Y) &= \text{law}(\hat{Y}), \end{aligned}$$

and $\{\hat{X}, \hat{Y}\}$ is a supermartingale, that is,

$$\mathbb{E}[\hat{Y} \mid \hat{X}] \geq \hat{X} \quad \text{almost surely.}$$

Using this criterion, we reformulate the **(mj)**-constraints as a condition of the existence of the joint distribution of \hat{X} and \hat{Y} . In what follows, we fix an item $k \in \mathcal{I}$.

Proposition 11. *Denote $X = T = [0, 1]$, $J = \{1, 2, \dots, n^I\}$, and $\bar{J} = \{0\} \cup J$. The following statements are equivalent:*

- (a) *The majorization condition $\mu_0^{(n)} \cdot \delta(x = 0) + \sum_j \mu_j^{(n)} \cdot \delta(x = p_{j,k}^{(n)}) \preceq \eta$ holds.*

(b) *There exists a probability distribution π concentrated on $X \times T$ such that*

$$\begin{aligned} \text{pr}_T \pi &= \eta, \\ \text{pr}_X \pi &= \mu_0^{(n)} \cdot \delta(x=0) + \sum_j \mu_j^{(n)} \cdot \delta(x=p_{j,k}^{(n)}), \\ \int_T t \cdot d\pi(x, t) &\geq x \cdot \text{pr}_X \pi(\{x\}) \quad \text{for all } x \in X, \end{aligned}$$

where $\text{pr}_X \pi$ and $\text{pr}_T \pi$ denotes the marginals of π on X and T , respectively.

(c) *There exists a probability distribution π concentrated on $\bar{J} \times T$ such that*

$$\begin{aligned} \text{pr}_T \pi &= \eta, \\ \text{pr}_{\bar{J}} \pi(j) &= \mu_j^{(n)} \quad \text{for all } j \in \bar{J}, \\ \int_T t \cdot d\pi(j, t) &\geq p_{j,k}^{(n)} \cdot \mu_j^{(n)} \quad \text{for all } j \in J. \end{aligned}$$

(d) *There exists a (not necessary probability) measure π concentrated on $J \times T$ such that*

$$\begin{aligned} \text{pr}_T \pi &\leq \eta, \\ \text{pr}_J \pi(j) &\leq \mu_j^{(n)} \quad \text{for all } j \in J, \\ \int_T t \cdot d\pi(j, t) &\geq p_{j,k}^{(n)} \cdot \mu_j^{(n)} \quad \text{for all } j \in J. \end{aligned}$$

Proof. The equivalence (a) \Leftrightarrow (b) is a reformulation of [Theorem 8](#). The distribution π can be considered as the joint law of \hat{X} and \hat{Y} .

(c) \Rightarrow (b). Let π be a distribution satisfying all the conditions of (c). Consider a mapping $f: \bar{J} \rightarrow X$ defined as $f(0) = 0$ and $f(j) = p_{j,k}^{(n)}$ for all $j \in J$. Define by $\hat{\pi}$ the pushforward measure $f_{\#} \pi$ concentrated on $X \times T$. It follows directly from the construction that

$$\text{pr}_T \hat{\pi} = \eta \quad \text{and} \quad \text{pr}_X \hat{\pi} = \mu_0^{(n)} \cdot \delta(x=0) + \sum_j \mu_j^{(n)} \cdot \delta(x=p_{j,k}^{(n)}).$$

Finally, we need to check the inequality $\int_T t \cdot d\pi(x, t) \geq x \cdot \text{pr}_X \pi(\{x\})$. If $x = 0$, there is nothing to prove. Otherwise,

$$\int_T t d\hat{\pi}(x, t) = \sum_{j \in J: f(j)=x} \int_T t d\pi(j, t) \geq \sum_{j \in J: p_{j,k}^{(n)}=x} p_{j,k}^{(n)} \cdot \mu_j^{(n)} = x \cdot \text{pr}_X \pi(\{x\}).$$

Thus $\hat{\pi}$ satisfies all the restrictions of (b).

(b) \Rightarrow (c). Let π be a distribution satisfying all the conditions of (b). For each $j \in J$, define

$$\pi_j = \frac{\mu_j^{(n)}}{\text{pr}_X \pi(\{p_{j,k}^{(n)}\})} \cdot \pi|_{x=p_{j,k}^{(n)}},$$

and

$$\pi_0 = \frac{\mu_0^{(n)}}{\text{pr}_X \pi(\{0\})} \cdot \pi|_{x=0}.$$

One can check easily that $\pi = \delta_0 \otimes \pi_0 + \sum_{j \in J} \delta_{p_{j,k}^{(n)}} \otimes \pi_j$, where δ_x is the Dirac delta measure concentrated at a point x .

Define a measure $\hat{\pi} = \sum_{j \in \bar{J}} \delta_j \otimes \pi_j$ concentrated on $\bar{J} \times T$. We have

$$\begin{aligned} \text{pr}_T \hat{\pi} &= \sum_{j \in \bar{J}} \text{pr}_T \pi_j = \text{pr}_T \pi = \eta, \\ \text{pr}_{\bar{J}} \hat{\pi}(j) &= |\pi_j| = \mu_j^{(n)} \quad \text{for all } j \in \bar{J}, \\ \int_T t \cdot d\hat{\pi}(j, t) &= \int_T t \cdot d\pi_j(t) = \\ &= \frac{\mu_j^{(n)}}{\text{pr}_X(\{p_{j,k}^{(n)}\})} \int_T t \cdot d\pi(p_{j,k}^{(n)}, t) \geq \mu_j^{(n)} \cdot p_{j,k}^{(n)} \quad \text{for all } j \in J. \end{aligned}$$

Thus $\hat{\pi}$ satisfies all the restrictions of (c).

(c) \Rightarrow (d). If a distribution π satisfies all the restrictions of (c), then the restriction of π to the set $J \times T$ satisfies all the restrictions of (d).

(d) \Rightarrow (c). Let π be a measure concentrated on $J \times T \subset \bar{J} \times T$ satisfying all the restrictions of (d). One can easily prove that there exists a distribution $\hat{\pi}$ concentrated on $\bar{J} \times T$ such that $\pi \leq \hat{\pi}$, $\text{pr}_T \hat{\pi} = \eta$, and $\text{pr}_{\bar{J}} \hat{\pi}(j) = \mu_j^{(n)}$ for all $j \in \bar{J}$. As a consequence,

$$\int_T t d\hat{\pi}(j, t) \geq \int_T t d\pi(j, t) \geq \mu_j^{(n)} \cdot p_{j,k}^{(n)} \quad \text{for all } j \in J.$$

Thus the distribution $\hat{\pi}$ satisfies all the restrictions of (c). \square

The measure π obtained in [Proposition 11](#)(d) is not discrete. To discretize this measure, we discretize the space T . Let $0 = q_0 < q_1 < \dots < q_M = 1$ be any partition of the space $T = [0, 1]$. For each $1 \leq m \leq M$, denote

$$w_m = \eta([q_{m-1}, q_m]), \quad t_m = \frac{1}{w_m} \int_{q_{m-1}}^{q_m} t d\eta, \quad \eta_m = \frac{1}{w_m} \cdot \eta|_{[q_{m-1}, q_m]}.$$

As a discrete approximation of π , we will only consider measures of the form

$$\pi = \sum_{1 \leq j \leq n^I, 1 \leq m \leq M} \pi_{j,m} \cdot \delta_j \otimes \eta_m \quad (95)$$

for some non-negative coefficients $\pi_{j,m}$. The following statement characterizes all such measures that satisfy the restrictions of [Proposition 11](#)(d).

Lemma 13. *The measure π defined in (95) satisfies all the restrictions of [Proposition 11](#)(d) if and only if the following inequalities hold:*

$$\begin{aligned} \sum_{1 \leq j \leq n^I} \pi_{j,m} &\leq w_m && \text{for all } 1 \leq m \leq M, \\ \sum_{1 \leq m \leq M} \pi_{j,m} &\leq \mu_j^{(n)} && \text{for all } 1 \leq j \leq n^I, \\ \sum_{1 \leq m \leq M} t_m \cdot \pi_{j,m} &\geq p_{j,k}^{(n)} \cdot \mu_j^{(n)} && \text{for all } 1 \leq j \leq n^I. \end{aligned}$$

This suggests considering the following linear problem.

Definition 2. Given a partition $0 = q_0 < q_1 < \dots < q_M = 1$, consider the following linear problem $\mathcal{D}_{n,M}$:

$$\text{maximize: } \sum_j \mu_j^{(n)} \cdot \left(\langle \theta_j^{(n)}, p_j^{(n)} \rangle - u_j^{(n)} \right) \quad (\mathcal{D}_{n,M})$$

subject to:

$$\begin{aligned} \text{(ir)} \quad u_j^{(n)} &\geq 0 && \text{for all } 1 \leq j \leq n^I; \\ \text{(fs)} \quad 0 &\leq p_{j,k}^{(n)} \leq 1 && \text{for all } 1 \leq j \leq n^I, k \in \mathcal{I}; \\ \text{(ic)} \quad u_i^{(n)} - u_j^{(n)} &\geq \langle \theta_i^{(n)} - \theta_j^{(n)}, p_j^{(n)} \rangle && \text{for all } 1 \leq i, j \leq n^I; \\ \text{(mj-T)} \quad \sum_{1 \leq j \leq n^I} \pi_{j,m,k}^{(n)} &\leq w_m && \text{for all } 1 \leq m \leq M, k \in \mathcal{I}; \\ \text{(mj-J)} \quad \sum_{1 \leq m \leq M} \pi_{j,m,k}^{(n)} &\leq \mu_j^{(n)} && \text{for all } 1 \leq j \leq n^I, k \in \mathcal{I}; \\ \text{(mj-E)} \quad \sum_{1 \leq m \leq M} t_m \cdot \pi_{j,m,k}^{(n)} &\geq p_{j,k}^{(n)} \cdot \mu_j^{(n)} && \text{for all } 1 \leq j \leq n^I, k \in \mathcal{I}; \\ \text{(mj-P)} \quad \pi_{j,m,k}^{(n)} &\geq 0 && \text{for all } j, m, k. \end{aligned}$$

A direct consequence of [Proposition 11](#) and [Lemma 13](#) is the following connection between the problems \mathcal{D}_n and $\mathcal{D}_{n,M}$.

Corollary 5. If $(u_j^{(n)}, p_j^{(n)}, \pi_{j,m,k}^{(n)})$ satisfies all the constraints of the problem $\mathcal{D}_{n,M}$, then $(u_j^{(n)}, p_j^{(n)})$ satisfies all the constraints of \mathcal{D}_n .

The constraint **(fs)** partially follows from the constraints **(mj-J)**, **(mj-E)**, and **(mj-P)**:

$$\begin{aligned} p_{j,k}^{(n)} \cdot \mu_j^{(n)} &\leq \sum_{1 \leq m \leq M} t_m \cdot \pi_{j,m,k}^{(n)} \leq \max_{1 \leq m \leq M} t_m \cdot \sum_{1 \leq m \leq M} \pi_{j,m,k}^{(n)} \leq t_M \cdot \mu_j^{(n)} \\ &\Rightarrow p_{j,k}^{(n)} \leq t_M \leq 1. \end{aligned}$$

Thus the linear problem $\mathcal{D}_{n,M}$ is equivalent to the problem $\mathcal{D}'_{n,M}$, where the **(fs)**-constraints are replaced with the following:

$$\textbf{(fs')} \quad p_{j,k}^{(n)} \geq 0 \quad \text{for all } 1 \leq j \leq n^I, k \in \mathcal{I}.$$

Our goal is to prove that the sequence $(\bar{u}_j^{(n,M)}, \bar{p}_j^{(n,M)}, \bar{\pi}_{j,m,k}^{(n,M)})$ of optimal solutions to the problem $\mathcal{D}_{n,M}$ contains a maximizing subsequence to the problem \mathcal{D}_n as $M \rightarrow \infty$. In order to do it, we formulate a dual problem to $\mathcal{D}'_{n,M}$.

Lemma 14. Consider the following finite-dimensional convex programs:

$$\begin{array}{lll} (\mathcal{P}) \max & c^T x & (\mathcal{D}) \min \quad b^T y_1 \quad (\mathcal{C}) \min \quad b^T y \\ \text{s.t.} & Ax \leq b, & \text{s.t.} \quad A^T y_1 + Q^T y_2 \geq c \quad \text{s.t.} \quad y^T A x \geq c^T x \\ & Qx \leq 0, & y_1 \geq 0, y_2 \geq 0; & \forall x \geq 0: Qx \leq 0, \\ & x \geq 0; & & y \geq 0. \end{array}$$

Assume that the problem (\mathcal{P}) is feasible and bounded. Then

- (a) if (y_1, y_2) satisfies all the restrictions of (\mathcal{D}) , then y_1 satisfies all the restrictions of (\mathcal{C}) ;
- (b) if (\bar{y}_1, \bar{y}_2) solves the problem (\mathcal{D}) , then \bar{y}_1 solves the problem (\mathcal{C}) ;
- (c) the strong duality holds: $\max_{\mathcal{P}} c^T x = \min_{\mathcal{D}} b^T y_1 = \min_{\mathcal{C}} b^T y$.

Proof. (a) Consider any couple of non-negative vectors (y_1, y_2) satisfying the inequality

$$A^T y_1 + Q^T y_2 \geq c \quad \Leftrightarrow \quad y_1^T A + y_2^T Q \geq c^T.$$

Then, for any vector $x \geq 0$, we have

$$y_1^T A x + y_2^T Q x \geq c^T x.$$

Assume in addition that $Qx \leq 0$. Then it follows from the non-negativity of y_2 that $y_2^T Qx \leq 0$; therefore,

$$y_1^T Ax \geq y_1^T Ax + y_2^T Qx \geq c^T x.$$

So, the vector y_1 satisfies all the restrictions of the problem (\mathcal{C}) .

(b) and (c). First, we prove the weak duality

$$\min_{\mathcal{C}} b^T y \geq \max_{\mathcal{P}} c^T x.$$

Let \bar{x} be a solution to (\mathcal{P}) and let y be any vector satisfying all the restrictions of the problem (\mathcal{C}) . We have

$$y^T A\bar{x} \geq c^T \bar{x}.$$

Since $A\bar{x} \leq b$ and $y \geq 0$, we conclude that $y^T b \geq y^T A\bar{x}$. Thus

$$b^T y = y^T b \geq c^T \bar{x} = \max_{\mathcal{P}} c^T x \quad \Rightarrow \quad \min_{\mathcal{C}} b^T y \geq \max_{\mathcal{P}} c^T x.$$

Let (\bar{y}_1, \bar{y}_2) be a solution to the problem (\mathcal{D}) . It follows from the duality theorem that $b^T \bar{y}_1 = \max_{\mathcal{P}} c^T x$. In addition, \bar{y}_1 satisfies all the restrictions of the problem (\mathcal{C}) ; therefore,

$$\max_{\mathcal{P}} c^T x = b^T \bar{y}_1 = \min_{\mathcal{C}} b^T y.$$

Thus \bar{y}_1 is a solution to (\mathcal{C}) and the strong duality holds. \square

Using this lemma, we formulate a dual convex program to $\mathcal{D}'_{n,m}$.

Definition 3. Given a positive integer n and a partition $0 = q_0 < q_1 < \dots < q_M = 1$, we define a convex program $\mathcal{D}^*_{n,M}$ as follows:

$$\begin{aligned} \text{minimize: } & \sum_{k \in \mathcal{I}} \left(\sum_{1 \leq m \leq M} \varphi_{m,k}^{(n)} \cdot w_m + \sum_{1 \leq j \leq n^I} \psi_{j,k}^{(n)} \cdot \mu_j^{(n)} \right) & (\mathcal{D}^*_{n,M}) \\ \text{subject to: } & \varphi_{m,k}^{(n)} \geq 0, \psi_{j,k}^{(n)} \geq 0, c_{j,k}^{(n)} \geq 0 \\ & \text{(lt)} \quad \varphi_{m,k}^{(n)} + \psi_{j,k}^{(n)} \geq t_m \cdot c_{j,k}^{(n)} \\ & \text{(c-def)} \quad \sum_j \mu_j^{(n)} \cdot \left(\langle \theta_j^{(n)}, p_j^{(n)} \rangle - u_j^{(n)} \right) \leq \sum_j \mu_j^{(n)} \cdot \langle c_j^{(n)}, p_j^{(n)} \rangle \\ & \text{for all } (u_j^{(n)}, p_j^{(n)}) \text{ satisfying (ir), (fs'), and (ic).} \end{aligned}$$

Here, the shortcut **(lt)** indicates a relation to the Legendre transform, and **(c-def)**, to the set of vector fields \mathcal{C} from Appendix C defined by

$$\int (\langle x, \nabla u(x) \rangle - u(x)) d\mu \leq \int \langle c(x), \nabla u(x) \rangle d\mu$$

for all convex non-decreasing u with $u(0) = 0$.

Proposition 12. *The strong duality holds: $\max \mathcal{D}'_{n,M} = \min \mathcal{D}^*_{n,M}$.*

Proof. Let $\varphi_{m,k}^{(n)}$ be a dual variable for the **(mj-T)**-constraint $\sum_j \pi_{j,m,k}^{(n)} \leq w_m$, let $\psi_{j,m}^{(n)}$ be a dual variable for the **(mj-J)**-constraint $\sum_m \pi_{j,m,k}^{(n)} \leq \mu_j^{(n)}$, and let $c_{j,k}^{(n)}$ be a dual variable for the **(mj-E)**-constraint $p_{j,k}^{(n)} \cdot \mu_j^{(n)} - \sum_m t_m \cdot \pi_{j,m,k}^{(n)} \leq 0$. The following duality equation follows from Lemma 14 applied to the linear program $\mathcal{D}'_{n,M}$:

$$\max \mathcal{D}'_{n,M} = \min \sum_{k=1}^{\mathcal{I}} \left(\sum_{1 \leq m \leq M} \varphi_{m,k}^{(n)} \cdot w_m + \sum_{1 \leq j \leq n^I} \psi_{j,k}^{(n)} \cdot \mu_j^{(n)} \right),$$

where minimum in the right-hand side is taken over all non-negative variables $(\varphi_{m,k}^{(n)}, \psi_{j,k}^{(n)}, c_{j,k}^{(n)})$ such that the inequality

$$\begin{aligned} \sum_{m,k} \varphi_{m,k}^{(n)} \cdot \sum_j \pi_{j,m,k}^{(n)} + \sum_{j,k} \psi_{j,k}^{(n)} \cdot \sum_m \pi_{j,m,k}^{(n)} + \sum_{j,k} c_{j,k}^{(n)} \cdot \left(p_{j,k}^{(n)} \cdot \mu_j^{(n)} - \sum_m t_m \cdot \pi_{j,m,k}^{(n)} \right) \\ \geq \sum_j \mu_j^{(n)} \cdot \left(\langle \theta_j^{(n)}, p_j^{(n)} \rangle - u_j^{(n)} \right) \end{aligned} \quad (96)$$

holds for all $(u_j^{(n)}, p_j^{(n)}, \pi_{j,m,k}^{(n)})$ satisfying the constraints **(ir)**, **(fs')**, and **(ic)**.

After the rearrangement of the left-hand side terms, inequality (96) transforms into

$$\begin{aligned} \sum_{j,m,k} \pi_{j,m,k}^{(n)} \cdot \left(\varphi_{m,k}^{(n)} + \psi_{j,k}^{(n)} - t_m \cdot c_{j,k}^{(n)} \right) + \sum_j \mu_j^{(n)} \cdot \langle p_j^{(n)}, c_j^{(n)} \rangle \geq \\ \geq \sum_j \mu_j^{(n)} \cdot \left(\langle \theta_j^{(n)}, p_j^{(n)} \rangle - u_j^{(n)} \right) \end{aligned} \quad (97)$$

Since the constraints **(ir)**, **(fs')**, and **(ic)** do not contain any restrictions on $\pi_{j,m,k}^{(n)}$, it follows from (97) that the **(lt)**-constraints hold:

$$\varphi_{m,k}^{(n)} + \psi_{j,k}^{(n)} - t_m \cdot c_{j,k}^{(n)} \geq 0 \quad \text{for all } 1 \leq j \leq n^I, 1 \leq m \leq M, k \in \mathcal{I}.$$

Substituting $\pi_{j,m,k}^{(n)} = 0$ into (97), we conclude that **(c-def)**-constraint holds. Vice versa, if all the **(lt)**- and **(c-def)**-constraints hold, then the inequality (97) holds for all $(u_j^{(n)}, p_j^{(n)})$ satisfying the constraints **(ir)**, **(fs')**, and **(ic)**, and for all non-negative $\pi_{j,m,k}^{(n)}$. \square

Remark 2. The problem $\mathcal{D}_{n,M}^*$ can be seen as a discrete approximation of the dual problem described in Theorem 5. The variable $c_{j,k}^{(n)}$ corresponds to the value of the component c_k of the vector field $c = (c_1, \dots, c_I)$ at the point $\theta_j^{(n)}$; the constraint **(c-def)** is a discrete approximation of the inequality

$$\int (\langle x, \nabla u(x) \rangle - u(x)) d\mu \leq \int \langle c(x), \nabla u(x) \rangle d\mu.$$

The variable $\varphi_{m,k}^{(n)}$ corresponds to the value of the function $\varphi_k(x)$ at the point t_m . By the constraint **(lt)**, the optimal value of $\varphi_{m,k}^{(n)}$ is equal to

$$\bar{\varphi}_{m,k}^{(n)} = \max \left\{ 0, \max_j (t_m \cdot \bar{c}_{j,k}^{(n)} - \bar{\psi}_{j,k}^{(n)}) \right\} = \bar{\varphi}_k(t_m),$$

where the function $\bar{\varphi}(x)$ is convex and non-negative as a maximum of non-decreasing linear functions. Similarly, the optimal value of $\psi_{j,k}^{(n)}$ is equal to $\bar{\varphi}_k^*(c_{j,k}^{(n)}) \approx \bar{\varphi}_k^*(c_k(\theta_j^{(n)}))$. The term $\sum_m \bar{\varphi}_{m,k}^{(n)} \cdot w_k$ in the objective function is an approximation of the integral $\int_0^1 \bar{\varphi}_k(t) d\eta(t)$, and the term $\sum_j \bar{\psi}_{j,k}^{(n)} \cdot \mu_j^{(n)}$ approximates the integral $\int \bar{\varphi}_k^*(c_k(x)) d\mu$.

Remark 3. Lemma 14(b) provides a practical way of solving the problem $\mathcal{D}_{n,M}^*$: we need to solve the problem $\mathcal{D}'_{n,M}$ and extract the optimal values of the dual variables that correspond to the constraints **(mj-T)**, **(mj-J)**, and **(mj-E)**.

Next, we formulate a weak duality for the problem \mathcal{D}_n .

Proposition 13. Consider a family of functions $\varphi_k(x): [0, 1] \rightarrow \mathbb{R}$, $k \in \mathcal{I}$, a family of variables $c_{j,k}^{(n)} \geq 0$ satisfying the constraint **(c-def)**, and a family of variables $\psi_{j,k}^{(n)}$ satisfying the inequality

$$\varphi_k(t) + \psi_{j,k}^{(n)} \geq t \cdot c_{j,k}^{(n)} \quad \text{for all } t \in [0, 1], \quad 1 \leq j \leq n^I, \quad k \in \mathcal{I}.$$

Then

$$\max \mathcal{D}_n \leq \sum_{k=1}^{\mathcal{I}} \left(\int_0^1 \varphi_k(x) d\eta(x) + \sum_{j=1}^{n^I} \psi_{j,k}^{(n)} \cdot \mu_j^{(n)} \right).$$

Proof. Consider any family of variables $(u_j^{(n)}, p_j^{(n)})$ satisfying all the constraints of \mathcal{D}_n . First, by the **(c-def)**-constraint, we have

$$\sum_j \mu_j^{(n)} \cdot \left(\langle \theta_j^{(n)}, p_j^{(n)} \rangle - u_j^{(n)} \right) \leq \sum_j \mu_j^{(n)} \cdot \langle c_j^{(n)}, p_j^{(n)} \rangle = \sum_{k \in \mathcal{I}} \sum_{j=1}^{n^I} c_{j,k}^{(n)} \cdot p_{j,k}^{(n)} \cdot \mu_j^{(n)}.$$

Since $(u_j^{(n)}, p_j^{(n)})$ satisfy the **(mj)**-constraint, we can find a family of measures $\pi_k^{(n)}$ satisfying all the constraints of [Proposition 11\(d\)](#). We have

$$c_{j,k}^{(n)} \cdot p_{j,k}^{(n)} \cdot \mu_j^{(n)} \leq \int_T t \cdot c_{j,k}^{(n)} \, d\pi_k(j, t) \quad \Rightarrow \quad \sum_{j=1}^{n^I} c_{j,k}^{(n)} \cdot p_{j,k}^{(n)} \cdot \mu_j^{(n)} \leq \int_{J \times T} t \cdot c_{j,k}^{(n)} \, d\pi_k(j, t).$$

Next, it follows from the inequality $\varphi_k(t) + \psi_{j,k}^{(n)} \geq t \cdot c_{j,k}^{(n)}$ that

$$\begin{aligned} \int_{J \times T} t \cdot c_{j,k}^{(n)} \, d\pi_k(j, t) &\leq \int_{J \times T} \left(\varphi_k(t) + \psi_{j,k}^{(n)} \right) \, d\pi_k(j, t) = \\ &= \int_T \varphi_k(t) \, \text{dpr}_T \pi_k(t) + \int_J \psi_{j,k}^{(n)} \, \text{dpr}_J \pi_k(j). \end{aligned}$$

Finally, since $\text{pr}_T \pi_k \leq \eta$ and $\text{pr}_J \pi_k(j) \leq \mu_j^{(n)}$,

$$\int_T \varphi_k(t) \, \text{pr}_T \, d\pi_k(t) \leq \int_T \varphi_k(t) \, d\eta(t), \quad \int_J \psi_{j,k}^{(n)} \, \text{pr}_J \, d\pi_k(j) \leq \sum_{j=1}^{n^I} \psi_{j,k}^{(n)} \cdot \mu_j^{(n)}.$$

Summing up, we conclude that for all $(u_j^{(n)}, p_j^{(n)})$ satisfying all the constraints of \mathcal{D}_n , the following inequality holds:

$$\begin{aligned} \sum_j \mu_j^{(n)} \cdot \left(\langle \theta_j^{(n)}, p_j^{(n)} \rangle - u_j^{(n)} \right) &\leq \sum_{k=1}^{\mathcal{I}} \sum_{j=1}^{n^I} c_{j,k}^{(n)} \cdot p_{j,k}^{(n)} \cdot \mu_j^{(n)} \leq \\ &\leq \sum_{k \in \mathcal{I}} \int_{J \times T} t \cdot c_{j,k}^{(n)} \, d\pi_k(j, t) \leq \sum_{k \in \mathcal{I}} \left(\int_0^1 \varphi_k(x) \, d\eta(x) + \sum_{j=1}^{n^I} \psi_{j,k}^{(n)} \cdot \mu_j^{(n)} \right). \end{aligned}$$

□

Remark 4. In fact, the following strong duality holds:

$$\max \mathcal{D}_n = \min \sum_{k \in \mathcal{I}} \left(\int_0^1 \varphi_k(x) \, d\eta(x) + \sum_{j=1}^{n^I} \psi_{j,k}^{(n)} \cdot \mu_j^{(n)} \right).$$

This duality can be proven similarly to [Theorem 6](#).

Finally, using a solution to the dual problem, we can estimate how well the problem $\mathcal{D}_{n,M}$ approximates the problem \mathcal{D}_n .

Theorem 9. Denote $\varepsilon = \max_m |q_m - q_{m-1}|$. Then for all $\varepsilon \leq \frac{1}{6}$,

$$\max \mathcal{D}_{n,M} \leq \max \mathcal{D}_n \leq (1 + \varepsilon C) \cdot \max \mathcal{D}_{n,M},$$

where C is a constant that depends only on η and is independent of n , μ , and I .

Proof. Consider an optimal solution $(\bar{u}_j^{(n)}, \bar{p}_j^{(n)}, \bar{\pi}_{j,m,k}^{(n)})$ to the problem $\mathcal{D}_{n,M}$. By [Corollary 5](#), the variables $(\bar{u}_j^{(n)}, \bar{p}_j^{(n)})$ satisfy all the constraints of the problem \mathcal{D}_n ; therefore, since the objective functions of \mathcal{D}_n and $\mathcal{D}_{n,M}$ are identical,

$$\max \mathcal{D}_{n,M} \leq \max \mathcal{D}_n.$$

Let $(\bar{\varphi}_{m,k}^{(n)}, \bar{\psi}_{j,k}^{(n)}, \bar{c}_{j,k}^{(n)})$ be an optimal solution to $\mathcal{D}_{n,M}^*$. Denote

$$R_k = \sum_{m=1}^M \bar{\varphi}_{m,k}^{(n)} \cdot w_m + \sum_{j=1}^{n^I} \bar{\psi}_{j,k}^{(n)} \cdot \mu_j^{(n)}.$$

By the strong duality, $\sum_{k \in \mathcal{I}} R_k = \max \mathcal{D}_{n,M}$. By the **(lt)**-constraint,

$$\bar{c}_{j,k}^{(n)} \leq \frac{1}{t_m} \left(\bar{\psi}_{j,k}^{(n)} + \bar{\varphi}_{m,k}^{(n)} \right);$$

therefore, for each index $1 \leq m \leq M$, we have

$$\sum_{j=1}^{n^I} \bar{c}_{j,k}^{(n)} \cdot \mu_j^{(n)} \leq \frac{1}{t_m} \left(\sum_{j=1}^{n^I} \bar{\psi}_{j,k}^{(n)} \cdot \mu_j^{(n)} + \bar{\varphi}_{m,k}^{(n)} \right) \leq \frac{1}{t_m} \left(R_k + \bar{\varphi}_{m,k}^{(n)} \right). \quad (98)$$

As we mentioned in [Remark 2](#), the values $\bar{\varphi}_{m,k}^{(n)}$ are equal to the values of the non-decreasing function $\bar{\varphi}_k$ at the points t_m ; hence, the sequence $\{\bar{\varphi}_{m,k}^{(n)}\}_m$ is non-decreasing, and

$$R_k \geq \sum_{m=1}^M \bar{\varphi}_{m,k}^{(n)} \geq \bar{\varphi}_{m,k}^{(n)} \cdot (w_m + w_{m+1} + \dots + w_M) = \bar{\varphi}_{m,k}^{(n)} \cdot \eta([q_{m-1}, 1]).$$

Choose an index m such that $q_{m-1} \leq \frac{1}{2} < q_m$. Since $q_m - q_{m-1} \leq \frac{1}{6}$, we can estimate $t_m \geq q_{m-1} \geq \frac{1}{3}$. In addition, $\eta([q_{m-1}, 1]) \geq \eta([1/2, 1])$; therefore,

$$\sum_{j=1}^{n^I} \bar{c}_{j,k}^{(n)} \cdot \mu_j^{(n)} \leq \frac{1}{t_m} \left(R_k + \bar{\varphi}_{m,k}^{(n)} \right) \leq 3 \left(1 + \frac{1}{\eta([1/2, 1])} \right) R_k = C \cdot R_k,$$

where C is a constant that depends only on η .

For each $k \in \mathcal{I}$, consider a function $\widehat{\varphi}_k: [0, 1] \rightarrow \mathbb{R}$, $\widehat{\varphi}_k(t) = \overline{\varphi}_{m,k}^{(n)}$ for all $t \in [q_{m-1}, q_m)$, and $\widehat{\varphi}_k(1) = \overline{\varphi}_{M,k}^{(n)}$. In addition, consider a family of variables $\widehat{\psi}_{j,k}^{(n)}$ defined as follows:

$$\widehat{\psi}_{j,k}^{(n)} = \overline{\psi}_{j,k}^{(n)} + \varepsilon \cdot \overline{c}_{j,k}^{(n)}.$$

We claim that the **(It)**-constraint holds:

$$\widehat{\psi}_{j,k}^{(n)} + \widehat{\varphi}_k(t) \geq t \cdot \overline{c}_{j,k}^{(n)} \quad \text{for all } t \in [0, 1], 1 \leq j \leq n^I.$$

Indeed, for each $t \in [0, 1]$, there exists an index m such that $t_m \geq t - \varepsilon$ and $\widehat{\varphi}_k(t) = \overline{\varphi}_{m,k}^{(n)}$. Hence,

$$\widehat{\psi}_{j,k}^{(n)} + \widehat{\varphi}_k(t) = \overline{\psi}_{j,k}^{(n)} + \overline{\varphi}_{m,k}^{(n)} + \varepsilon \cdot \overline{c}_{j,k}^{(n)} \geq (t_m + \varepsilon) \cdot \overline{c}_{j,k}^{(n)} \geq t \cdot \overline{c}_{j,k}^{(n)}.$$

Thus $(\widehat{\varphi}_k, \widehat{\psi}_{j,k}^{(n)}, \overline{c}_{j,k}^{(n)})$ satisfies all the constraints of [Proposition 13](#); therefore,

$$\begin{aligned} \max \mathcal{D}_n &\leq \sum_{k \in \mathcal{I}} \left(\int_0^1 \widehat{\varphi}_k(t) \, d\eta(t) + \sum_{j=1}^{n^I} \widehat{\psi}_{j,k}^{(n)} \cdot \mu_j^{(n)} \right) \\ &= \sum_{k \in \mathcal{I}} \left(\sum_{m=1}^M \overline{\varphi}_{m,k}^{(n)} \cdot w_m + \sum_{j=1}^{n^I} \overline{\psi}_{j,k}^{(n)} \cdot \mu_j^{(n)} + \varepsilon \sum_{j=1}^{n^I} \overline{c}_{j,k}^{(n)} \cdot \mu_j^{(n)} \right) \\ &\leq \sum_{k \in \mathcal{I}} (R_k + \varepsilon C \cdot R_k) = (1 + \varepsilon C) \cdot \max \mathcal{D}_{n,M}. \end{aligned}$$

□

We conclude that the following convergence result holds.

Corollary 6. *Let $(\overline{u}_j^{(n,M)}, \overline{p}_j^{(n,M)}, \overline{\pi}_{j,m,k}^{(n,M)})$ be an optimal solution to the problem $\mathcal{D}_{n,M}$. Assume that $\max_{1 \leq m \leq M} |q_m - q_{m-1}| \rightarrow 0$ as $M \rightarrow \infty$. Then the sequence $\left\{ (\overline{u}_j^{(n,M)}, \overline{p}_j^{(n,M)}) \right\}_M$ contains a subsequence converging to an optimum of \mathcal{D}_n .*

E.3 Additional empirical optimizations

The total number of **(ic)**-constraints is n^{2I} , which constitutes the majority of all the constraints in the problem \mathcal{D}_n . We list some heuristics that allow to get rid of the redundant constraints improving the run time in practice.

Definition 4. A couple of points $\theta_i^{(n)}, \theta_j^{(n)} \in \Theta_n$ is called *irreducible* if the interval $(\theta_i^{(n)}, \theta_j^{(n)})$ does not contain any elements of Θ_n (by the interval $(\theta_i^{(n)}, \theta_j^{(n)})$ we mean the multi-dimensional linear segment with endpoints $\theta_i^{(n)}$ and $\theta_j^{(n)}$).

The following proposition shows that the incentive compatibility constraints from the problem \mathcal{D}_n can be verified only for irreducible couples of points.

Proposition 14. *Suppose that the inequality*

$$u_i^{(n)} - u_j^{(n)} \geq \langle \theta_i^{(n)} - \theta_j^{(n)}, p_j^{(n)} \rangle$$

holds for all couples of irreducible couple of points $\theta_i^{(n)}, \theta_j^{(n)}$. Then such an inequality holds for all $1 \leq i, j \leq n^I$.

Proof. Assume the converse and choose a pair of points $\theta_i^{(n)}, \theta_j^{(n)}$ such that

$$u_i^{(n)} - u_j^{(n)} < \langle \theta_i^{(n)} - \theta_j^{(n)}, p_j^{(n)} \rangle$$

and the interval $(\theta_i^{(n)}, \theta_j^{(n)})$ contains the smallest number of elements of Θ_n . Since the couple of points $\theta_i^{(n)}, \theta_j^{(n)}$ is not irreducible, there exists a point $\theta_k^{(n)} \in \Theta_n \cap (\theta_i^{(n)}, \theta_j^{(n)})$. By the construction, the following inequalities hold:

$$u_i^{(n)} - u_k^{(n)} \geq \langle \theta_i^{(n)} - \theta_k^{(n)}, p_k^{(n)} \rangle, \quad (99)$$

$$u_k^{(n)} - u_j^{(n)} \geq \langle \theta_k^{(n)} - \theta_j^{(n)}, p_j^{(n)} \rangle, \quad (100)$$

$$u_j^{(n)} - u_k^{(n)} \geq \langle \theta_j^{(n)} - \theta_k^{(n)}, p_k^{(n)} \rangle. \quad (101)$$

Combining inequalities (100) and (101), we conclude that

$$\langle \theta_j^{(n)} - \theta_k^{(n)}, p_j^{(n)} \rangle \geq \langle \theta_j^{(n)} - \theta_k^{(n)}, p_k^{(n)} \rangle \Rightarrow \langle \theta_j^{(n)} - \theta_k^{(n)}, p_j^{(n)} - p_k^{(n)} \rangle \geq 0.$$

Since $\theta_k^{(n)} \in (\theta_i^{(n)}, \theta_j^{(n)})$, there exists a constant $\gamma > 0$ such that $\theta_k^{(n)} - \theta_i^{(n)} = \gamma \cdot (\theta_j^{(n)} - \theta_i^{(n)})$; therefore,

$$\begin{aligned} \langle \theta_k^{(n)} - \theta_i^{(n)}, p_j^{(n)} - p_k^{(n)} \rangle &= \gamma \cdot \langle \theta_j^{(n)} - \theta_i^{(n)}, p_j^{(n)} - p_k^{(n)} \rangle \geq 0 \\ \Rightarrow \langle \theta_i^{(n)} - \theta_k^{(n)}, p_k^{(n)} \rangle &\geq \langle \theta_i^{(n)} - \theta_k^{(n)}, p_j^{(n)} \rangle. \end{aligned}$$

Thus, by inequality (99), we have

$$u_i^{(n)} - u_k^{(n)} \geq \langle \theta_i^{(n)} - \theta_k^{(n)}, p_k^{(n)} \rangle \geq \langle \theta_i^{(n)} - \theta_k^{(n)}, p_j^{(n)} \rangle.$$

Summing it up with (100), we conclude that $u_i^{(n)} - u_j^{(n)} \geq \langle \theta_i^{(n)} - \theta_j^{(n)}, p_j^{(n)} \rangle$. This contradiction proves the statement. \square

Remark 5. We state without a proof that the total number of irreducible couples of points is asymptotically equal to $\zeta(I)^{-1} \cdot n^{2I}$, where ζ is the Riemann zeta function. Informally, this can be show by the following argument. Rescale the points $\theta_j^{(n)}$ in such a way that the lattice $\{\theta_j^{(n)}\}_j$ coincides with the uniform integer lattice $\{1, 2, \dots, n\}^{\mathcal{I}}$. Then the couple of points $(\theta_i^{(n)}, \theta_j^{(n)})$ is irreducible if and only if

$$\gcd(\theta_{i,1}^{(n)} - \theta_{j,1}^{(n)}, \theta_{i,2}^{(n)} - \theta_{j,2}^{(n)}, \dots, \theta_{i,I}^{(n)} - \theta_{j,I}^{(n)}) = 1,$$

where \gcd denotes the greatest common divider.

For each prime p , consider the event

$$\gcd(\theta_{i,1}^{(n)} - \theta_{j,1}^{(n)}, \theta_{i,2}^{(n)} - \theta_{j,2}^{(n)}, \dots, \theta_{i,I}^{(n)} - \theta_{j,I}^{(n)}) \text{ is not divisible by } p. \quad (102)$$

For each k , the number p divides $\theta_{i,k}^{(n)} - \theta_{j,k}^{(n)}$ with the probability approximately equal to $\frac{1}{p}$. The events corresponding to different coordinate numbers k are mutually independent; therefore, the probability of the event (102) is approximately equal to $1 - p^{-I}$.

For a finite set of distinct prime numbers p , the events (102) can be considered as approximately mutually independent; therefore, the probability of the irreducibility of the couple $(\theta_i^{(n)}, \theta_j^{(n)})$ is approximately equal to

$$\prod_{p \text{ is prime}} (1 - p^{-I}) = \zeta(I)^{-1}.$$

In the case two items ($I = 2$), the total number of irreducible couples is approximately equal to $\frac{6}{\pi^2} n^4 \approx 0.61 n^4$; so, the heuristic described in Proposition 14 removes approximately 39% of all the (ic)-constraints.

Another heuristics that can be useful in practice is the following: we replace the global (ic)-constraints with the following local ones:

$$\begin{aligned} \text{(ic-local)} \quad u_i^{(n)} - u_j^{(n)} &\geq \langle \theta_i^{(n)} - \theta_j^{(n)}, p_j^{(n)} \rangle \\ &\text{for all irreducible } (\theta_i^{(n)}, \theta_j^{(n)}): \|\theta_i^{(n)} - \theta_j^{(n)}\| \leq \frac{c}{n}, \end{aligned}$$

where c is a small constant.

This definition is motivated by the notion of the directional convexity considered in Oberman (2013). Let V be a set of vectors v with integer coordinates such that $\|v\| \leq c$. Then $\|\theta_i^{(n)} - \theta_j^{(n)}\| \leq \frac{c}{n}$ if and only if $\theta_i^{(n)} - \theta_j^{(n)}$ is proportional to v for some $v \in V$.

Assume for simplicity that the function u defined on X is twice-differentiable. We say that u is *directionally convex* with respect to the set of direction vectors V if

$$\frac{\partial^2 u}{\partial v^2} \geq 0 \quad \text{for all } v \in V.$$

In other words, the function u is directionally convex if and only if the function $t \rightarrow u(x + tv)$ is convex for all $x \in X$ and all $v \in V$.

The constraint **(ic-local)** can be considered as a discrete version of the directional convexity; we state without the proof that if the sequence $\{(u^{(n)}, p^{(n)})\}_n$ satisfy the **(ic-local)** constraint, then the sequence $\{u^{(n)}\}_n$ contains a subsequence that converges weakly to the directionally convex function u .

[Oberman \(2013\)](#) proved that a directionally convex function is “nearly” convex.

Proposition (([Oberman, 2013](#), Proposition 3.1)). *Denote by $d\theta$ the directional resolution of the set V :*

$$d\theta = \max_{\|w\|=1} \min_{v \in V} \arccos \left(\frac{w \cdot v}{\|v\|} \right).$$

Assume that $d\theta \leq \frac{\pi}{4}$. Then every directionally convex function u is nearly convex, in the sense that

$$\frac{\lambda_1}{\lambda_I} \geq -\tan^2(d\theta),$$

where $\lambda_1 \leq \lambda_2 \leq \dots \leq \lambda_I$ are the eigenvalues of the Hessian matrix at the point x .

If $c \rightarrow \infty$, then $d\theta \rightarrow 0$, and the nearly convex function u becomes a convex one. This justifies the convenience of the suggested approach. In practice, an iterative [Algorithm 1](#) can be used: we start with a small number of initial **(ic)**-constraints, and, at each step, we add all the violated constraints to the linear program.

Finally, we improve the convergence rate of the majorization constraint approximation obtained in [Theorem 9](#) for the case of bounded optimal solution $\bar{c}_{j,k}^{(n)}$ of the dual problem $\mathcal{D}_{n,M}^*$.

Proposition 15. *Assume that the partition $0 = q_0 < q_1 < \dots < q_M = 1$ satisfies the following property: for each function $\varphi: [0, 1] \rightarrow \mathbb{R}$ which is constant on the intervals $[0, t_1]$, $[t_M, 1]$ and linear on each interval $[t_m, t_{m+1}]$ we have*

$$\int_0^1 \varphi(t) \, d\eta(t) = \sum_{m=1}^M \varphi(t_m) \cdot w_m.$$

Let $(\bar{\varphi}_{m,k}^{(n)}, \bar{\psi}_{j,k}^{(n)}, \bar{c}_{j,k}^{(n)})$ be an optimal solution to $\mathcal{D}_{n,M}^$. Then*

$$\max \mathcal{D}_n \leq \max \mathcal{D}_{n,M} + \int_{t_M}^1 (t - t_M) \, d\eta(t) \cdot \sum_{k \in \mathcal{I}} \max_j \bar{c}_{j,k}^{(n)}.$$

Remark 6. For the case of $B = 2$ bidders, the distribution η is uniform on the interval $[0, 1]$, and the uniform partition $q_m = m/M$ satisfies the restriction of

Algorithm 1: Iterative scheme with the local IC-constraints

Define the problem $\mathcal{D}_{n,M}^{(0)}$, where **(ic)**-constraints are replaced with **(ic-local)**;

for $s = 0, 1, \dots$ **do**

 find an optimal solution $(\bar{u}^{(n)}, \bar{p}^{(n)})$ of the problem $\mathcal{D}_{n,M}^{(s)}$;
 define the set $V^{(s)}$ of all the violated constraints:

$$V^{(s)} = \left\{ (i, j) : \bar{u}_i^{(n)} - \bar{u}_j^{(n)} < \langle \bar{p}_j, \theta_i^{(n)} - \theta_j^{(n)} \rangle \right\}$$

if $V^{(s)}$ *is not empty* **then**

 define $\mathcal{D}_{n,M}^{(s+1)} = \mathcal{D}_{n,M}^{(s)}$;

 add all the violated constraints to the problem $\mathcal{D}_{n,M}^{(s+1)}$:

$$u_i^{(n)} - u_j^{(n)} \geq \langle p_j, \theta_i^{(n)} - \theta_j^{(n)} \rangle \quad \text{for all } (i, j) \in V^{(s)}$$

else

 stop with the solution $(\bar{u}^{(n)}, \bar{p}^{(n)})$ of the problem $\mathcal{D}_{n,M}^{(s)}$;

end

end

Proposition 15. In the case of $B > 2$, such a partition can be efficiently found numerically. We checked numerically that for such a partition we have $t_M = 1 - O(M^{-1})$; therefore,

$$\int_{t_M}^1 (t - t_M) \, d\eta(t) = O(M^{-2}).$$

So, if $\bar{c}_{j,k}^{(n)}$ is uniformly bounded on M , the constructed partition provides a quadratic convergence rate (compared with the linear convergence rate obtained in [Theorem 9](#)).

Proof of Proposition 15. The proof is based on the same ideas as the proof of [Theorem 9](#). Define

$$\bar{\varphi}_k(t) = \max \left\{ 0, \max_j \left(t \cdot \bar{c}_{j,k}^{(n)} - \bar{\psi}_{j,k}^{(n)} \right) \right\}, \quad t \in [0, 1].$$

The function $\bar{\varphi}_k(t)$ is convex, non-decreasing, and non-negative. In addition,

$$\bar{\varphi}_k(t_m) = \bar{\varphi}_{m,k}^{(n)} \quad \text{for all } 1 \leq m \leq M.$$

Besides, by the construction, this function satisfies the **(It)**-constraint:

$$\bar{\varphi}_k(t) + \bar{\psi}_{j,k}^{(n)} \geq t \cdot \bar{c}_{j,k}^{(n)}.$$

For each k , let $\hat{\varphi}_k(x)$ be a unique function that is constant on the intervals $[0, t_1]$, $[t_M, 1]$, linear on each interval $[t_m, t_{m+1}]$, and equal to $\bar{\varphi}_{m,k}^{(n)}$ at the point t_m for all $1 \leq m \leq M$. We claim that $\hat{\varphi}_k(t) \geq \bar{\varphi}_k(t)$ for all $t \in [0, t_M]$. Indeed, for each m , the function $\hat{\varphi}_k(t)$ is linear on the interval $[t_m, t_{m+1}]$, the function $\bar{\varphi}_k(t)$ is convex on the same interval, and

$$\bar{\varphi}_k(t_m) = \hat{\varphi}_k(t_m), \quad \bar{\varphi}_k(t_{m+1}) = \hat{\varphi}_k(t_{m+1}).$$

Thus it follows from Jensen's inequality that $\hat{\varphi}_k(t) \geq \bar{\varphi}_k(t)$ for all $t \in [t_m, t_{m+1}]$. Finally, the function $\bar{\varphi}_k(t)$ is non-decreasing on the interval $[0, t_1]$, the function $\hat{\varphi}_k(t)$ is constant on $[0, t_1]$, and $\hat{\varphi}_k(t_1) = \bar{\varphi}_k(t_1)$; therefore, $\hat{\varphi}_k(t) \geq \bar{\varphi}_k(t)$ for all $t \in [0, t_1]$.

By the construction, the derivative of $\bar{\varphi}_k(t)$ cannot exceed the maximal slope in the family linear functions $t \cdot \bar{c}_{j,k}^{(n)} - \bar{\psi}_{j,k}^{(n)}$. Denoting by $C_k = \max_j \bar{c}_{j,k}^{(n)}$, we conclude that for all $t \in [t_M, 1]$:

$$\bar{\varphi}_k'(t) \leq C_k \quad \Rightarrow \quad \bar{\varphi}_k(t) \leq \bar{\varphi}_k(t_M) + C_k \cdot (t - t_M) = \hat{\varphi}_k(t) + C_k \cdot (t - t_M).$$

Thus $\hat{\varphi}_k(t) + C_k \cdot [t - t_M]_+ \geq \bar{\varphi}_k(t)$ for all $t \in [0, 1]$.

Finally, it follows from [Proposition 13](#) that

$$\begin{aligned}
\max \mathcal{D}_n &\leq \sum_{k \in \mathcal{I}} \left(\int_0^1 (\widehat{\varphi}_k(t) + C_k \cdot [t - t_M]_+) \, d\eta(t) + \sum_{j=1}^{n^I} \overline{\psi}_{j,k}^{(n)} \cdot \mu_j^{(n)} \right) = \\
&= \sum_{k \in \mathcal{I}} \left(\sum_{m=1}^M \overline{\varphi}_{m,k}^{(n)} \cdot w_m + C_k \int_{t_M}^1 (t - t_M) \, d\eta(t) + \sum_{j=1}^{n^I} \overline{\psi}_{j,k}^{(n)} \cdot \mu_j^{(n)} \right) = \\
&= \max \mathcal{D}_{n,M} + \sum_{k \in \mathcal{I}} C_k \cdot \int_{t_M}^1 (t - t_M) \, d\eta(t).
\end{aligned}$$

□

F Beckmann's problem, congested transport, and dynamic viewpoint

Beckmann's problem is equivalent to a Monge-Kantorovich-type problem called “congested optimal transport”; see ([Santambrogio, 2015](#)) for the detailed presentation and references. Let us describe the equivalence informally for Beckmann's problem with the weight $\rho \equiv 1$. Given a domain Ω of a Euclidean space, an absolutely continuous supply-demand imbalance measure π on Ω satisfying $\pi(\Omega) = 0$, and a convex function Φ , the following identity holds:

$$\inf_{c: \operatorname{div}[c] + \pi = 0} \int \Phi(c) \, dx = \inf_{Q: Q_1 - Q_0 + \pi = 0} \int_{\Omega} \Phi(i_Q) \, dx. \quad (103)$$

Here Q ranges over to the set of all probability measures on “curves”, i.e., continuous mappings $\gamma: [0, 1] \rightarrow \Omega$, and Q_t denotes the probability measure on Ω obtained as the image of Q under the map $\gamma \rightarrow \gamma(t) \in \Omega$. The object i_Q is the so-called traffic intensity function which is defined so that the following identity holds for any test function φ :

$$\int_{\Omega} \varphi(x) i_Q(x) \, dx = \int \left(\int_0^1 \varphi(\gamma(t)) |\gamma'(t)| \, dt \right) dQ(\gamma)$$

Formula (103) can be seen as a Lagrangian formulation of Beckmann's problem taking a form of a “problem for measures on curves.” It is a quasi-dynamical formulation, a version of which is well known for the Monge-Kantorovich transportation problem; see ([Villani, 2009](#)).

The problem (103) is equivalent to a version of the Monge-Kantorovich problem with the cost function depending on optimal Q (or c); see ([Santambrogio, 2015](#), Theorem 4.33). This form justifies the term “congested optimal transport.”

Finally, let us mention that the construction of Q relies on the so-called Dacorogna–Moser ([Santambrogio, 2015](#)) interpolation of probability measures, which is a solution to the following transport equation:

$$\frac{\partial}{\partial t}\rho + \operatorname{div}\left[\frac{c}{(1-t)f_+ + tf_-}\rho_t\right] = 0, \quad \rho_0 = f_+,$$

where f_+ and f_- are the densities of the positive and the negative components of $\pi = \pi_c - \pi_p$ with respect to the Lebesgue measure.

On the other hand, to the best of our knowledge, there is no natural variational / dynamical interpretation of congested optimal transport in the spirit of the Benamou–Brenier formula (“problem for curves of measures”; see [Villani \(2009\)](#); [Santambrogio \(2015\)](#)); see also remarks in Section 4.5 of ([Carlier, 2012](#)).